

Based on “Computer Networks: A Systems Approach”, 5e
By Larry L. Peterson and Bruce S. Davie, Morgan-Kaufmann / Elsevier

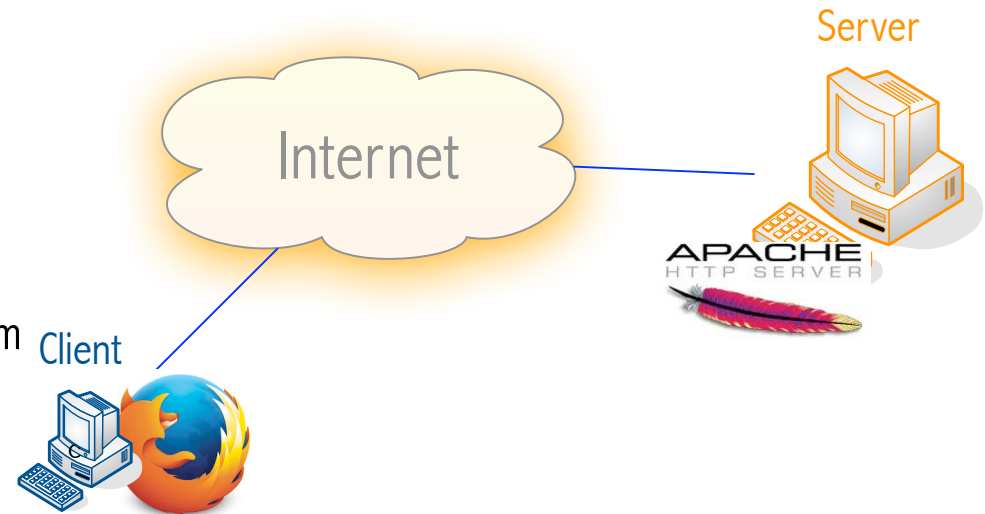
Extended and adapted by José María Foces Morán (2015)

CHAPTER 1: FOUNDATION

An introduction to Computer Networks

Essential example: www

- **www** = World Wide Web:
 - A distributed and Client/Server application
 - Server program (e.g., Apache)
 - Client program (e.g., Firefox)
- URL
 - Uniform Resource Locator
 - <http://paloalto.unileon.es/cn/index.html>
- Web pages are downloaded by the client from the **server**
 - Client and server speak the **http** protocol
 - http = Hyper Text Transfer Protocol
- HTTP, in turn uses the **TCP** protocol
 - TCP = Transmission Control Protocol
 - TCP provides reliability



Main metrics used in Networking



□ Bandwidth

- Directly related to the acceptable speed of bit transmission over some medium
- Number of bits transmitted in one second:
 - Bps (Bits Per Second = Bits/Sec)
- Since bandwidth is a rate, the multipliers take on the following values:
 - K (Kilo = 10^3)
 - M (Mega = 10^6)
 - G (Giga = 10^9)
 - T (Tera = 10^{12})

□ Delay

- Seconds
- How much time it takes to transport one bit from a source to a destination directly connected
- Propagation delay

□ Jitter

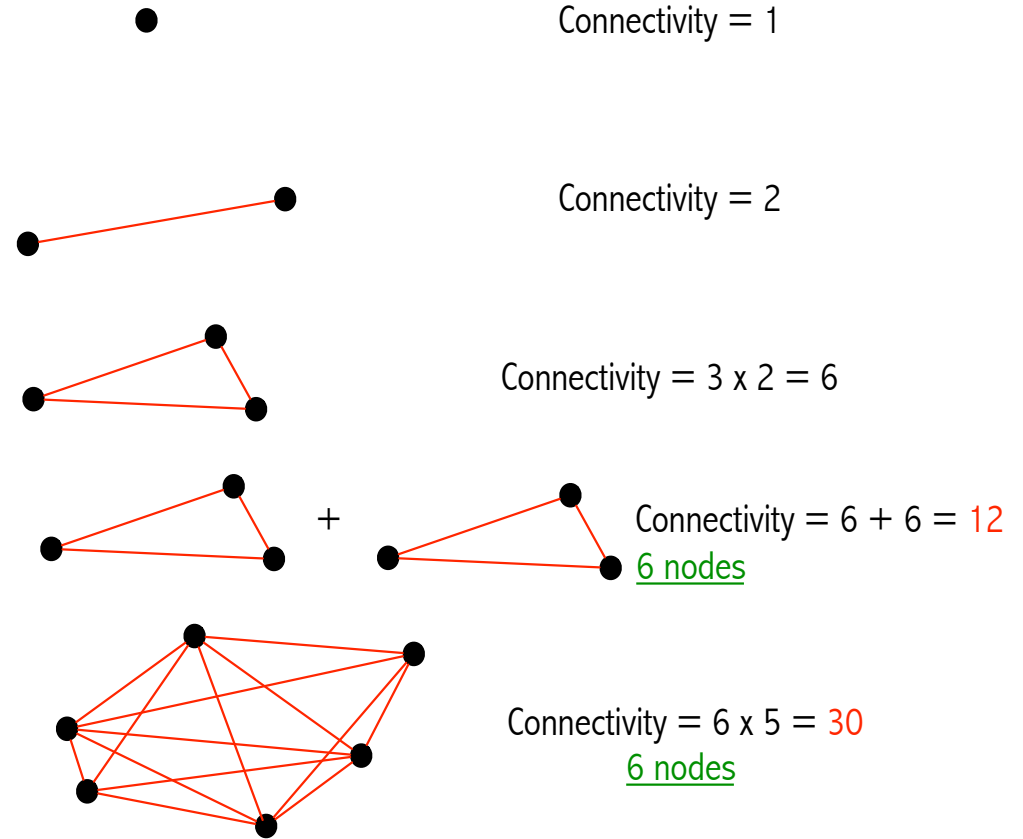
- The variance of the delay

Connectivity

Computer Networks connect computers

Theoretical connectivity

- Connectivity is the capacity of connection of a network
- If a network has N hosts, its connectivity is: $N \cdot (N-1) \cong N^2$
- Metcalfe's law: The connectivity of a network grows fast as we add more nodes (N^2)



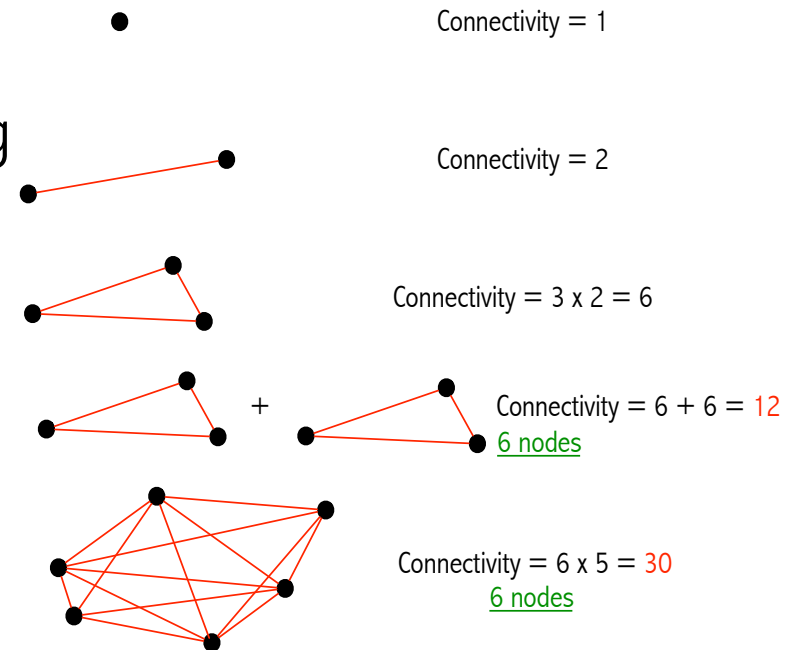
Theoretical connectivity is bounded by network *technology*

□ Metcalf's law: Represents the potential connectivity of a network

□ Nodes communicate by sending/receiving messages through the links which bandwidth is limited

□ Each node must 'have some knowledge' about the topology of the network

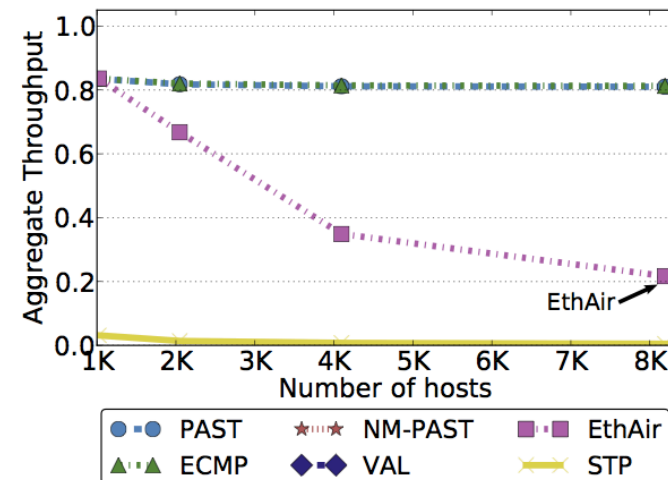
□ These factors limit the use of the potential connectivity



