

Chapter 1: Conceptual Basis

Section 1

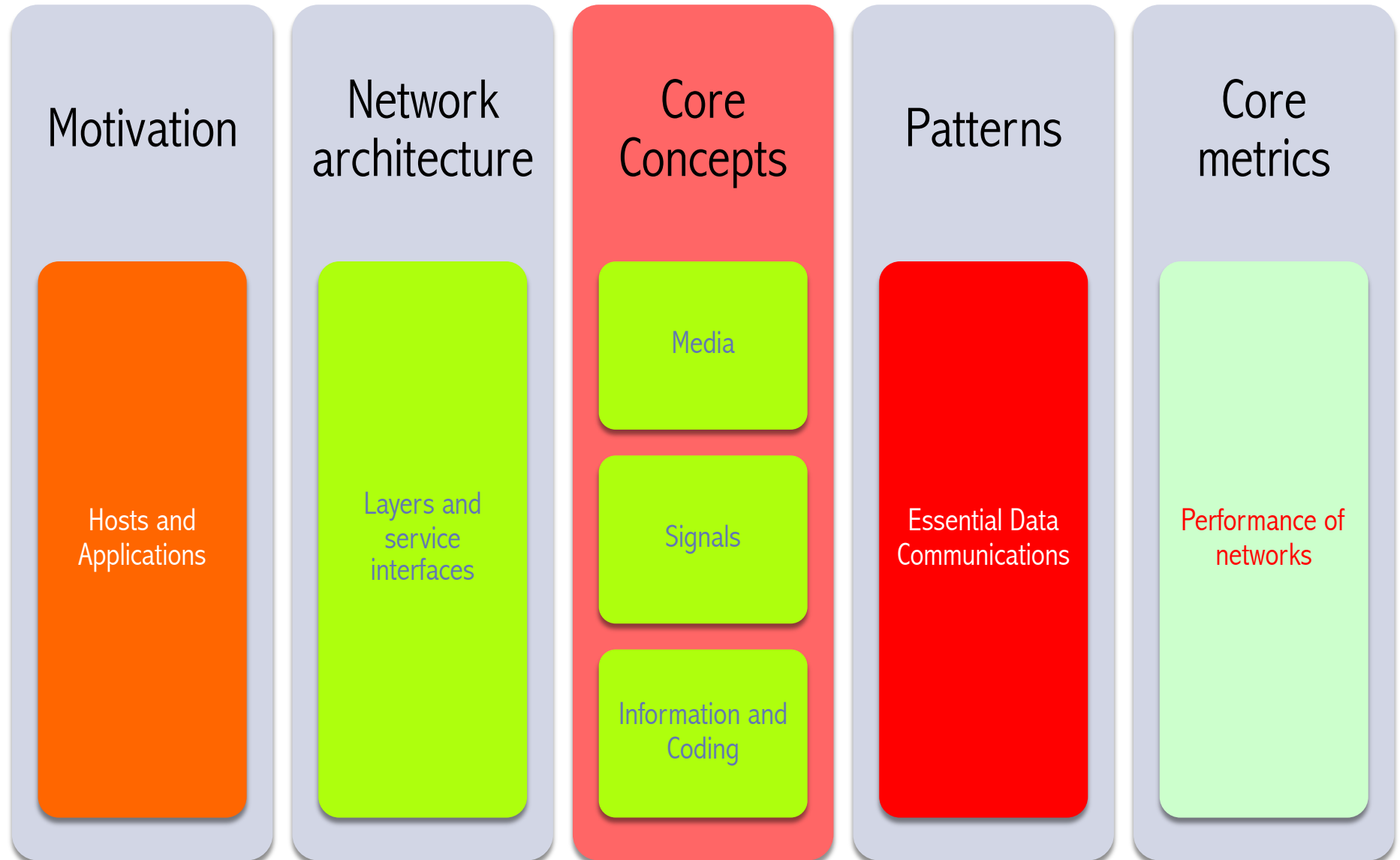
Leading questions

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- ☐ What are the principles guiding the communication between two parties?
- ☐ When can a communication be considered fast and efficient?
- ☐ What are the landmarks about the development of Internet?
- ☐ Why is networking essential for progress?
- ☐ What is a network architecture?

Flow of topics

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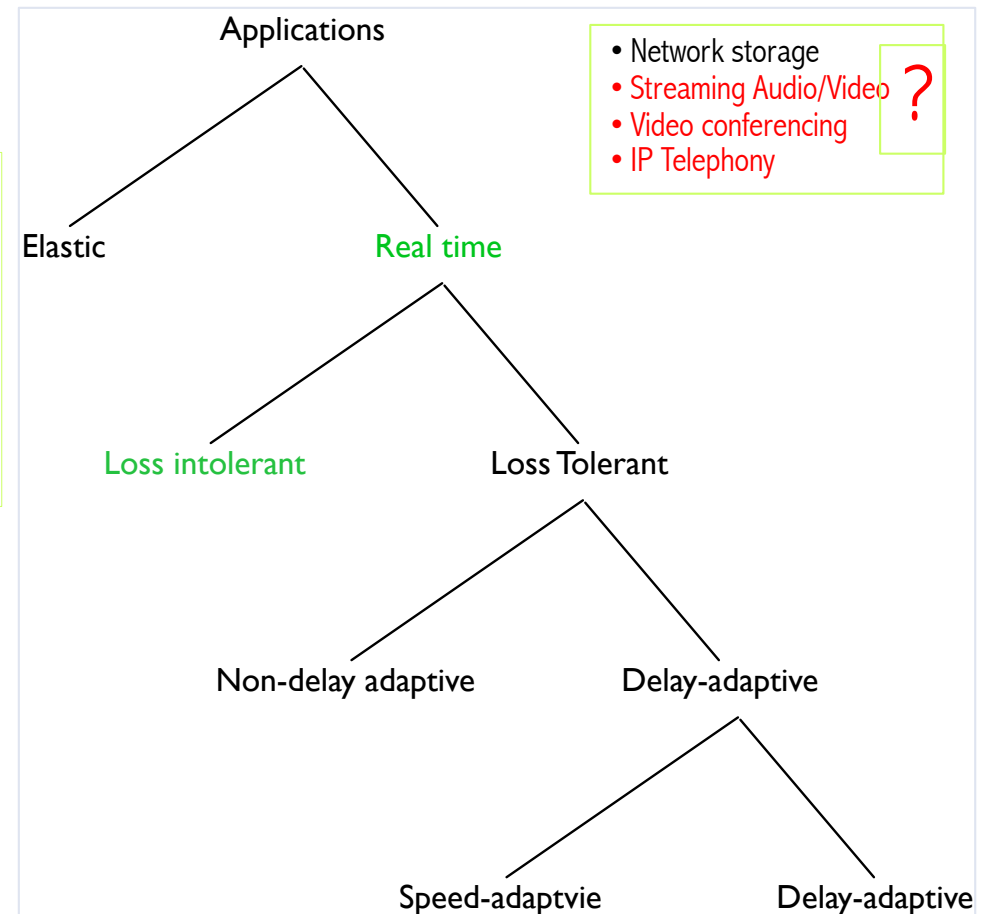


Cooperating host applications

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- Applications are computer programs
- Communicate over the Internet
- At work, at home and mobile
- Vastly differing requirements

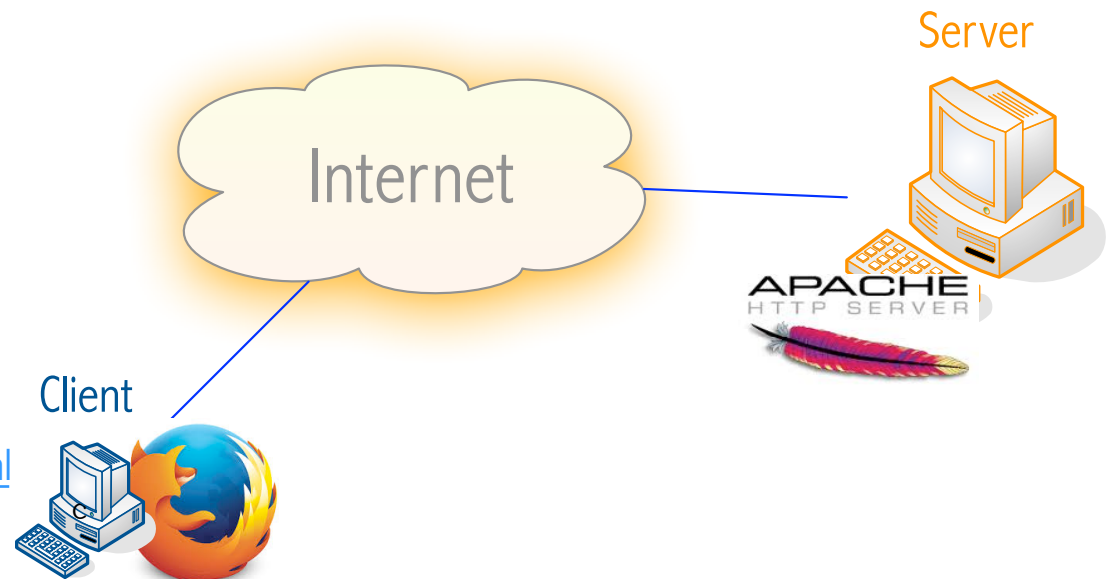
- e-mail: smtp, pop3, imap
- File Sharing, file transfer: ftp, rcp, scp
- Printer sharing
- Virtual terminal: telnet, ssh (Secure Shell)
- e-commerce
- Geolocation
- **World Wide Web (www)**
- Social Networks
- Instant Messaging (Whatsup, ...)