Scalable connectivity

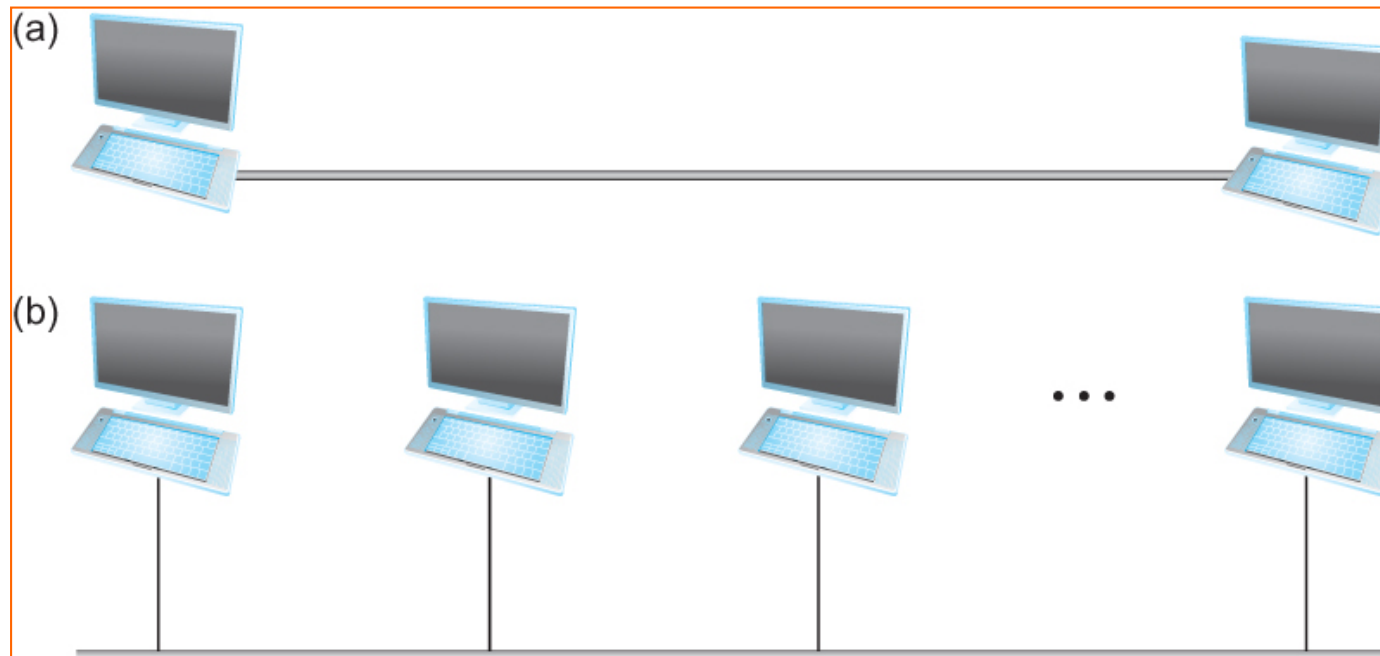
- Not all network technologies use the available connectivity with the same efficiency
- Ethernet can function efficiently up to certain network size: we say that Ethernet scales well up to that limit.
- Then, how come the Internet has 2000M hosts? How can the Internet scale to such a huge size so well?
 - IP protocol

EthAir scales poorly



Attaining scalable connectivity

- Something is scalable if it can grow to huge sizes preserving its functionality
- We want that the connectivity offered by a network be scalable, i.e., that it can grow as needed



Basic Network terms (i)

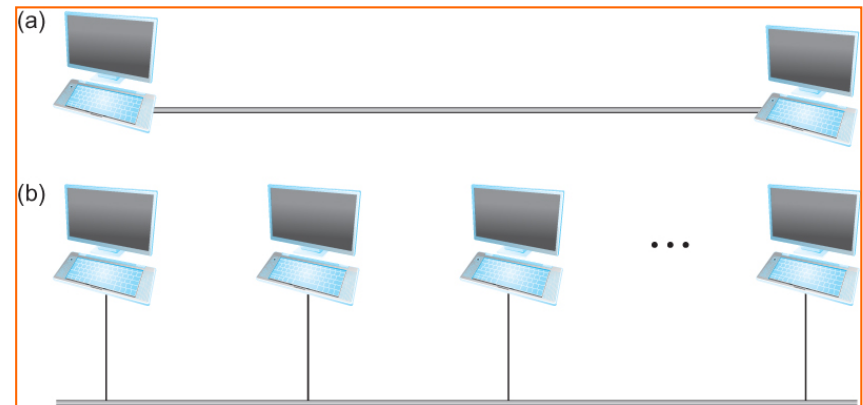
Scale (Size)

Host 

Link 

Communication links:

- ▣ (a) Point-to-point
- ▣ (b) Multiple access



Which one scales better?

Basic terms (ii)



□ Switch



□ Network



□ Switched Network

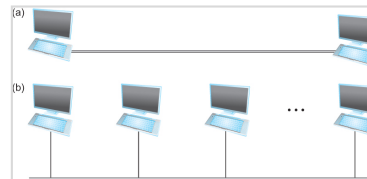
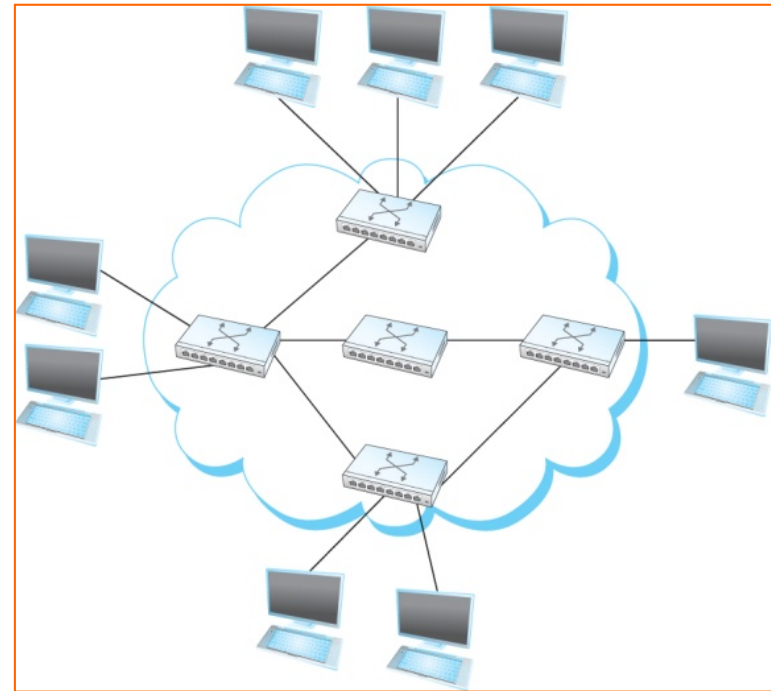
▣ Circuit Switched (Old telephone network)

▣ Packet Switched

□ Packet and message


□ Store-and-forward

Switched network (A single network !!!)

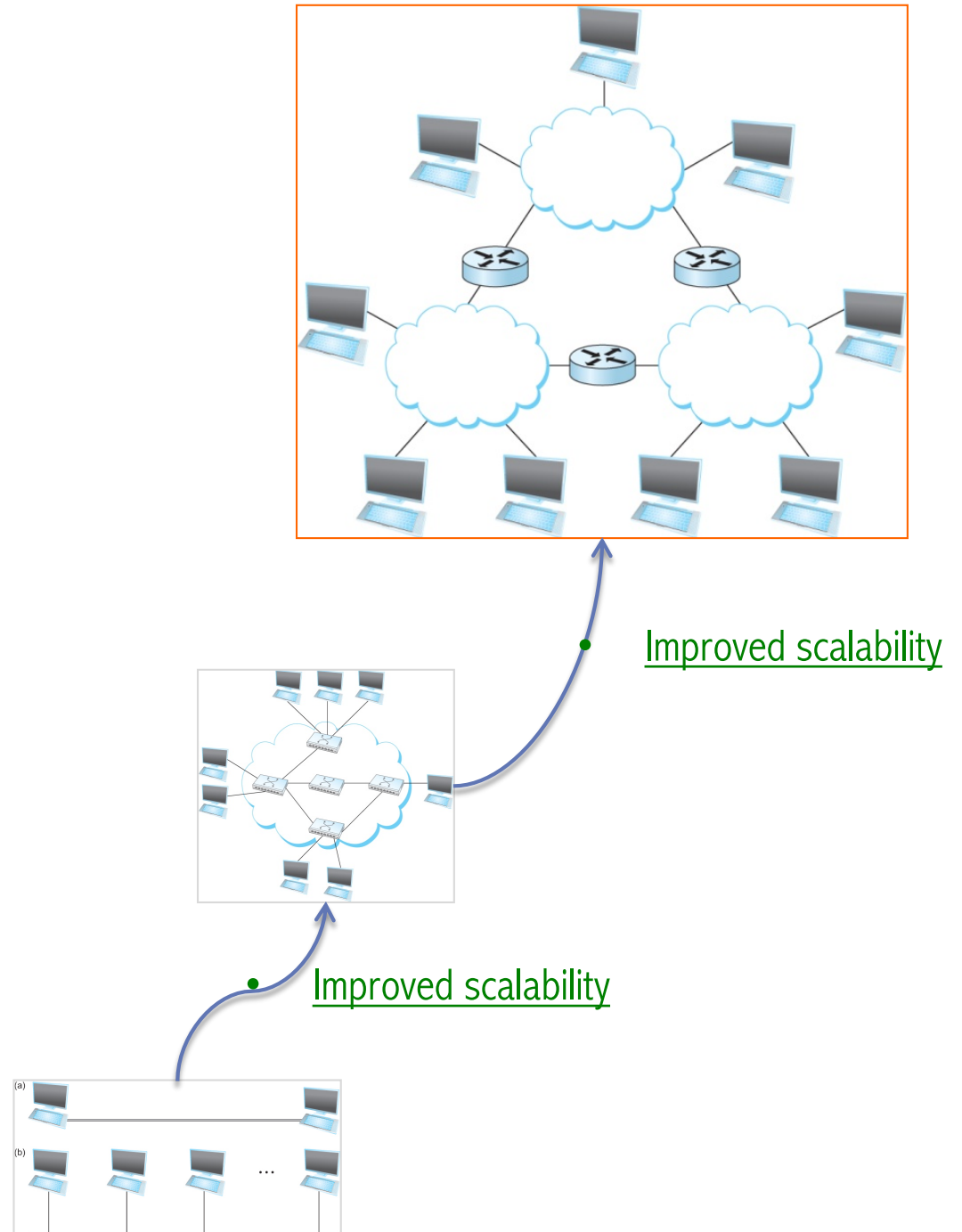


• Improved scalability

Basic terms (iii)

- Router/gateway 
- Host-to-host connectivity
- Address
- Routing
- Unicast/broadcast/multicast