Essential Internet service: www

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- Web pages are downloaded by the client from the **server**
 - Client and server speak the **http** protocol
 - http = Hyper Text Transfer Protocol
- **www** = World Wide Web:
 - A distributed, Client/Server application
 - **Server** program (e.g., Apache)
 - **Client** program (e.g., Firefox)
- **URL**
 - Uniform Resource Locator
 - <http://palauto.unileon.es/cn/index.html>
- HTTP, in turn uses the **TCP** protocol for reliability
 - TCP = Transmission Control Protocol
 - TCP provides reliability
 - In case of packet loss, duplication, errors, etc



Units and multipliers

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□ Bandwidth

- ▣ Directly related to the acceptable speed of bit transmission over some medium
- ▣ Number of bits transmitted in one second:
 - ▣ Bps (Bits Per Second = Bits/Sec)
- ▣ Since bandwidth is a rate, the multipliers take on the following values:
 - ▣ K (Kilo = 10^3)
 - ▣ M (Mega = 10^6)
 - ▣ G (Giga = 10^9)
 - ▣ T (Tera = 10^{12})

□ Delay

- ▣ Seconds
- ▣ How much time it takes to transport one bit from a source to a destination directly connected
- ▣ Propagation delay

□ Jitter

- ▣ The variance of the delay

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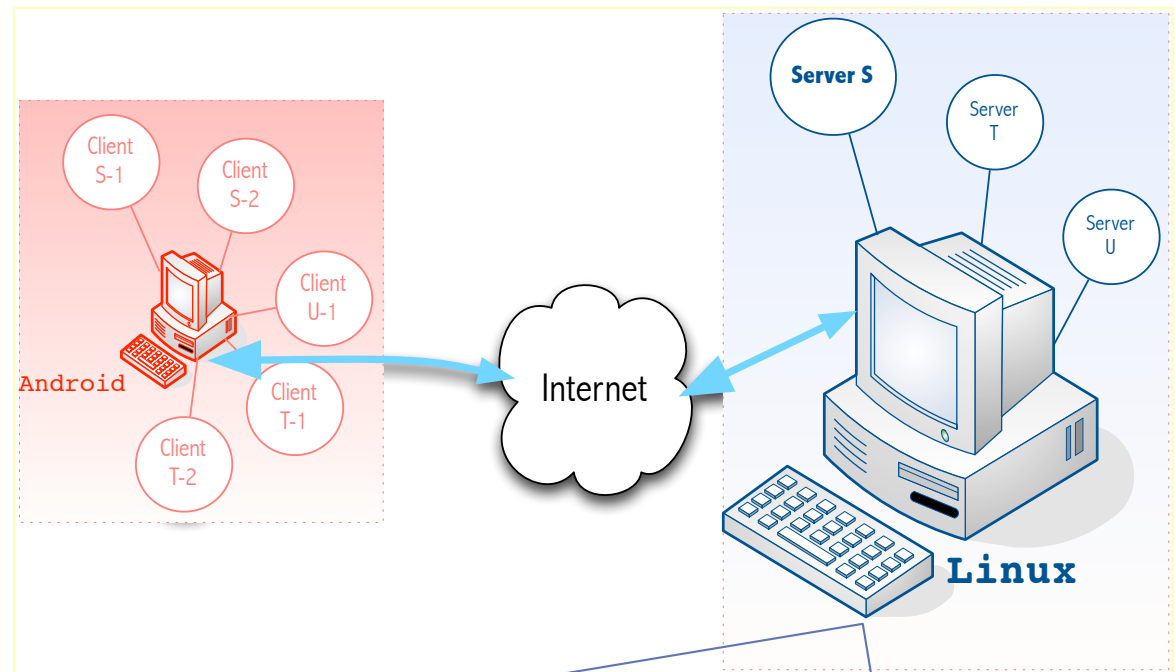
Network Architecture

Manage the complexity of networks

Logical channels

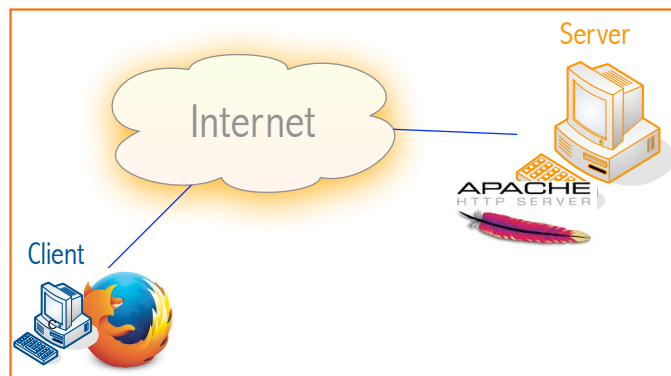
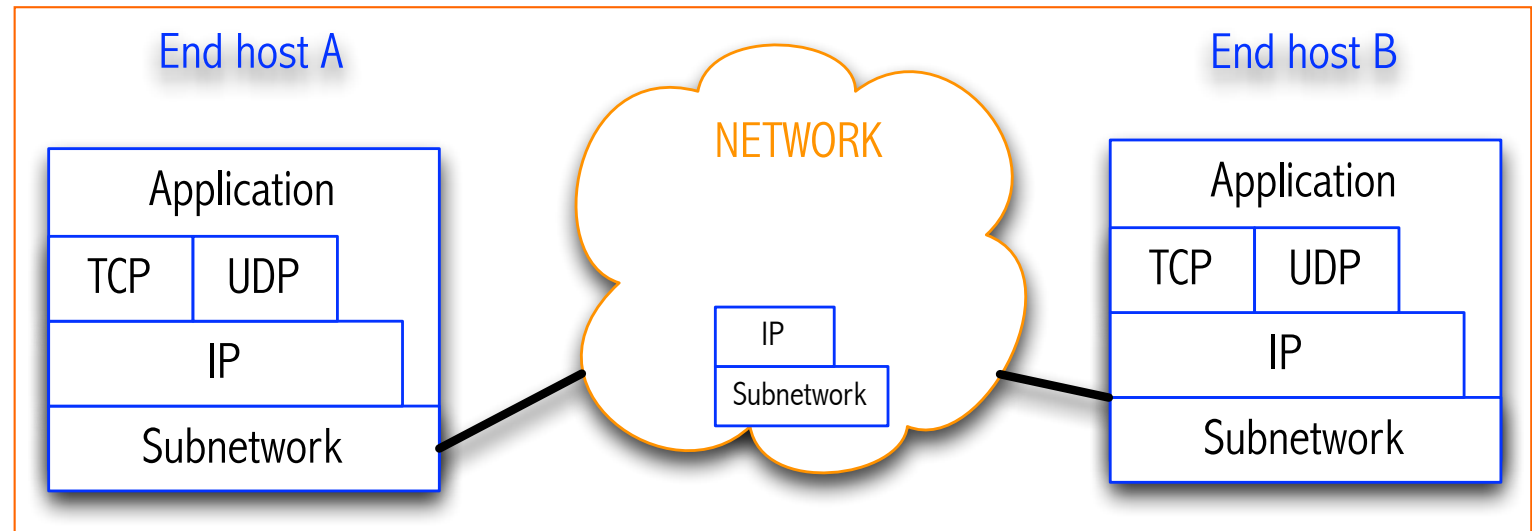
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- Applications communicate over the Internet
- The channel between two communicating applications is logical
- Each channel:
 - ▣ *Connects* two applications
 - ▣ Hosts must be identified:
 - IP address
 - ▣ Applications must be identified:
 - Port numbers



How to **hide** network **complexity** from application programmers?

Layering in hosts and network



Internet Architecture

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- Network complexity is organized into 4 layers

- Each *layer*

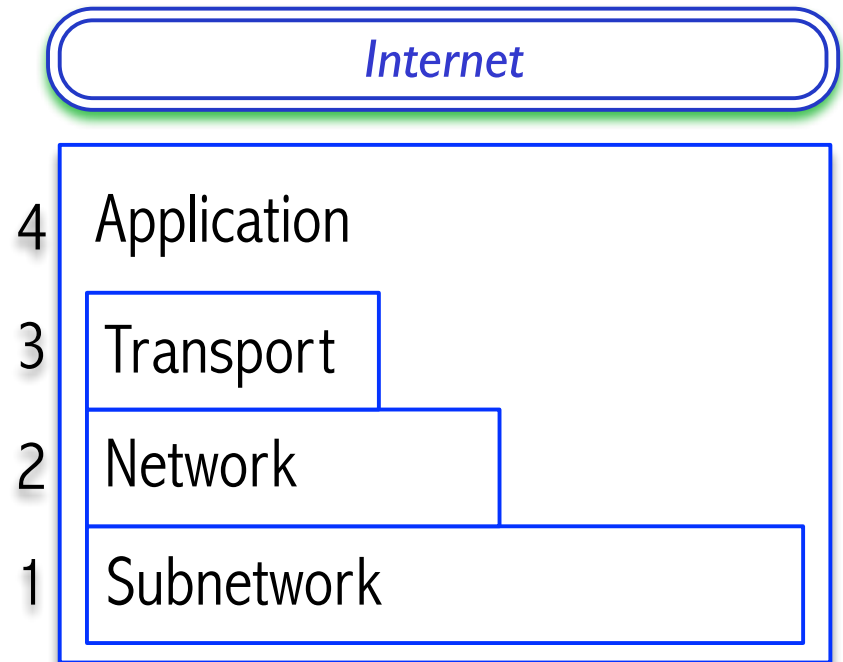
- Offers a set of services to the upper layers
 - The mechanism that attains each *service* is a protocol
 - An upper layer avails one service from a lower layer by calling its *interface*

- 1. Subnetwork: Ethernet, Wi-Fi, Bluetooth

- 2. Network: Only IP !!!