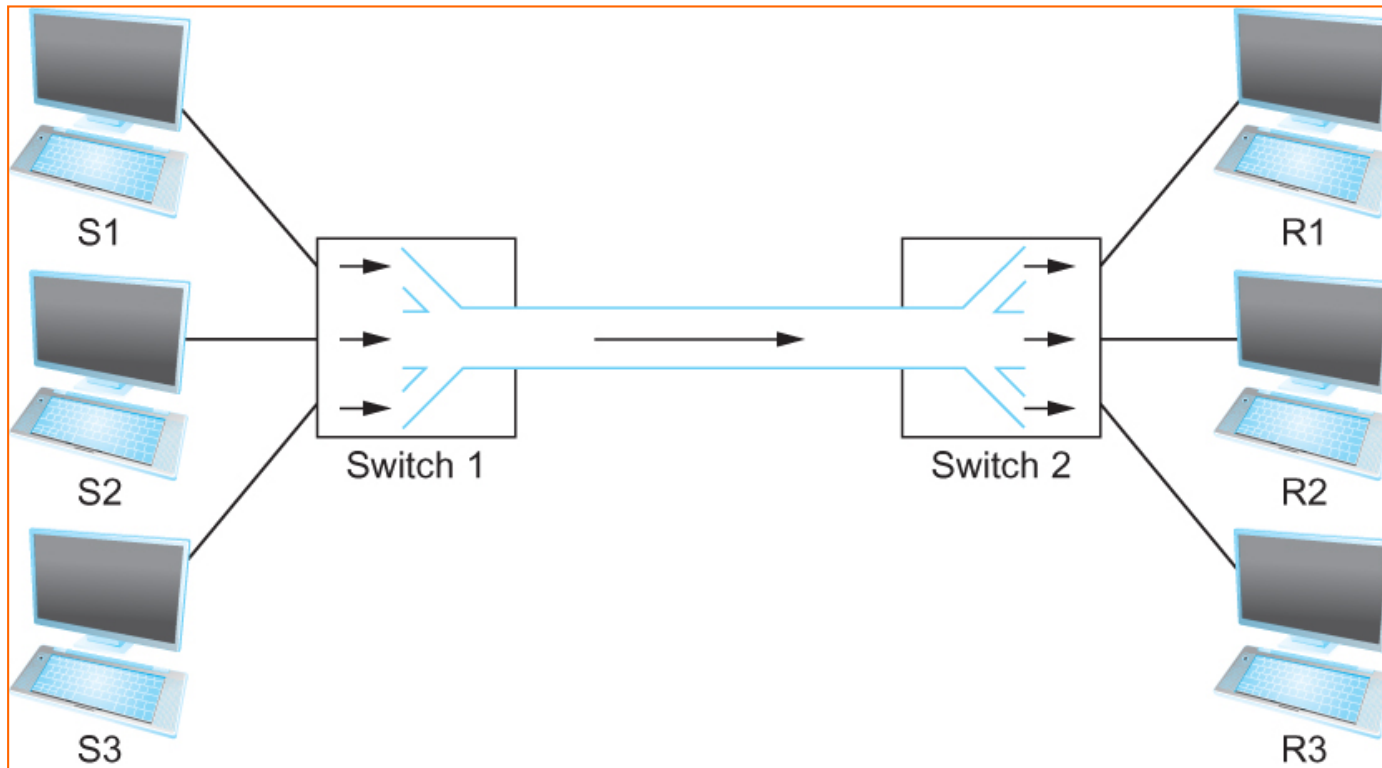
Interconnection of networks (Several networks interconnected)



Sharing connectivity

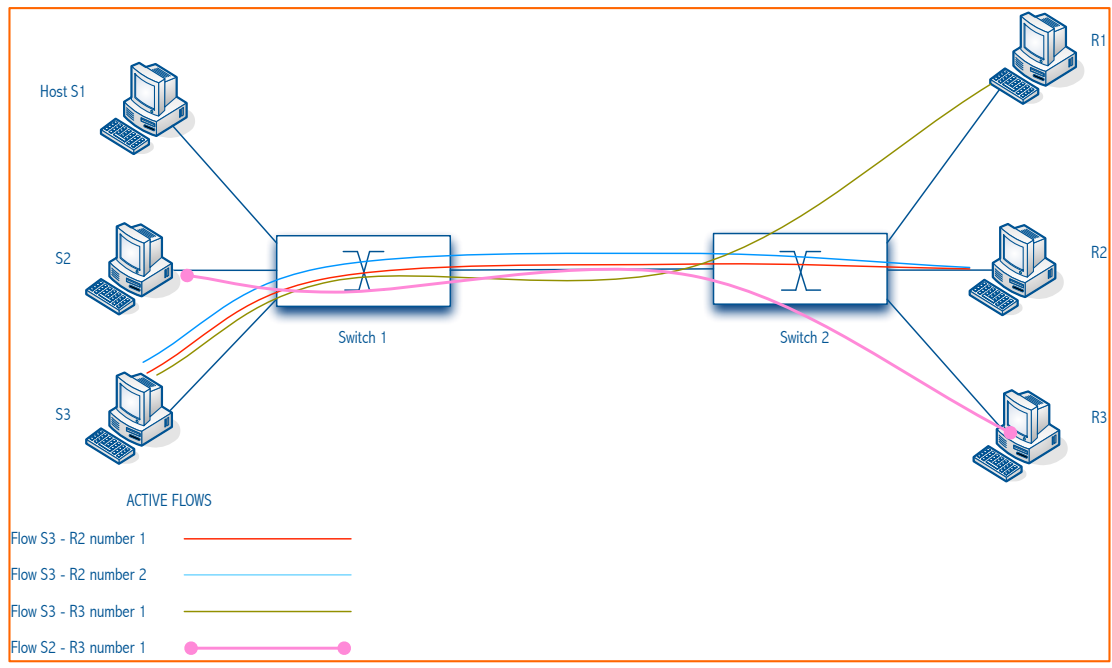
How to share a link, a network among multiple users

Cost-Effective Resource Sharing



Multiplexing multiple logical flows over a single physical link

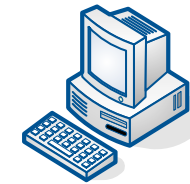
Review of terminology



Multiplexing multiple logical flows over a single physical link

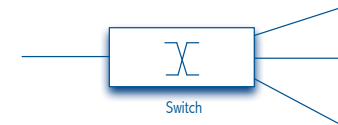
□ **Hosts:** the computers that run the application programs

- ▣ Clients
- ▣ Servers



□ **Network nodes**

- ▣ DataComm equipment (DCE) for building a network

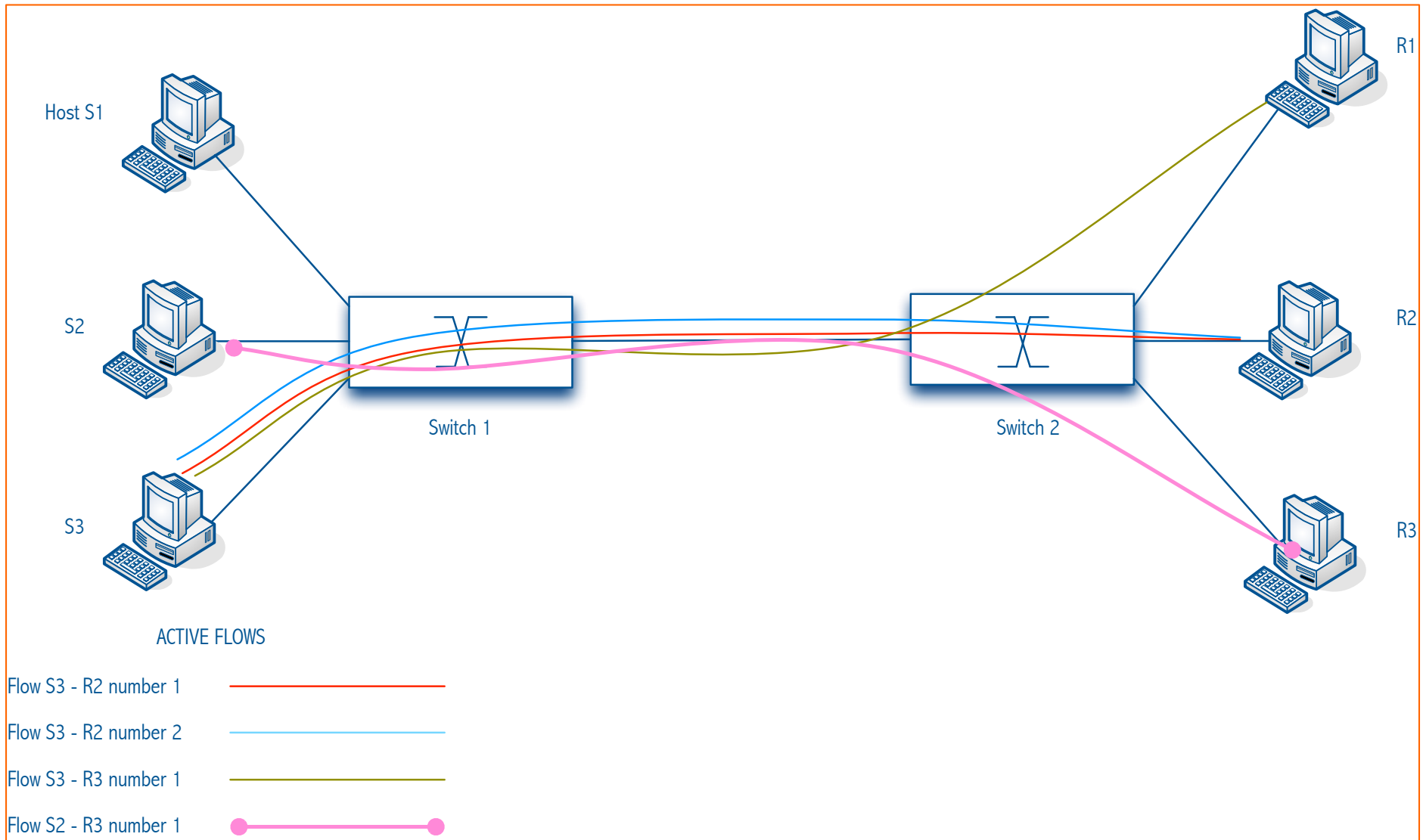


□ **Links:**

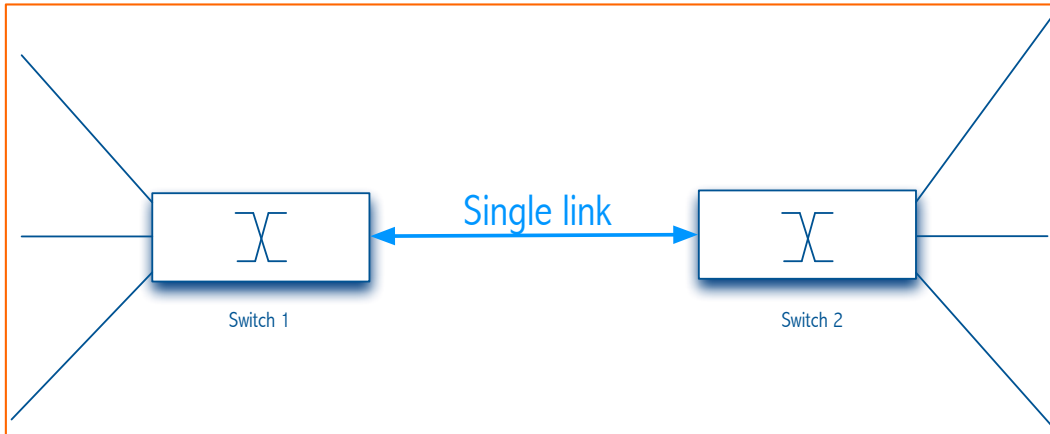
- ▣ The physical transmission media and other logical characteristics
 - ▣ Simplex
 - ▣ Half-duplex
 - ▣ Full-duplex
- ▣ *WiFi, Twisted pair cables, optical fibers, etc*