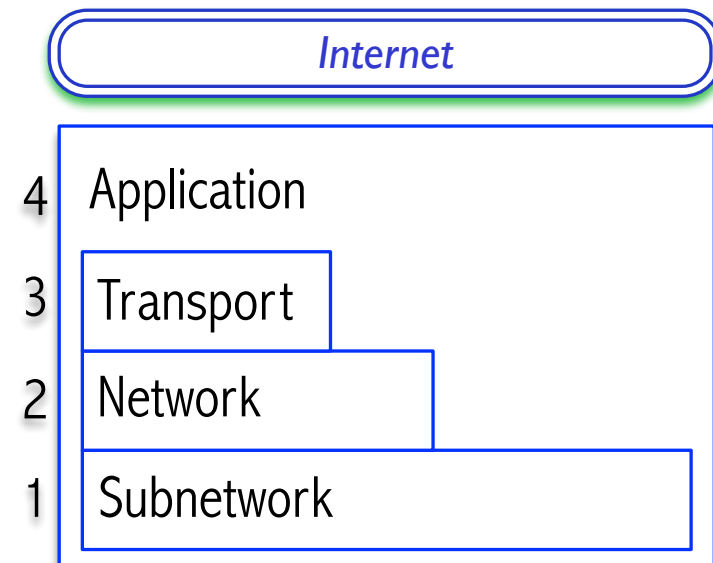
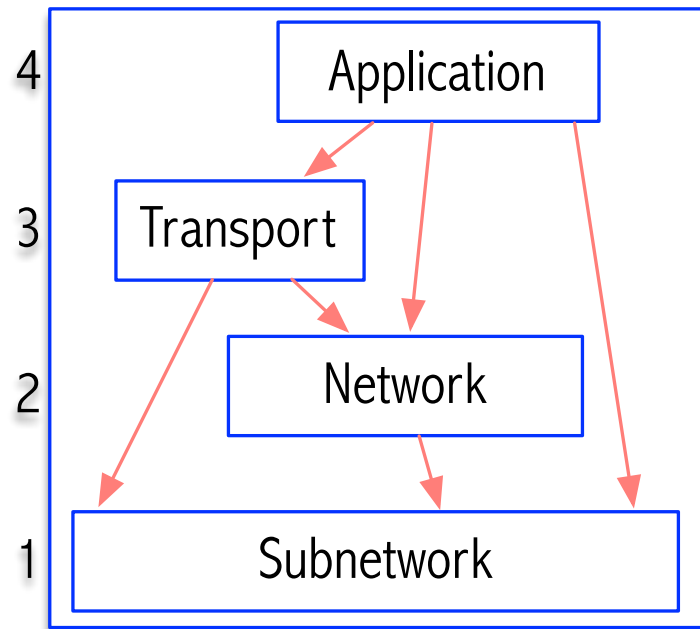
- 3. Transport: TCP and UDP

- 4. Application: Whatsup and innumerably others

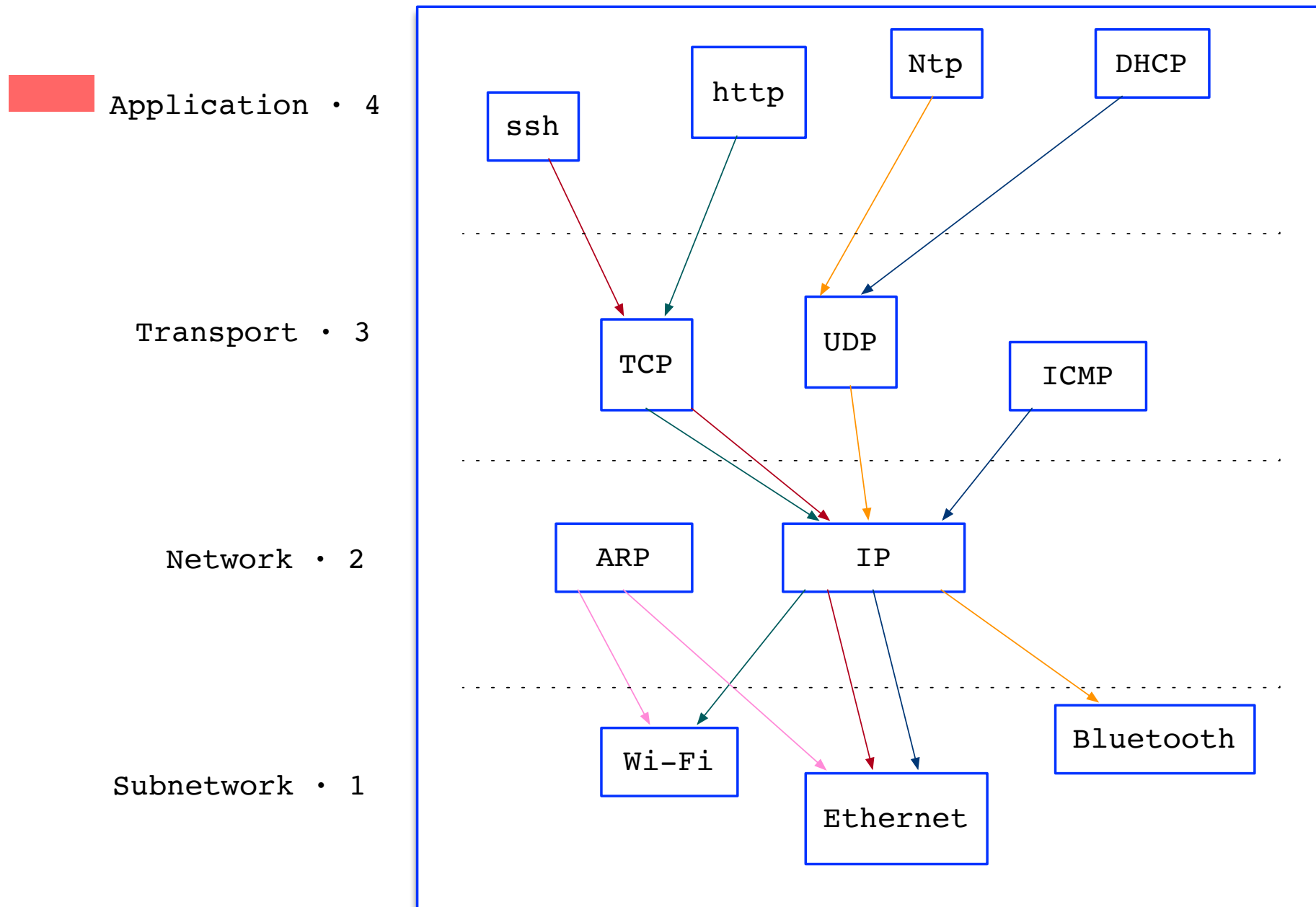


Internet Architecture

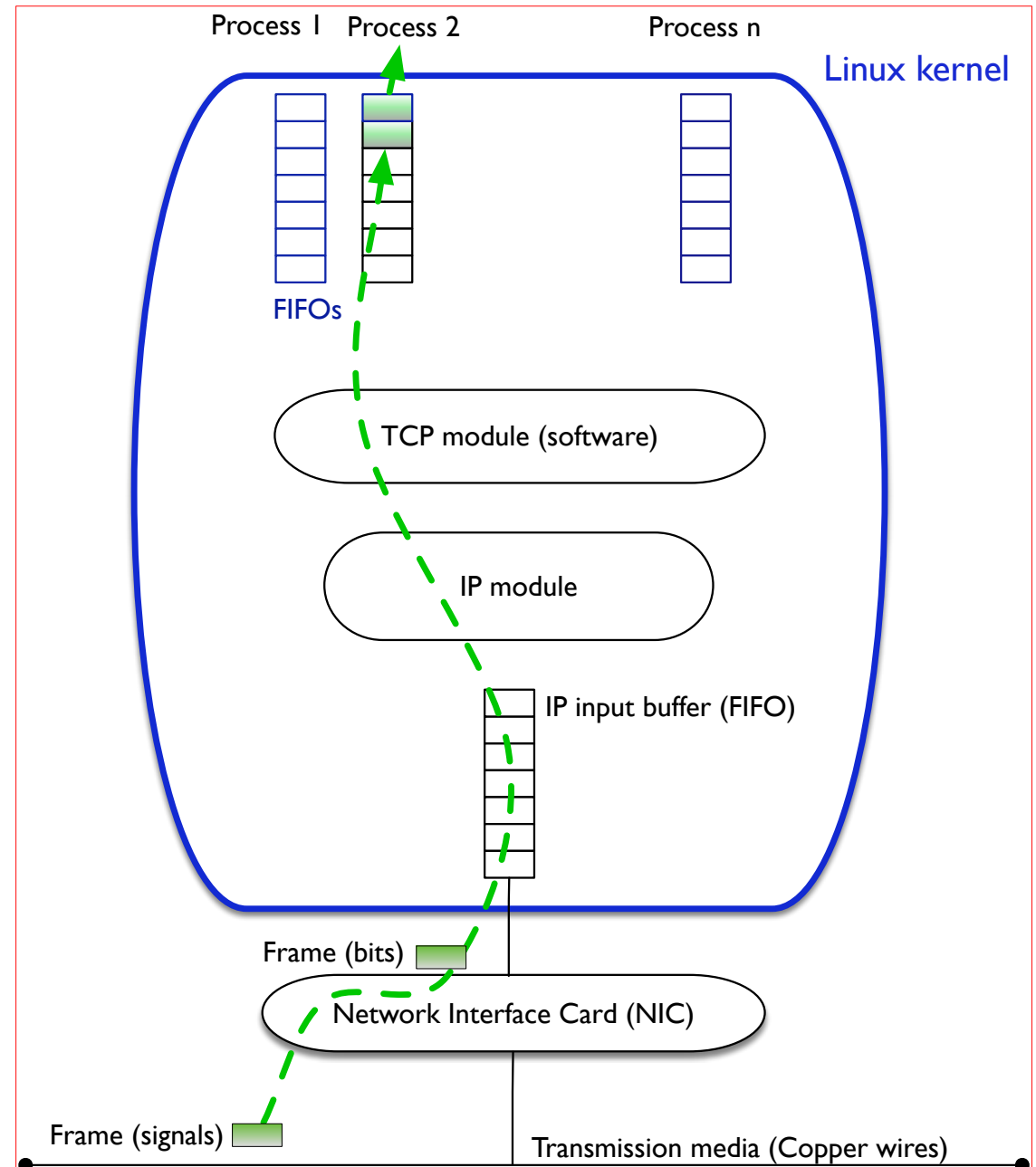
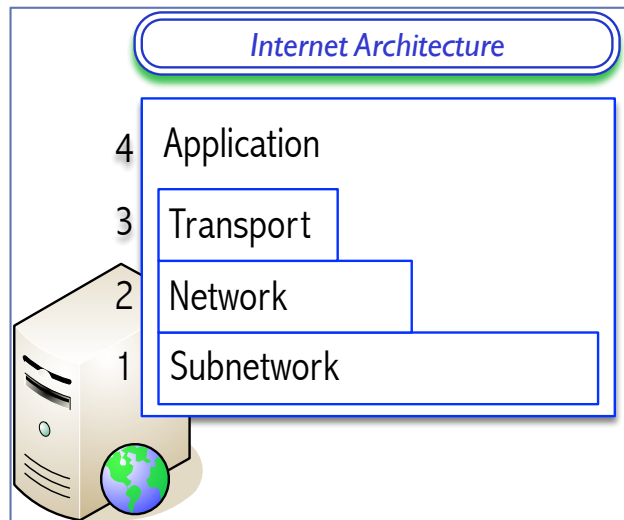
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Typical Internet Protocol Stack



Implementation of protocols



Internet Architecture

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- Specified by Internet Engineering Task Force in 1970
- RFC 1287
- Derived from the TCP/IP Protocol suite
 - ▣ In any implementation of IE programs can call any layer's service interface
 - ▣ Only one network protocol: IP
 - Many application protocols
 - A few transports
 - Many link protocols
 - A **glass-shaped architecture**
 - ▣ IETF accepts protocols officially if they offer one reference implementation at least

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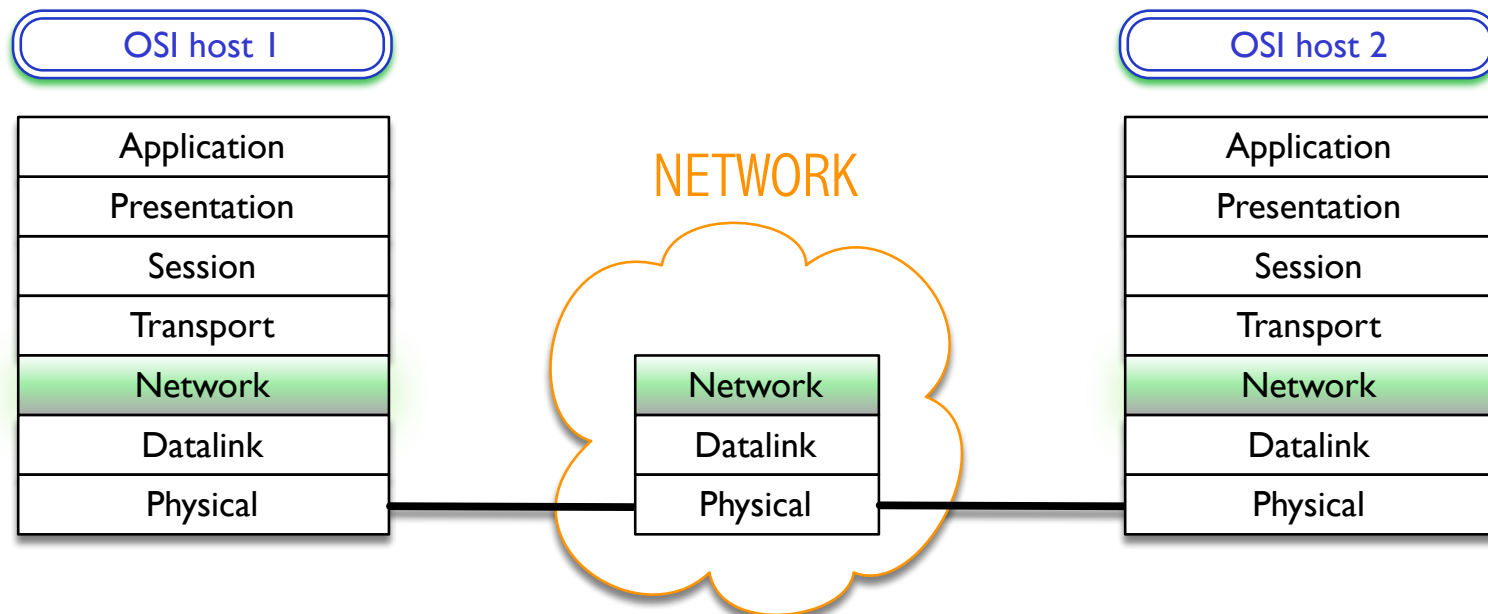
Protocols and their Services

Protocols offer services

7-layer OSI Architecture

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- A Reference Model
- Use of this architecture is limited to some specific protocols
 - ▣ IEEE 802.3, 802.1Q, 802.1P, 802.11 (Wifi), 802.15 (WiMax)
- Layering is strict



Description of OSI Layers

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□ Physical Layer

- ▣ Handles the transmission of raw bits over a communication link

□ Data Link Layer

- ▣ Collects a stream of bits into a larger aggregate called a *frame*
- ▣ Network adaptor along with device driver in OS implement the protocol in this layer
- ▣ Frames are actually delivered to hosts

□ Network Layer

- ▣ Handles routing among nodes within a packet-switched network
- ▣ Unit of data exchanged between nodes in this layer is called a *packet*

The lower three layers are implemented **on all network nodes**

OSI Architecture

Application
Presentation
Session
Transport
Network
Datalink
Physical

Description of OSI Layers

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- **Transport Layer**
 - ▣ Implements a process-to-process channel
 - ▣ Unit of data exchanges in this layer is called a *message*
- **Session Layer**
 - ▣ Provides a name space that is used to tie together the potentially different transport streams that are part of a single application
- **Presentation Layer**
 - ▣ Concerned about the format of data exchanged between peers
- **Application Layer**
 - ▣ Standardize common type of exchanges

OSI Architecture

Application
Presentation
Session
Transport
Network
Datalink
Physical

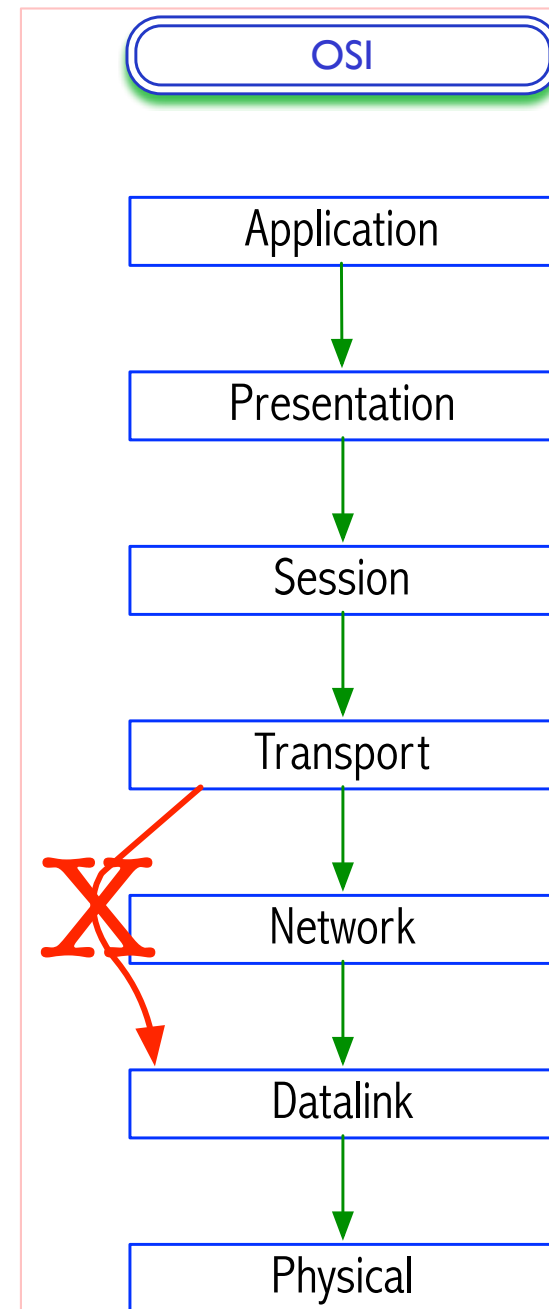
The transport layer and the higher layers typically run only **on end-hosts** and not on the intermediate switches and routers