□ **FLOW** comes next

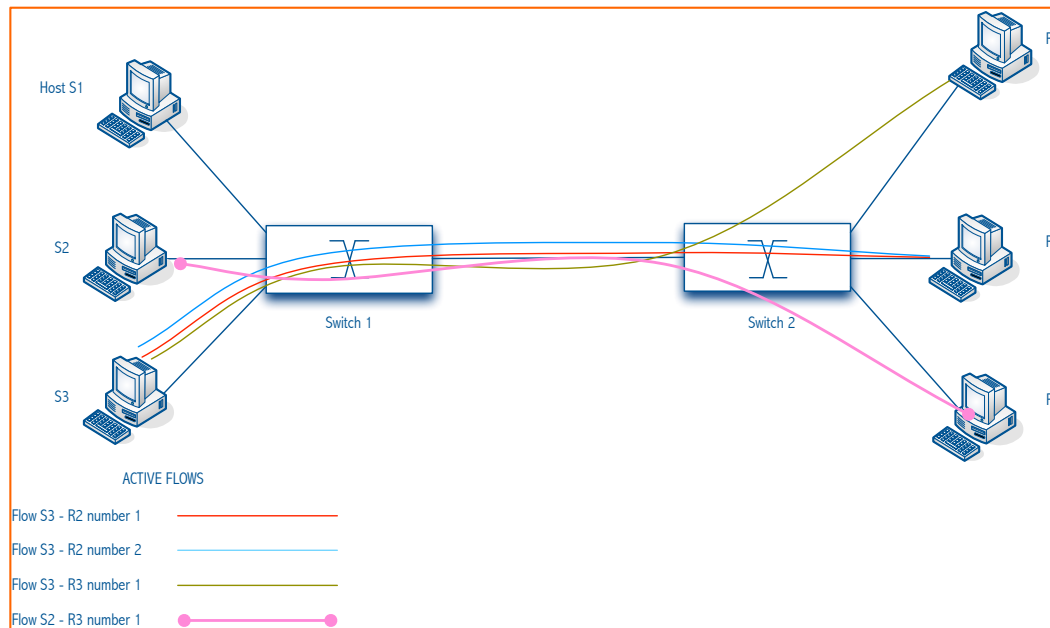
What is a *logical* flow?



Multiplexing = sharing



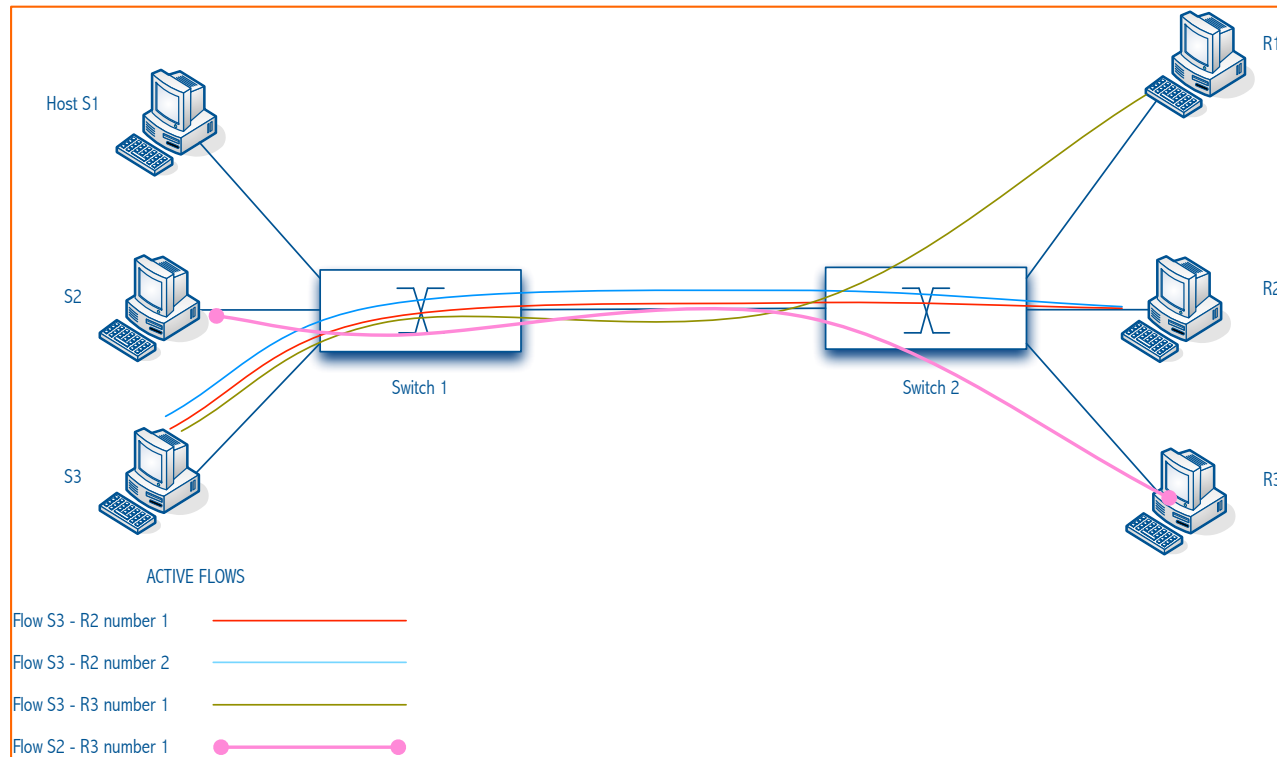
- Switch 1 and Switch 2 are linked by a single link
- How can that link be shared among the hosts?
- How can that sharing be done?
 - ▣ By **multiplexing** the flows of packets over the link that connects the switches



Multiplexing techniques

1. Synchronous Time-division Multiplexing (STDM)

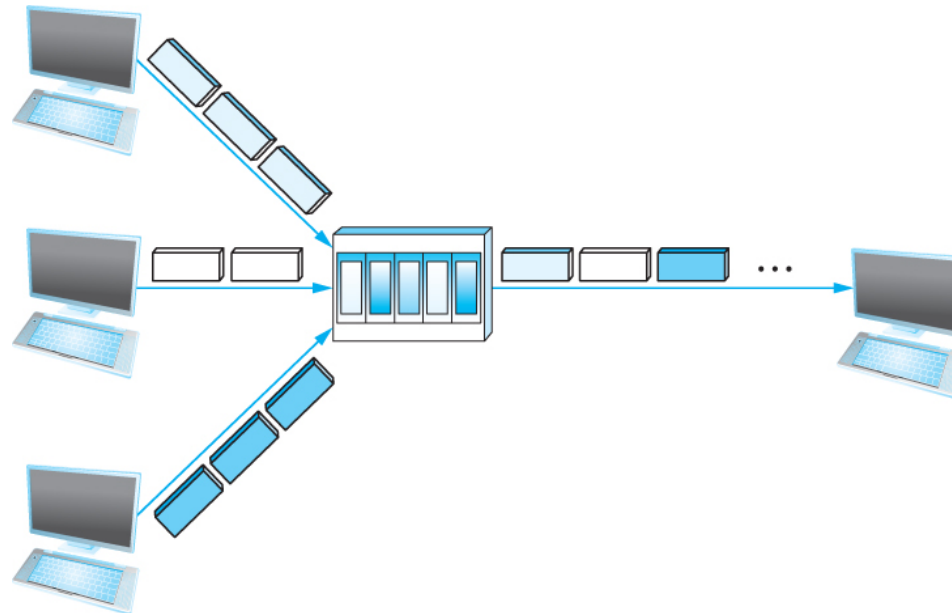
- ▣ Allocate a predetermined time slot to each flow
- ▣ Drawbacks
 - If a flow is not transmitting (idle), its slot remains allocated and is wasted



Multiplexing techniques

2. FDM: Frequency Division Multiplexing

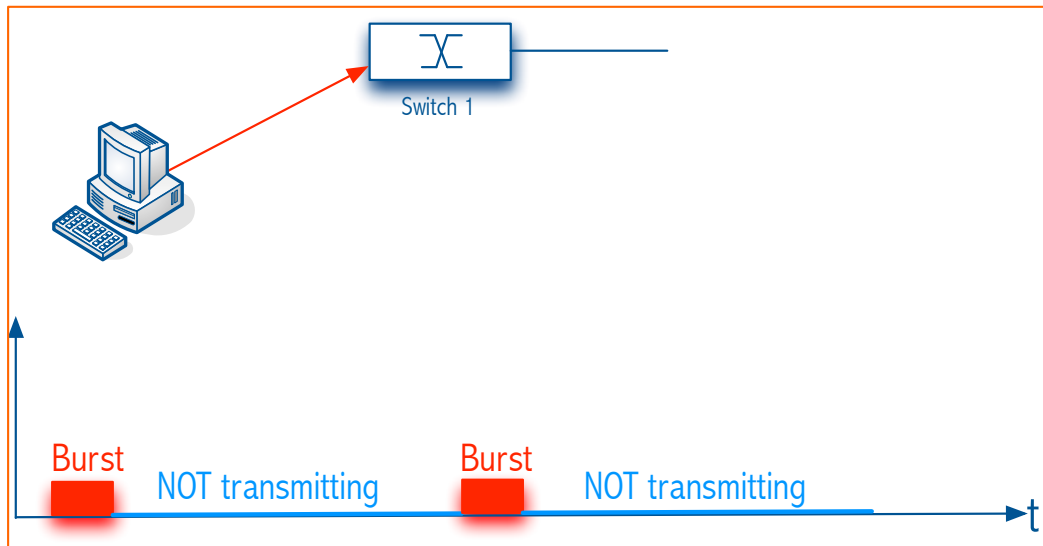
- ▣ *Similar* to FM radio
- ▣ Each flow is assigned a portion (a band) of the link spectrum, each portion is the same size (bandwidth)