OSI, strict layering

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A layer only uses the services provided by the layer below

- The internal mechanisms of each layer remain hidden
 - ▣ Layer N+1 knows nothing about the internal mechanisms of layer N
- Example:
 - ▣ Transport layer can only use the Network layer

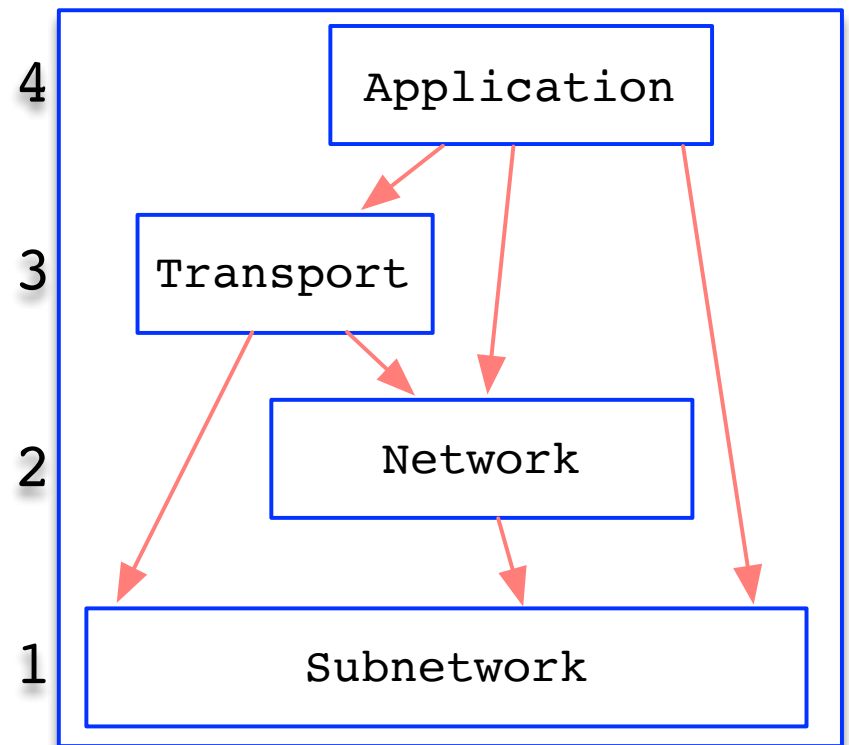


Layering in Internet, non-strict

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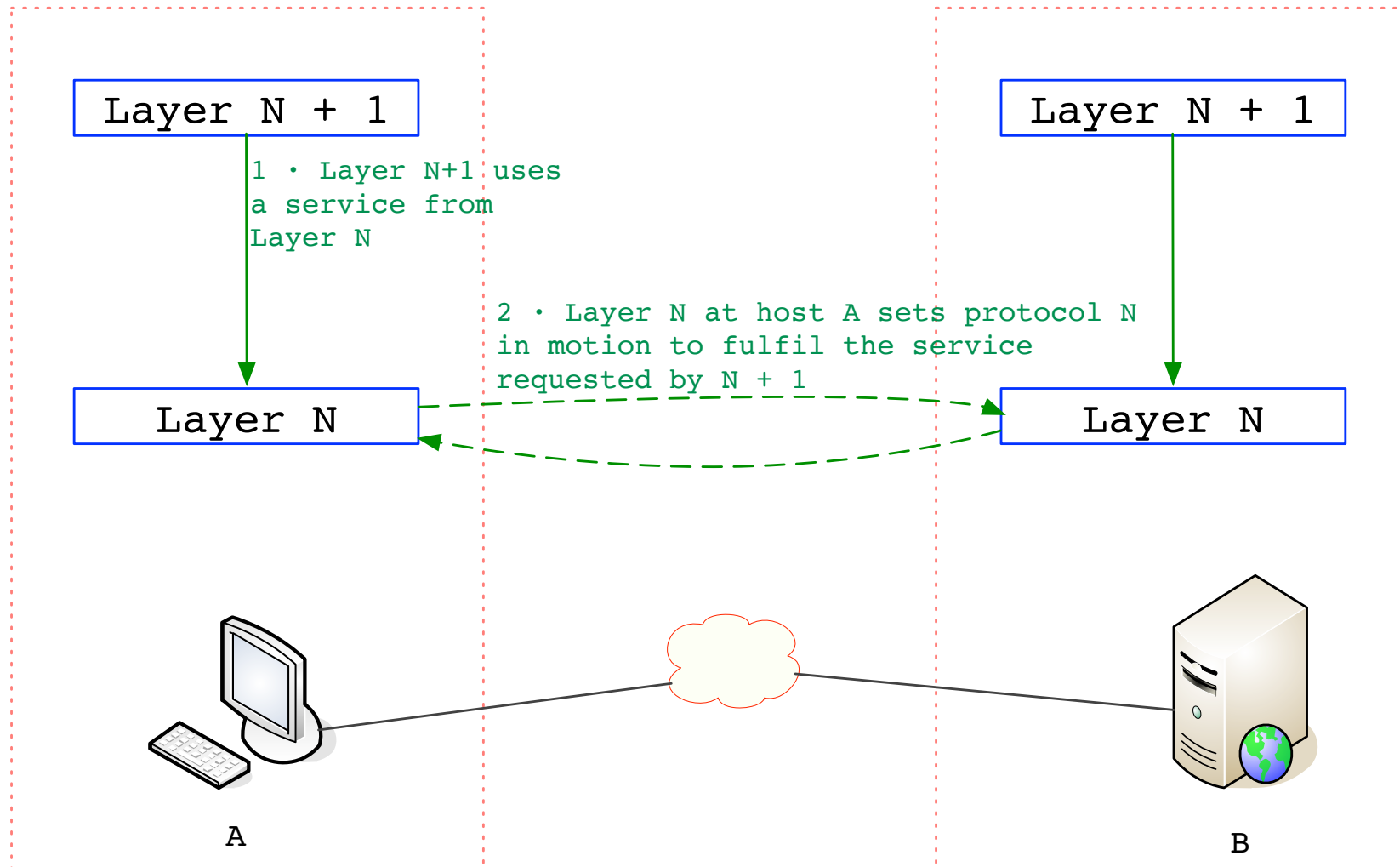
A layer may use the services provided by *any* layer below

- The internal mechanisms of each layer remain hidden
 - ▣ Layer N+1 knows nothing about the internal mechanisms of layer N
- Example:
 - ▣ An Application protocol may use whichever lower layer



Protocol: The foreman of a service

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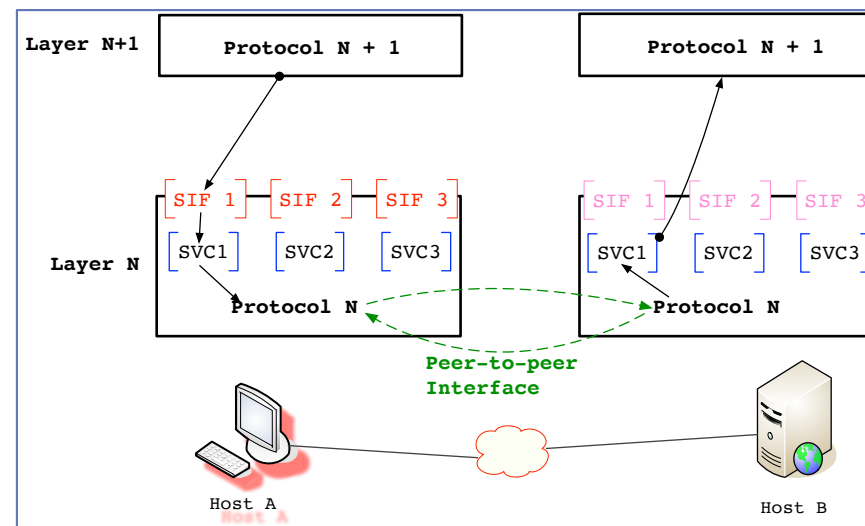
N+1 represents any upper layer in IA (Internet Architecture)

Layer N+1 uses a service at Layer N

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□ Layer N

- ▣ Several services: SVC1, SVC2
- ▣ Each service is accessed through its **Service Interface: SIF1, SIF2**
- ▣ The protocol N (Host A) fulfils the functionality offered by SVC by exchanging messages with protocol N at Host B
- ▣ These messages comprise the **Peer-to-Peer Interface**

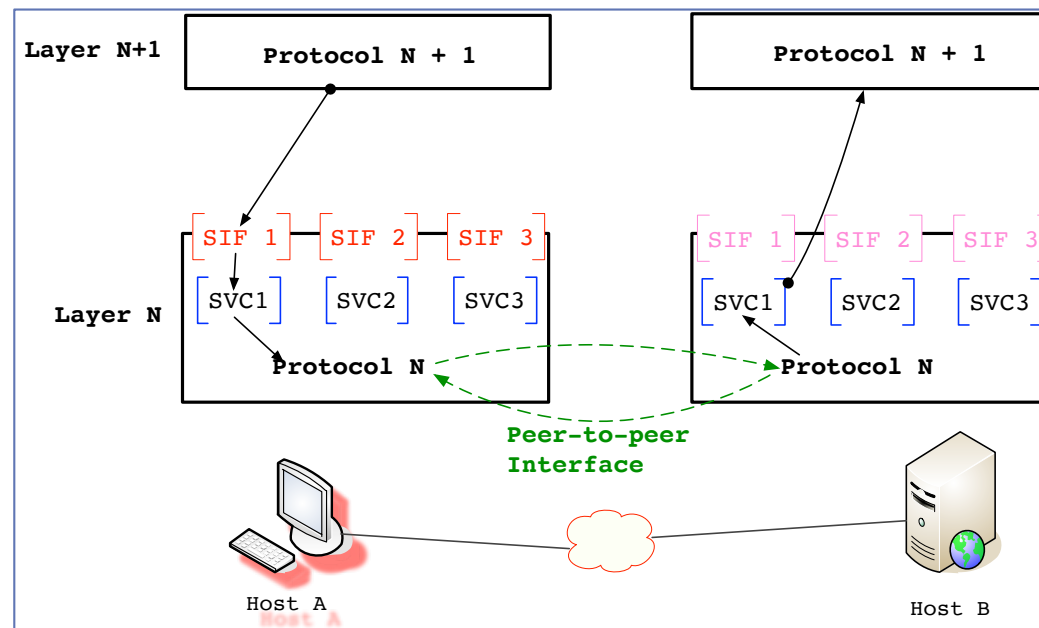


N+1 represents any upper layer in IA (Internet Architecture)

Example: A runs Linux; B runs Windows

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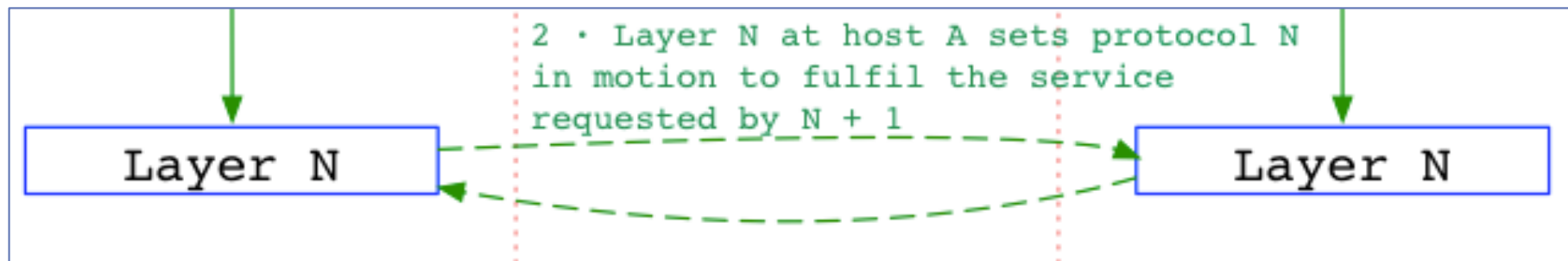
- Equal layers at A and B must implement the **same protocol**
 - ▣ *Same peer-to-peer interface*
- However, **Service Interfaces** at A and B might be present differences



Peer-to-peer interface

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- The syntax and the semantics of the messages exchanged by the two peers must follow a formal specification
 - ▣ ASN.1, Abstract Syntax Notation
- Normally, we refer to the peer-to-peer interface with the same word: protocol
- Protocols of Internet are specified by the IETF
 - ▣ RFC: Request For Comments
 - ▣ Example: The ICMP protocol is specified in RFC 792



Encapsulation and Multiplexing: *Concept*

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- What information is sent from N+1 to N through the SIF (Service Interface)?
 - ▣ Protocol N+1 sends a N+1 Data Unit to Protocol N
 - ▣ Protocol N encapsulates the N+1 Data Unit into a fresh N Data Unit:
 - Payload(N+1) + Header(N)
 - This scheme is reproduced at each service use
 - ▣ **Data Unit:** A bit string produced by a protocol
 - ▣ **Encapsulation:** Appending a Header to a Data Unit

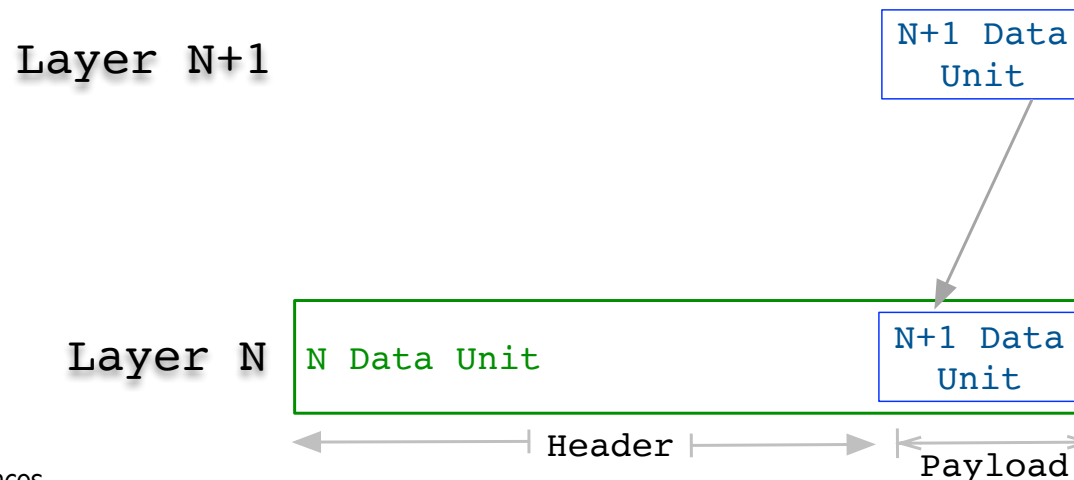
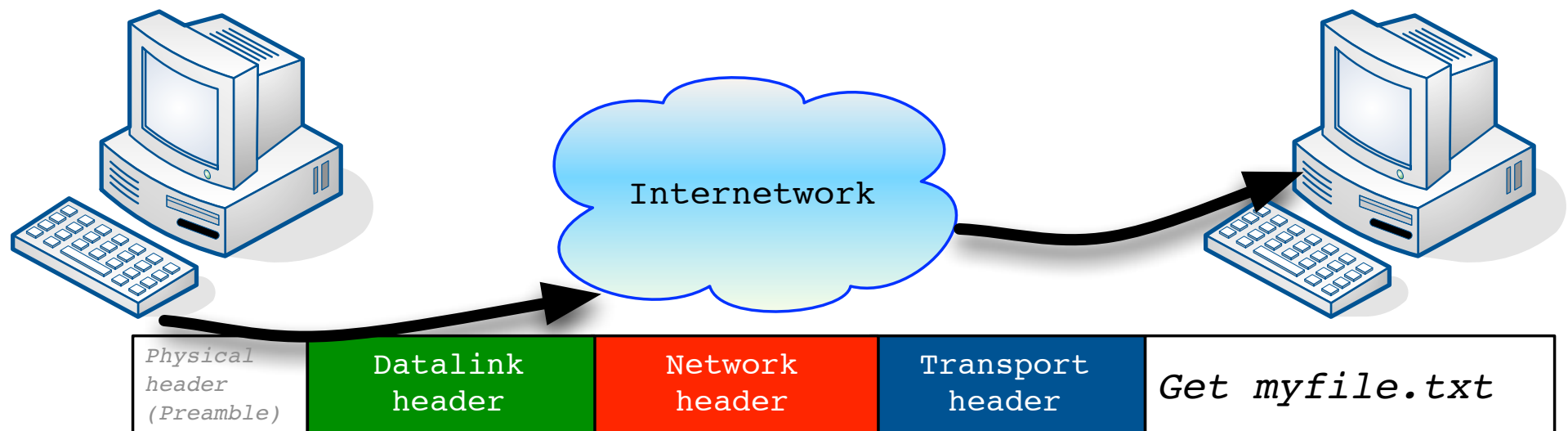