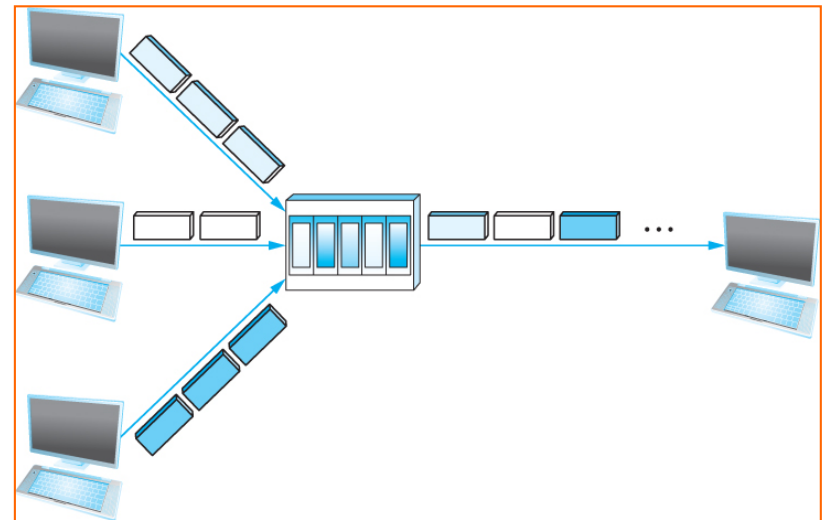
Multiplexing techniques

3. Statistical Multiplexing

- Data is transmitted based on the level of demand (offered load) of each flow
 - Sources are supposed to be bursty
 - The bandwidth of the shared link is a fraction of the peak aggregated bandwidth of the other links



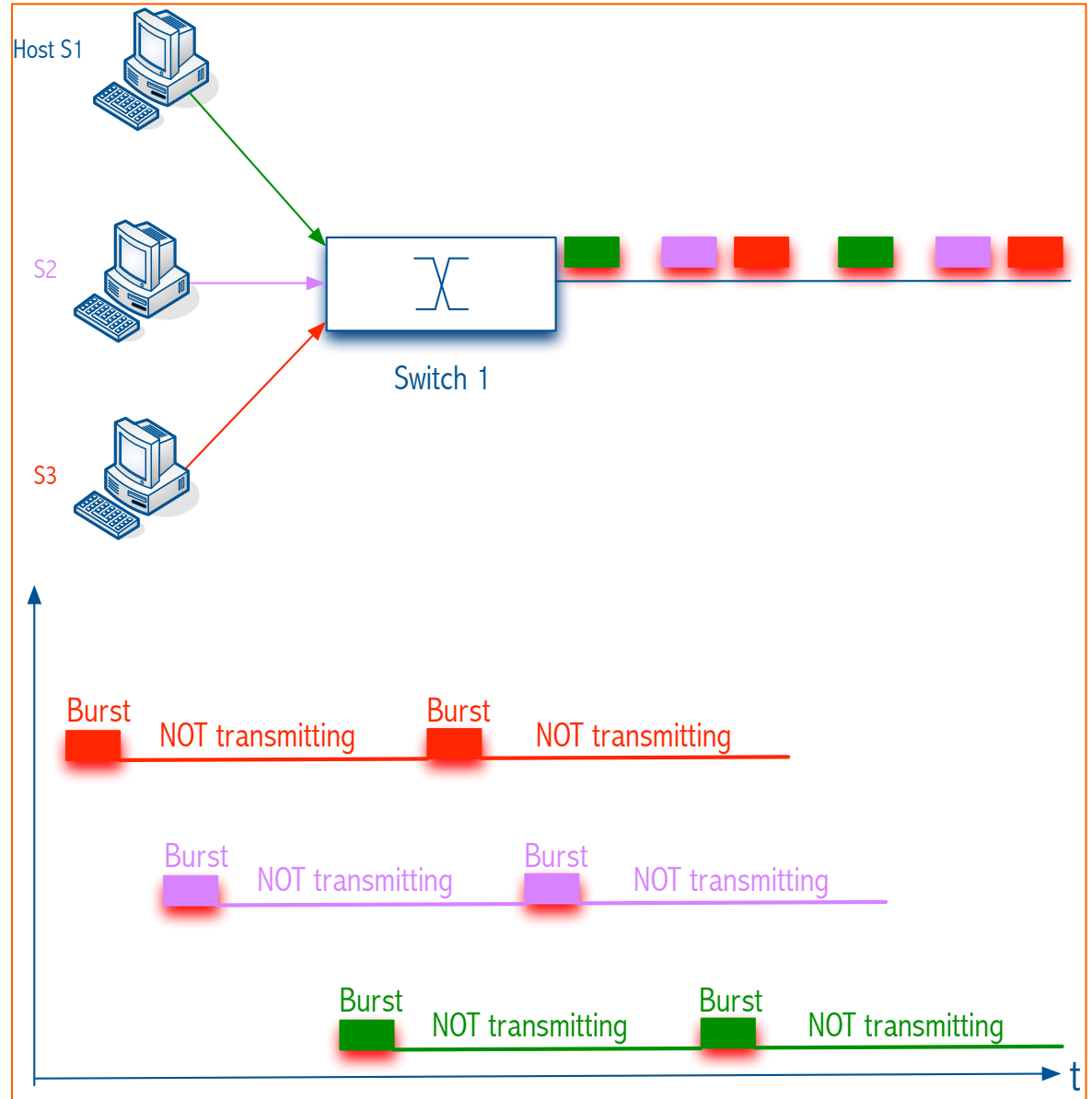
A bursty source



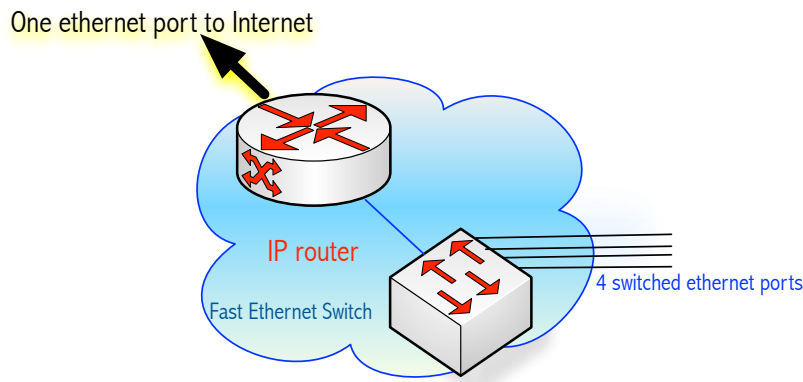
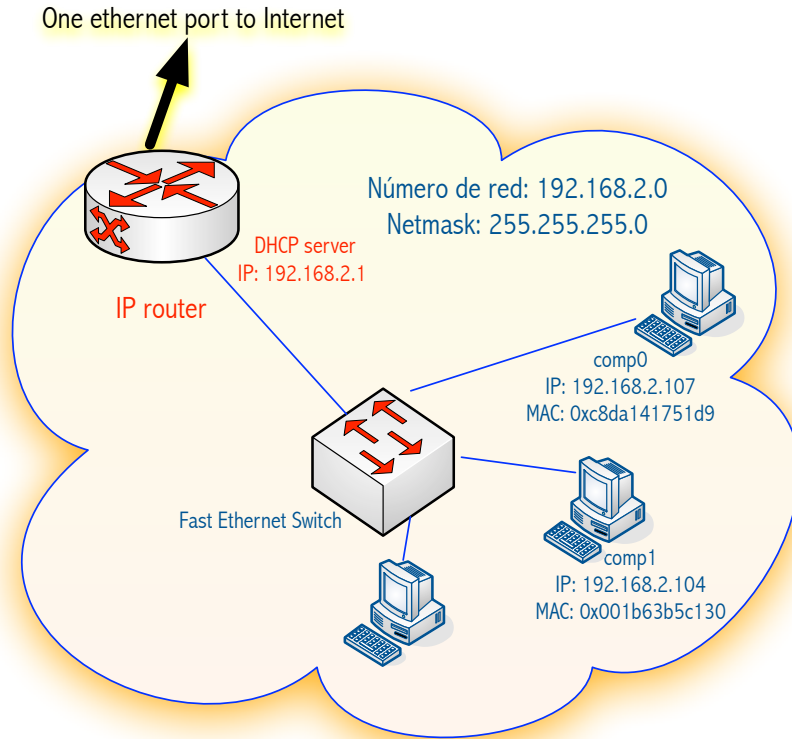
Multiplexing techniques

3. Statistical Multiplexing

- ▣ Data is transmitted based on the level of demand (offered load) of each flow
 - ▣ Sources are supposed to be bursty
 - ▣ The bandwidth of the shared link is a fraction of the peak aggregated bandwidth of the other links
- ▣ Packets vs. Messages
- ▣ Some switches may apply QoS (Quality of Service):
 - ▣ FIFO, Round-Robin, Priorities
- ▣ Congestion



Traditional types of networks



Depending on their geographical area

- LAN : Local Area Network
 - Ethernet
- MAN: Metropolitan
 - WiMax
- WAN: Wide
 - X.25
- SAN (System Area Networks)
 - Interconnect hard disks, network storage
 - High speed

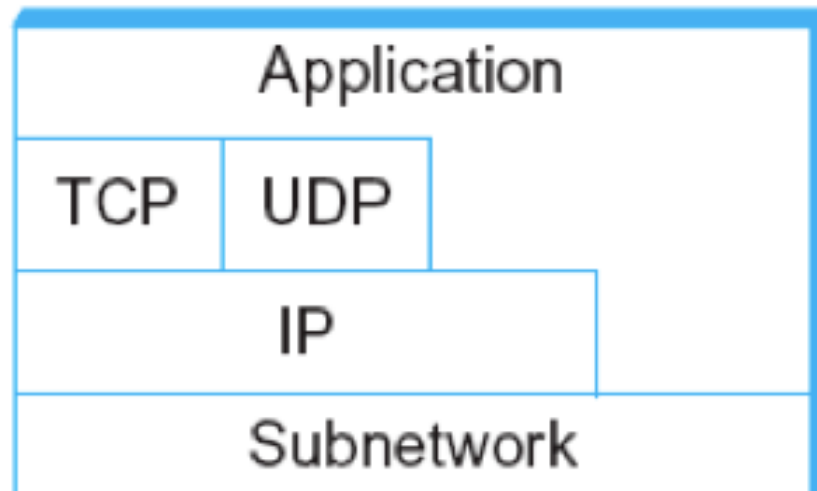
Network Architecture

How to hide the complexity of networks in order to promote the development of new network uses

Architecture



TCP/IP architecture!