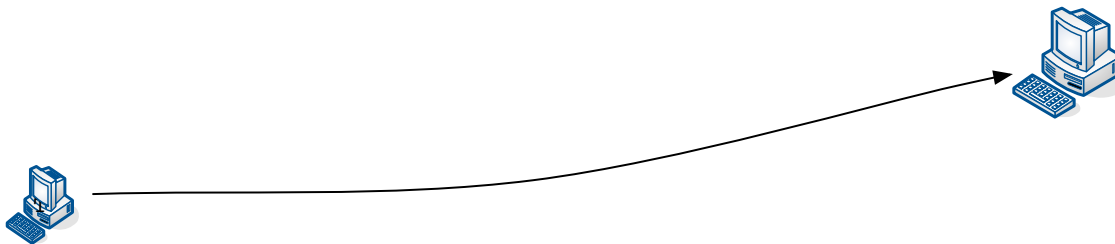
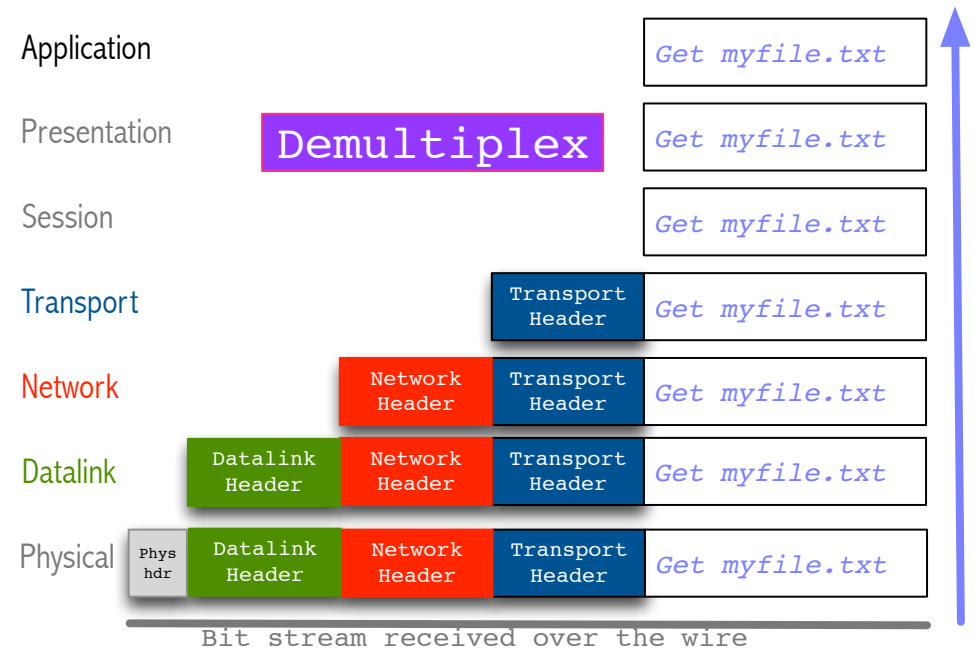
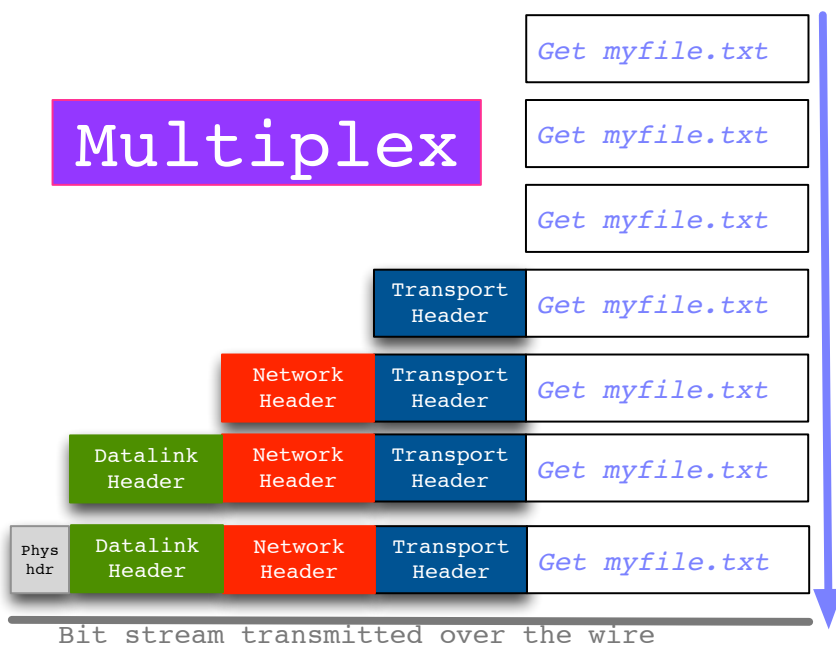
Illustration of encapsulation in OSI



Multiplexing

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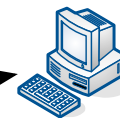
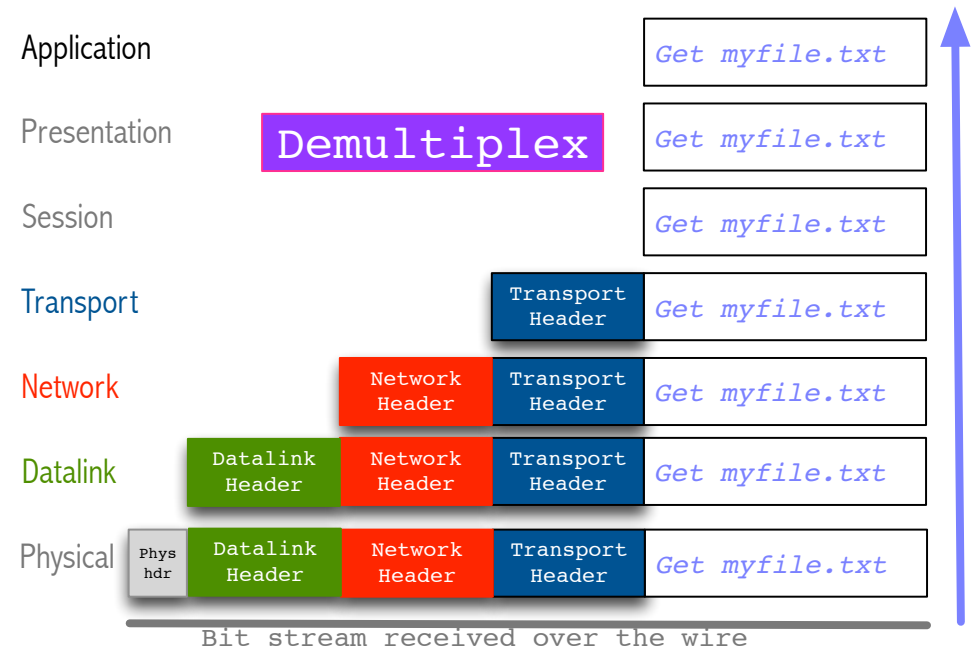
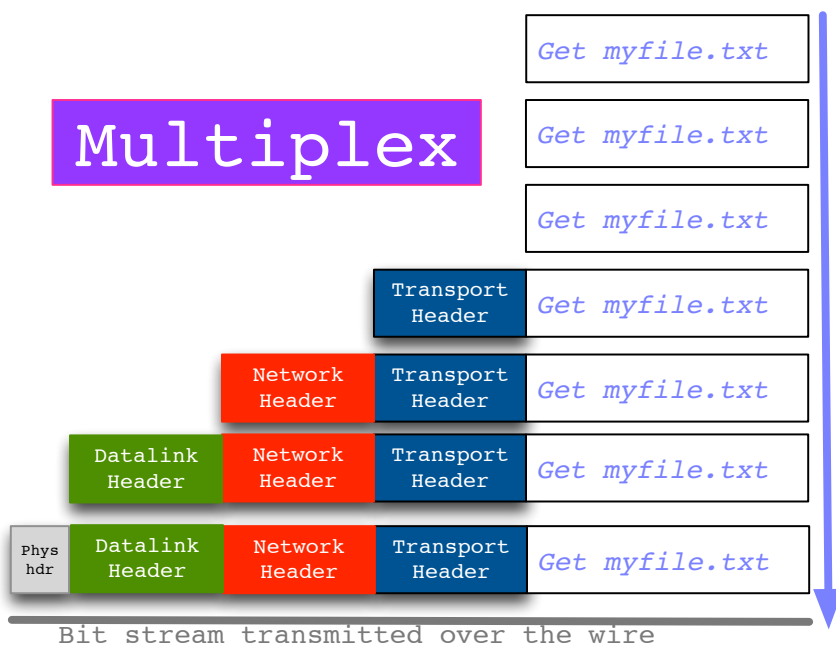
- Transmitter multiplexes several flows by having each layer add its header which contains addressing information



Demultiplexing

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- Receiver demultiplexes several flows by having each layer analyze its header which contains addressing information about the upper-layer protocol that is to receive the payload



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Connectivity

Computer Networks connect computers; the many more, the better, with a limit!

Theoretical connectivity

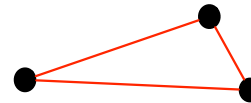
- Connectivity is the capacity of connection of a network
- If a network has N hosts, its connectivity is: $N \cdot (N-1) \cong N^2$
- Metcalf's law: The connectivity of a network grows fast as we add more nodes (N^2)



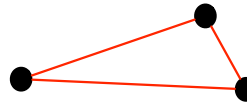
Connectivity = 1



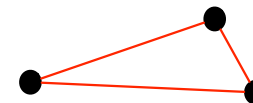
Connectivity = 2



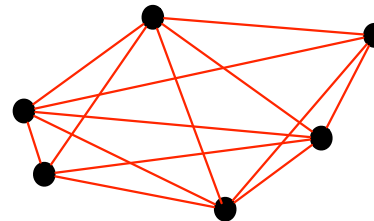
Connectivity = $3 \times 2 = 6$



+



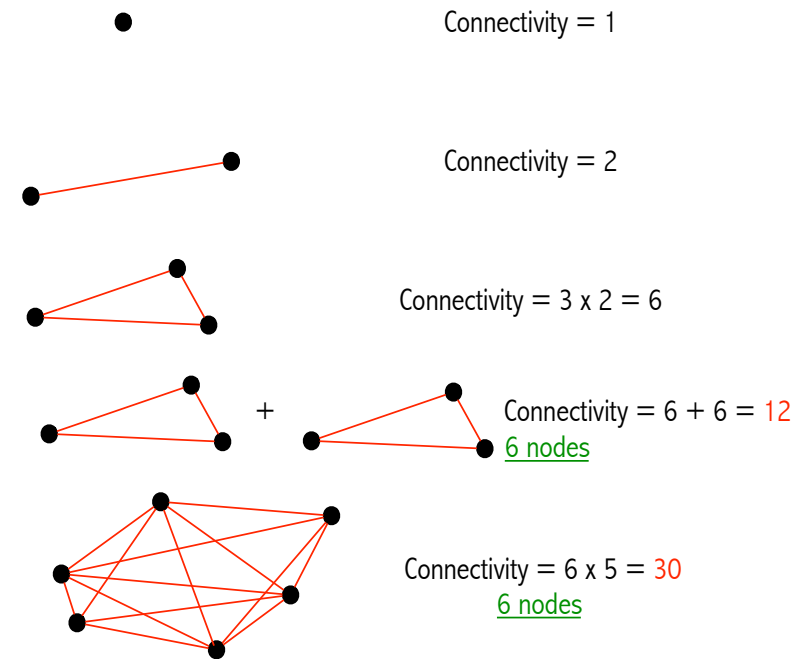
Connectivity = $6 + 6 = 12$
6 nodes



Connectivity = $6 \times 5 = 30$
6 nodes

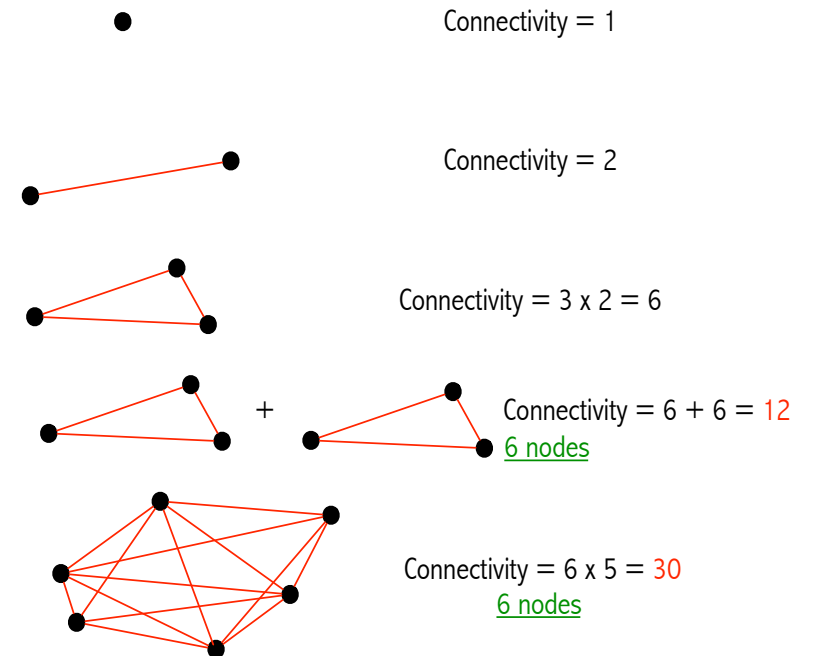
Increasing connectivity whilst preserving the capability for communication

- Metcalf's law
 - ▣ Increased connectivity means increased value
- Nodes communicate by sending/receiving messages
 - ▣ The bandwidth available at each link is limited
 - ▣ Links at highly demanded locations may become a bottleneck
- What's a figure of merit that will tell whether communication has been preserved after increasing the connectivity?
 - ▣ Is connectivity scalable?



Increasing connectivity whilst preserving the capability for communication

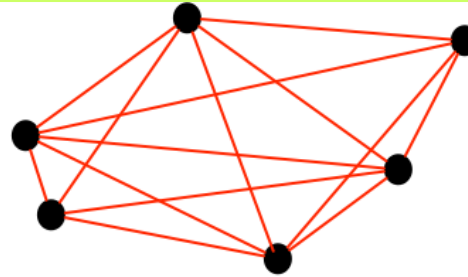
- The total number of packets per second that successfully make it to their destinations is known as:
 - ▣ Throughput, the figure of merit
 - ▣ Overall network productivity
 - ▣ Overall bps, or pps (packets per second), etc.



Scalable connectivity

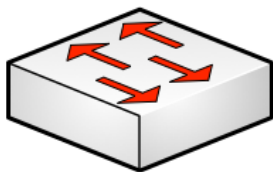
- Not all network technologies use the available connectivity with the same efficiency
- Ethernet can function efficiently up to certain network size: we say that Ethernet scales well up to that limit.
- Then, how come the Internet has 4000M hosts? How can the Internet scale to such a huge size so well?
 - ▣ Each network has a limited size
 - ▣ Interconnecting networks is the key:
 - With IP gateways
 - IP protocol

Switch won't scale to 9 hosts

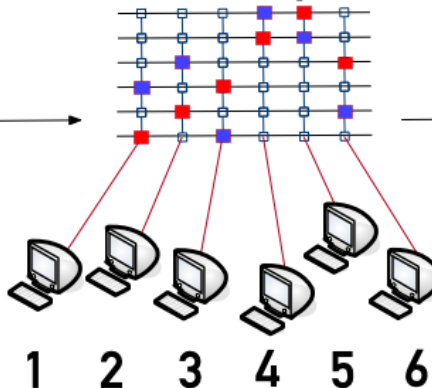


Full mesh comprised of 6 nodes
Connectivity = $6 \times (6 - 1)$
= 30 simplex links

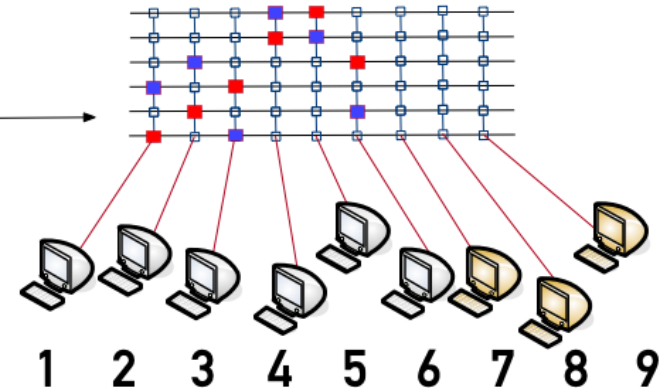
Implementation with a
9-port LAN Switch



Scalable connectivity
reaches 6 simplex links



Connectivity not available
for further hosts



Based on textbook *Conceptual Computer Networks* by:

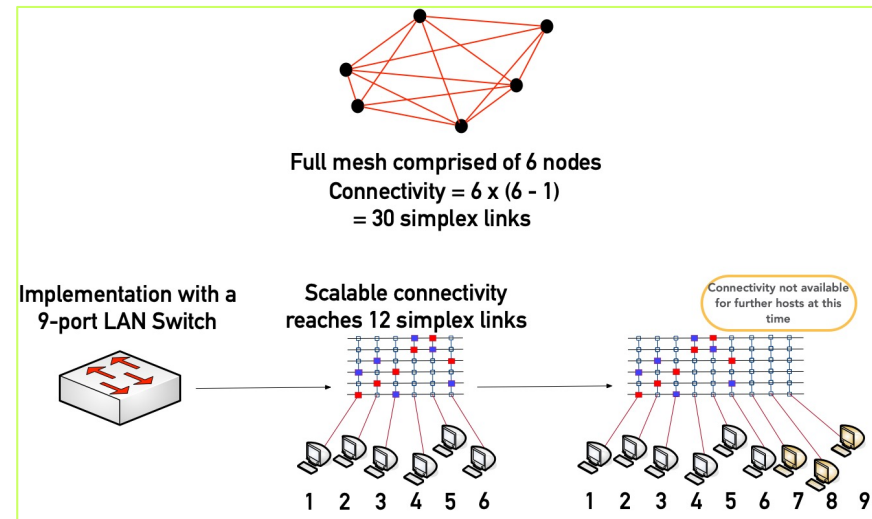
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Switch won't scale to 9 hosts

6 full-duplex communication flows are possible

- The overall productivity of this switch will be bounded by
 - ▣ The available number of connection points
 - ▣ The available number of horizontal lines
- Throughput, the total pps or bps will be bounded by the limited switch resources
 - ▣ PPS = Average number of Packets Per Second that the switch can successfully deliver
 - ▣ Bps = Average number of bits per second that the switch can successfully deliver

Physically connecting more than 6 hosts will not achieve a Throughput improvement



Network Throughput of various wireless technologies

EthAir scales poorly