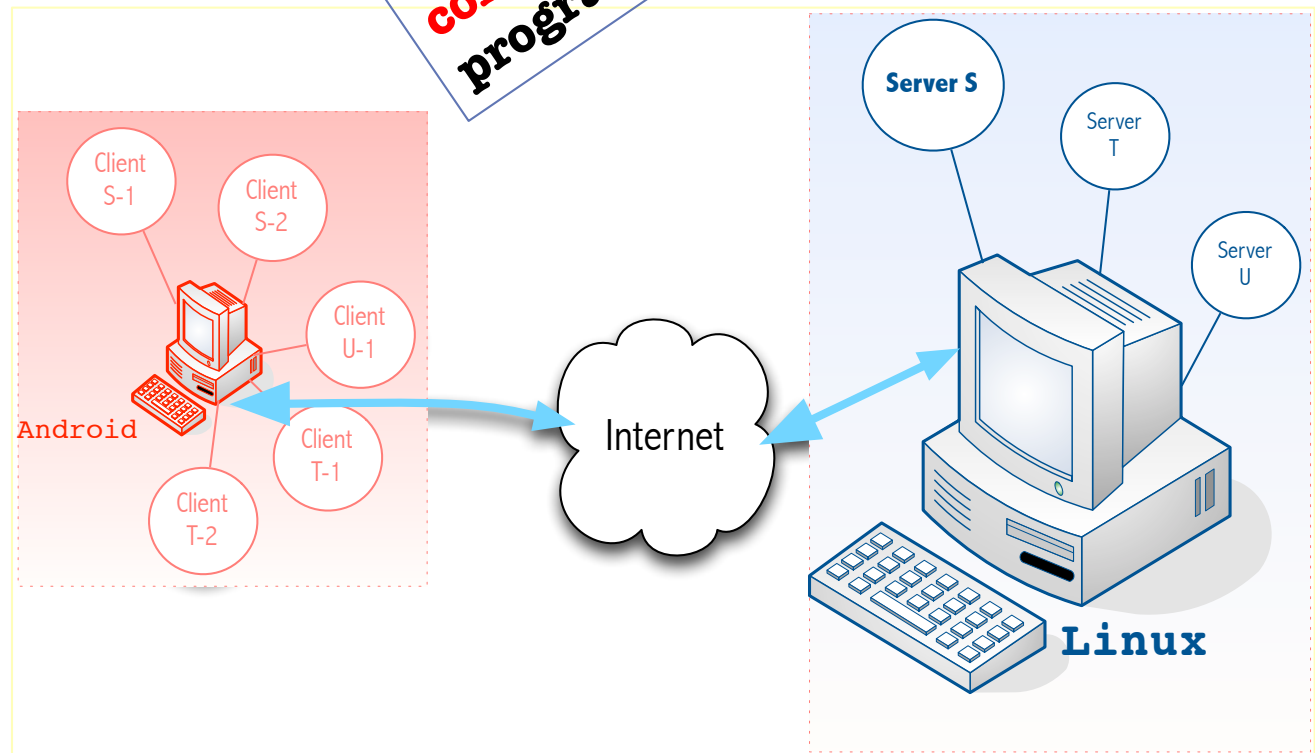


Network Architecture

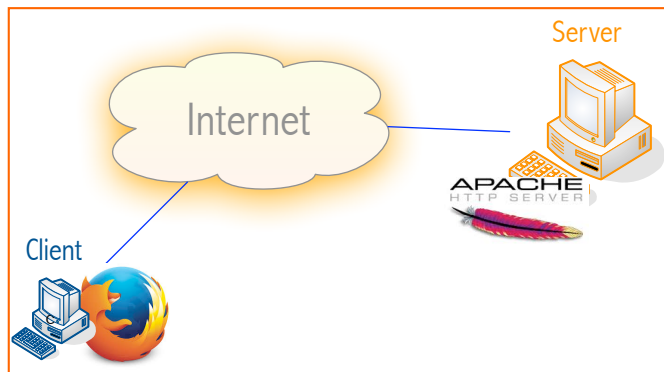
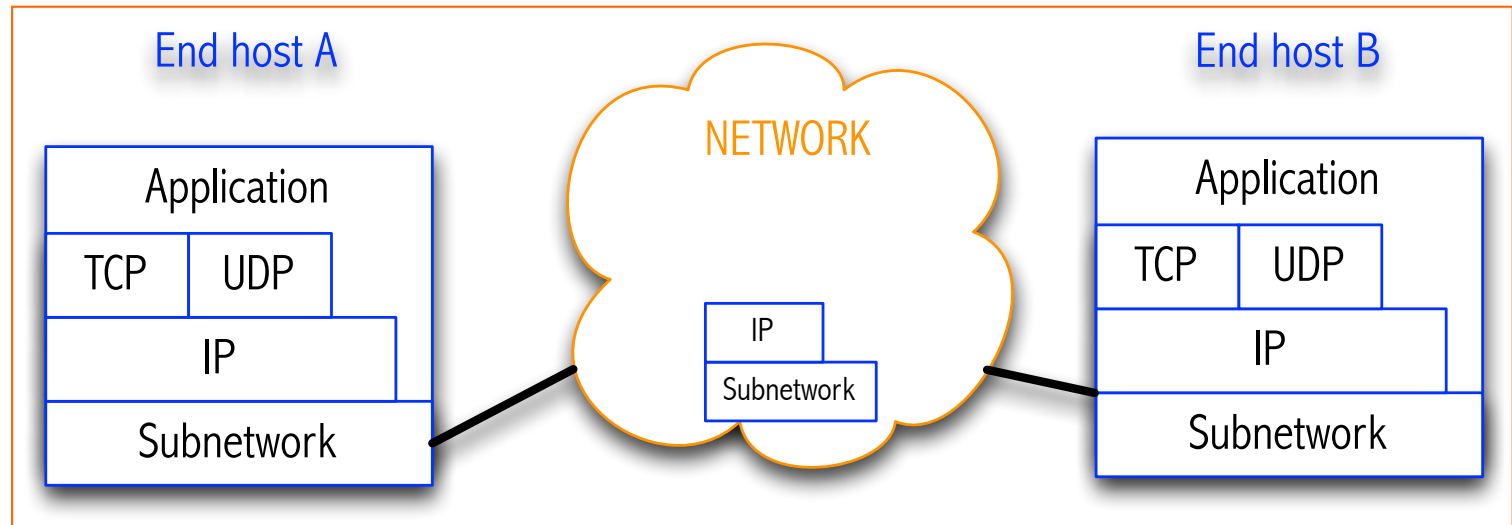
How to **hide** network **complexity** from application **programmers?**

CHANNEL

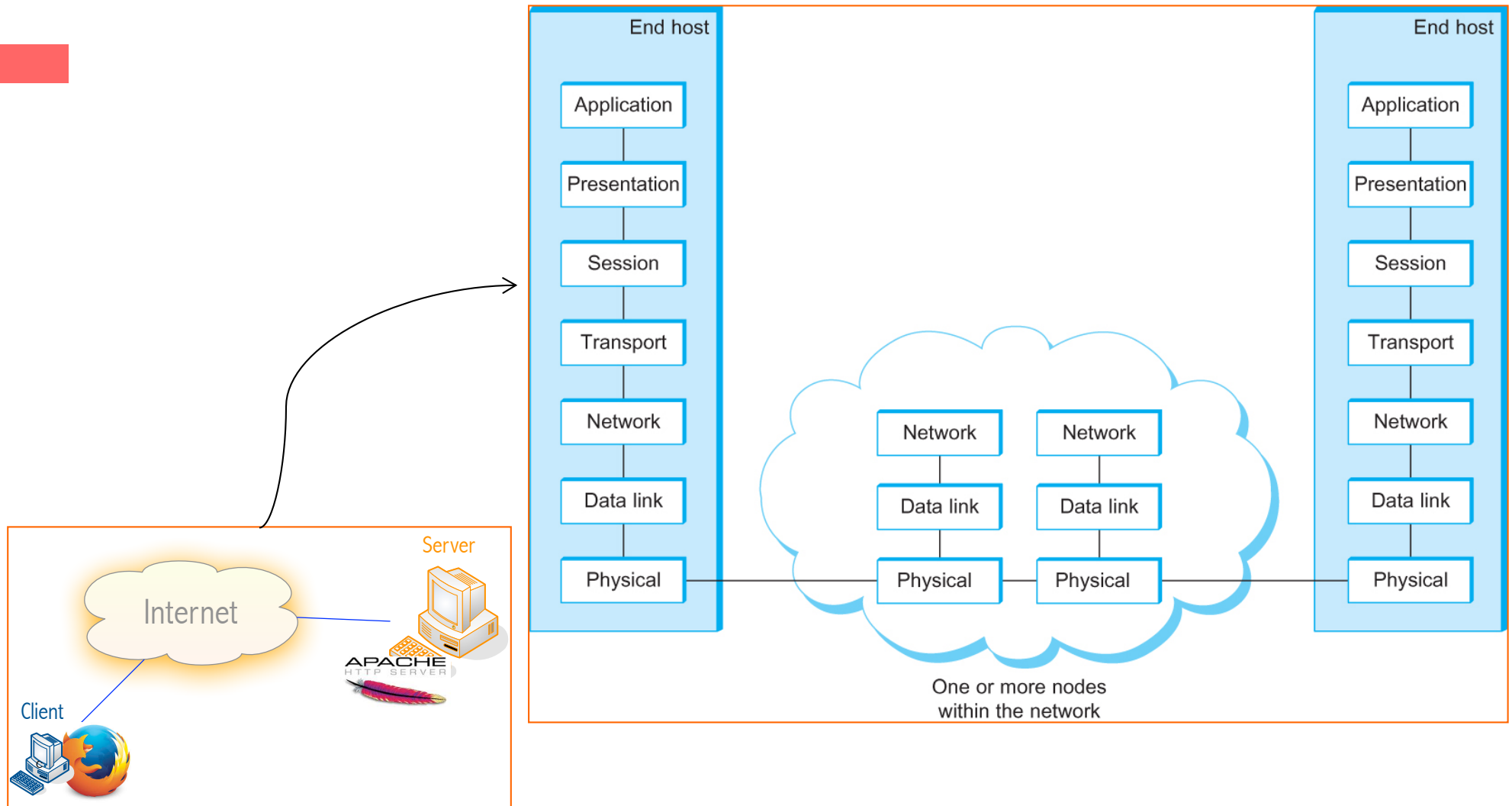
- Connects applications
- Logical channels
- Hosts must be identified:
 - IP address
- Applications/processes must be identified:
 - Port numbers



TCP/IP architecture



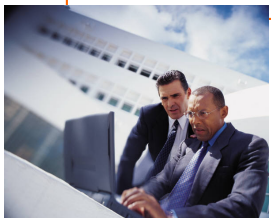
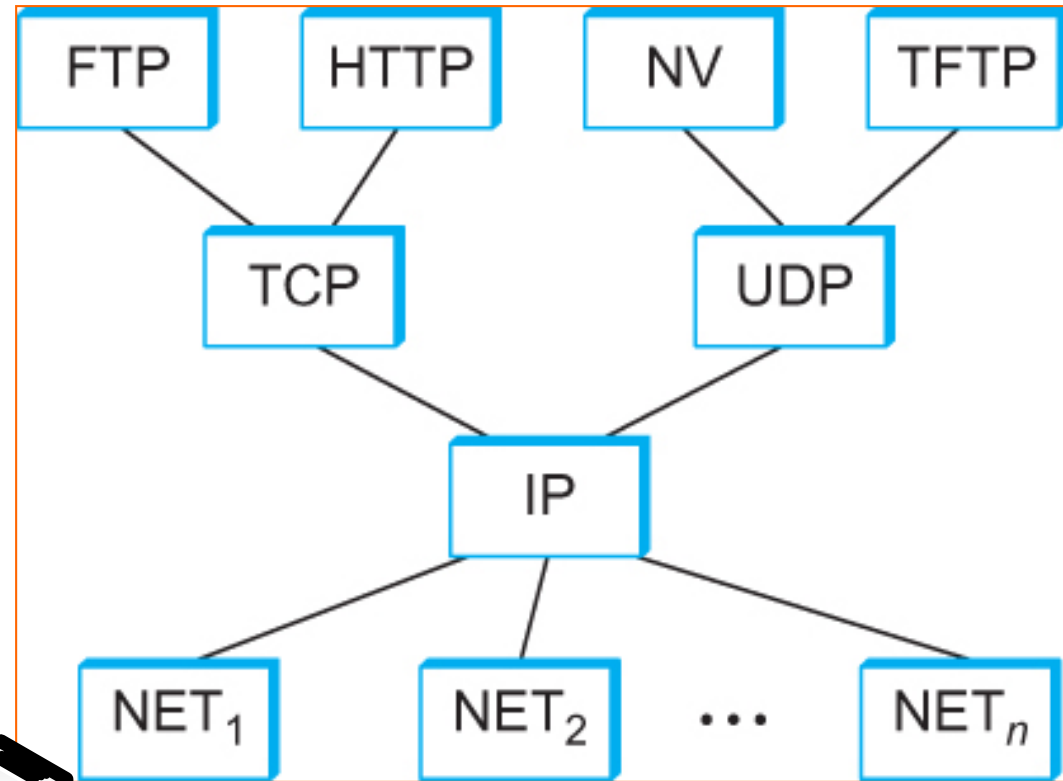
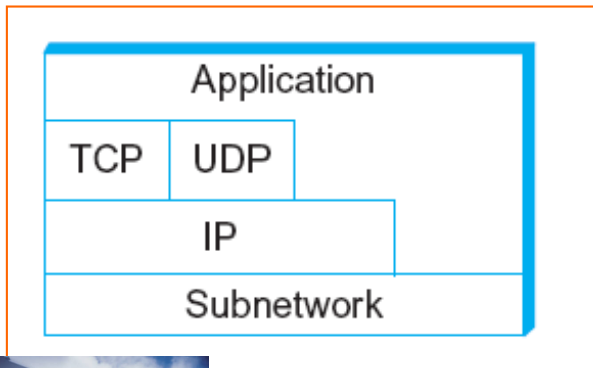
OSI Architecture



Example

A protocol graph of a host that implements the Internet architecture

Internet architecture

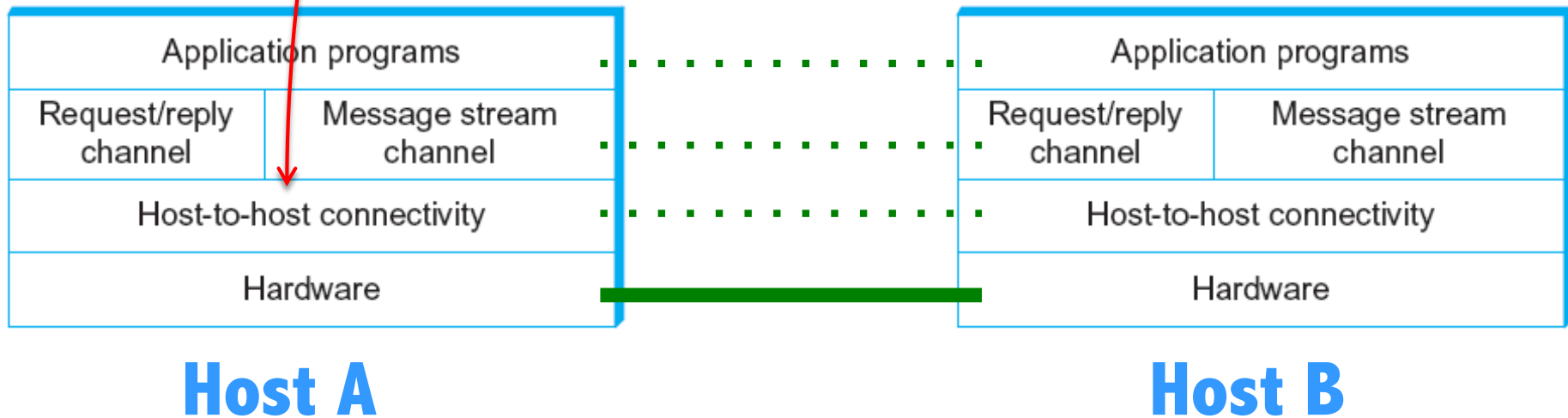


Protocols and Services

The abstractions exposed by the architecture

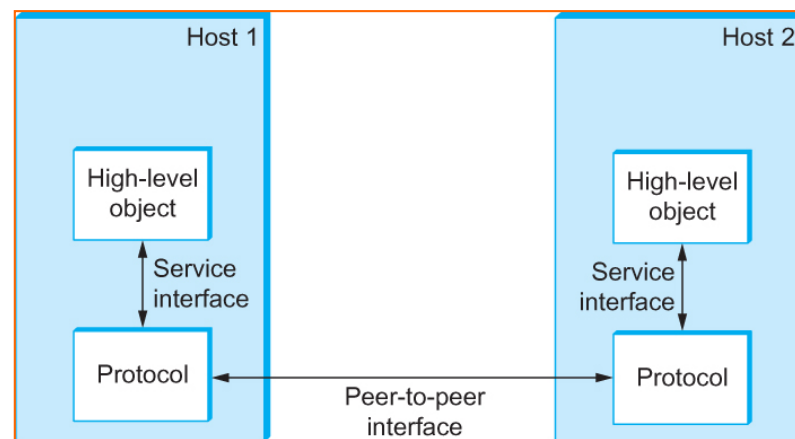
Protocols: building blocks of a network architecture

- A protocol defines two interfaces:
 - 1. Between the **layers in the same system**
 - Operations available at a specific layer
 - 2. Between **layers of peer systems (A – B)**
 - Messages exchanged between A and B, their syntax and semantics



Protocols

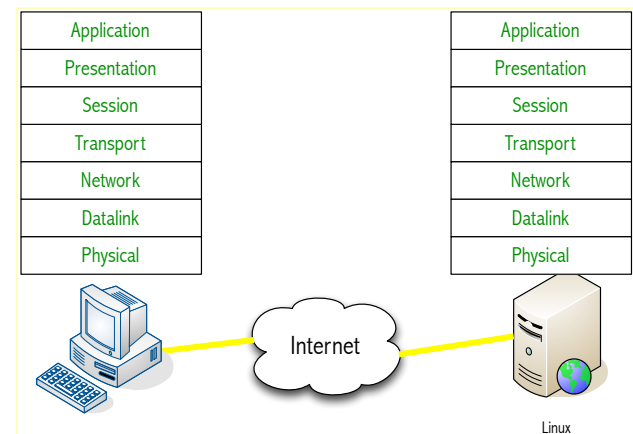
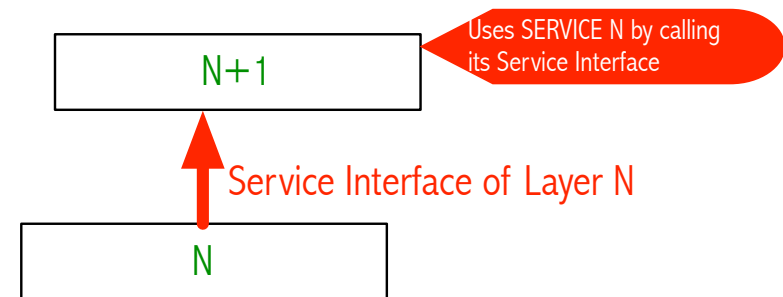
- Each **protocol object** has two different *interfaces*
 - ▣ **service** interface: operations on this protocol within this host
 - ▣ **peer-to-peer** interface: messages exchanged with peer (The peer host)
- Term “protocol” has two uses in CN:
 - ▣ Specification of **peer-to-peer** interface
 - ▣ Module (of software) that implements this interface



Layering in CN

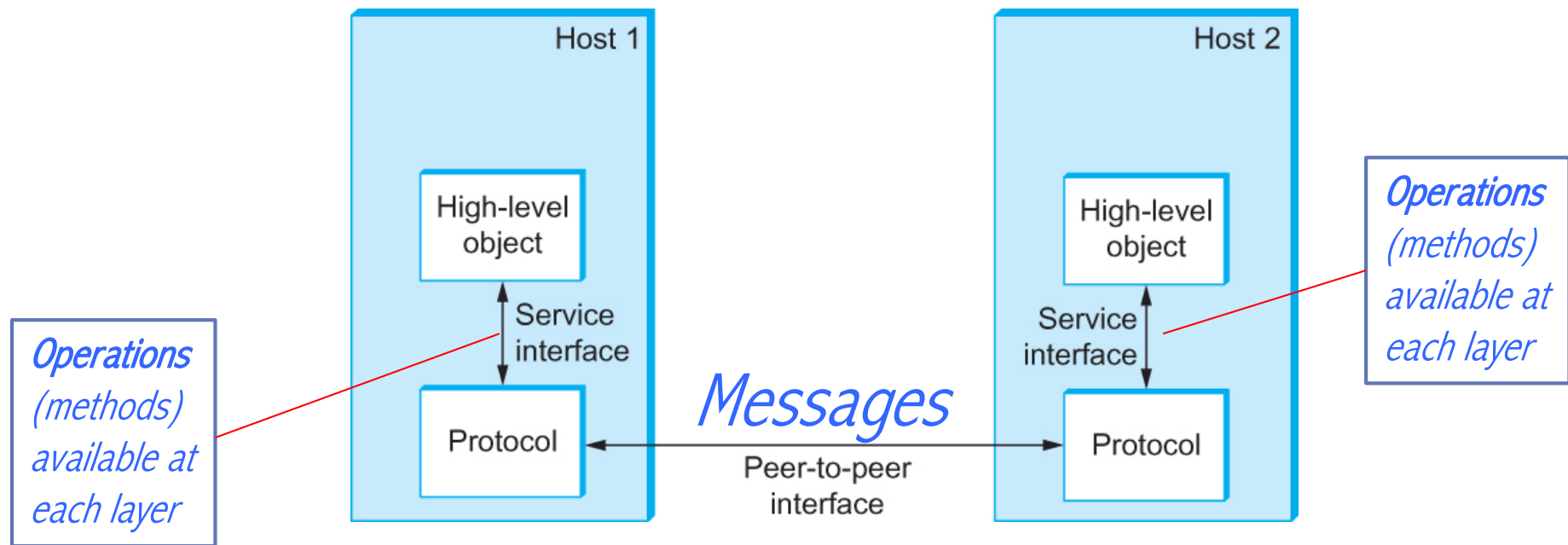
A layer uses the services provided by the layer below

- Implementation details remain hidden
 - ▣ Layer N+1 knows nothing about the implementation of layer N
- Service use via Service Interface, only
- OSI 7-layer model



Interfaces and Protocols

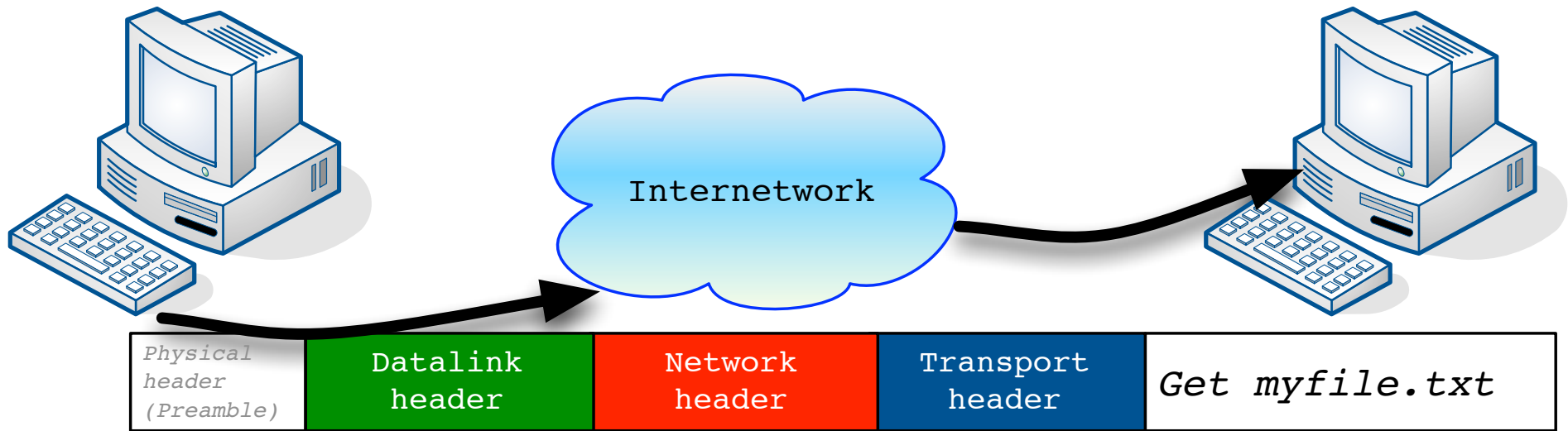
- □ This happens at any two-layer interface
 - ▣ 3/2
 - ▣ 2/1



Protocols

- - Protocol Specification: prose, pseudo-code, state transition diagram
 - Interoperable: when two or more protocols that implement the specification accurately
 - IETF: Internet Engineering Task Force

Illustration of encapsulation in TCP/IP





Get myfile.txt

Application

Presentation

Session

Transport

Network

Datalink

Physical



Get myfile.txt

Get myfile.txt

Application

Presentation

Session

Transport

Network

Datalink

Physical



Get myfile.txt

Application

Get myfile.txt

Presentation

Get myfile.txt

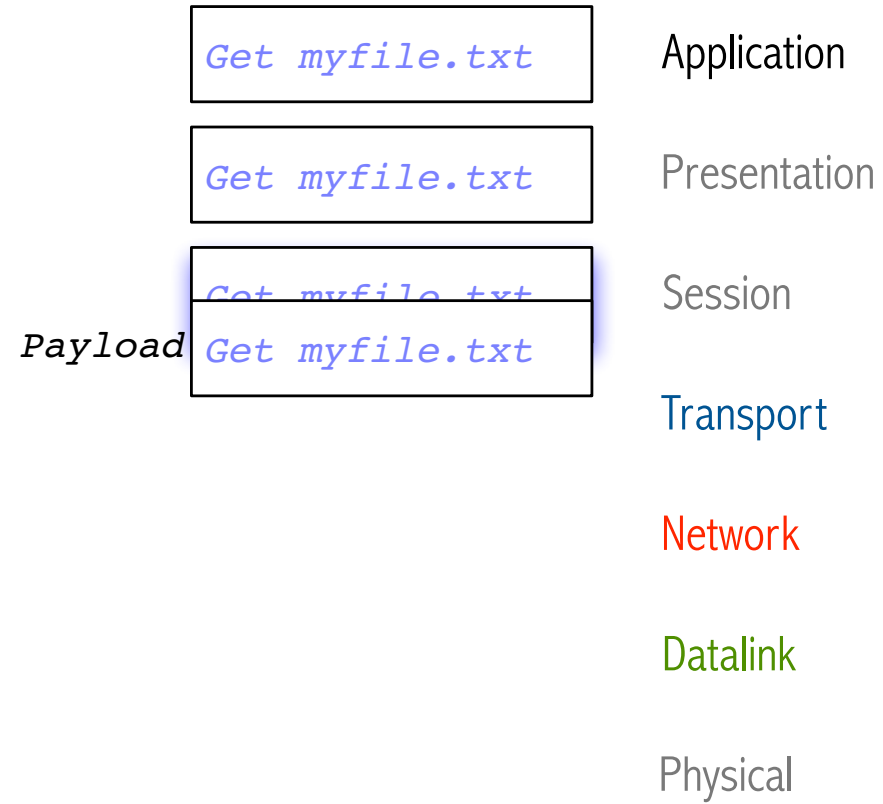
Session

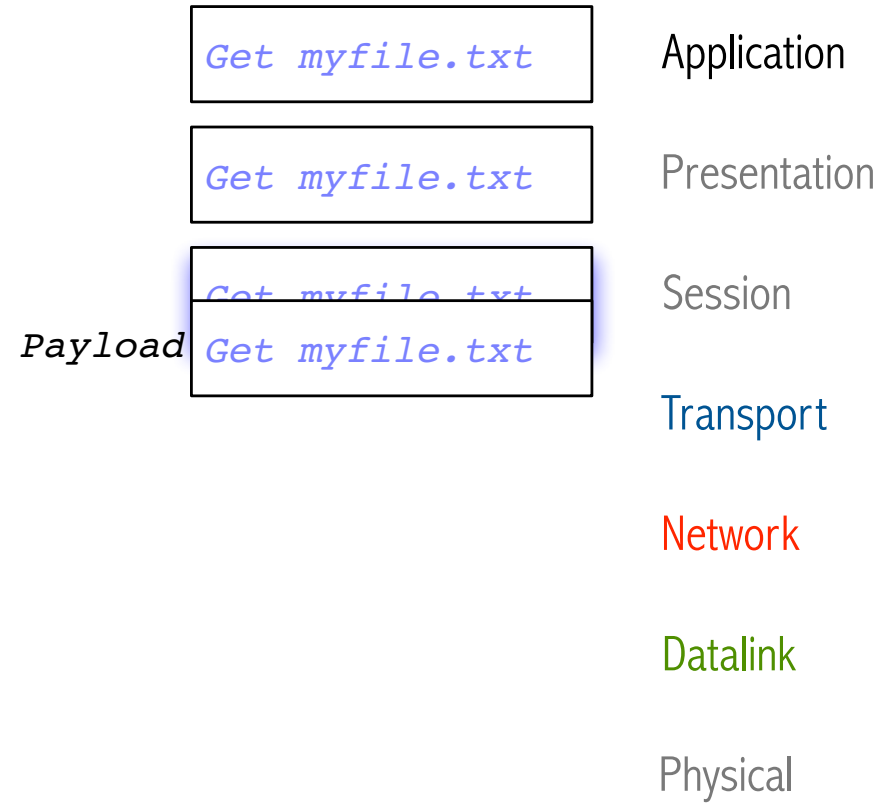
Transport

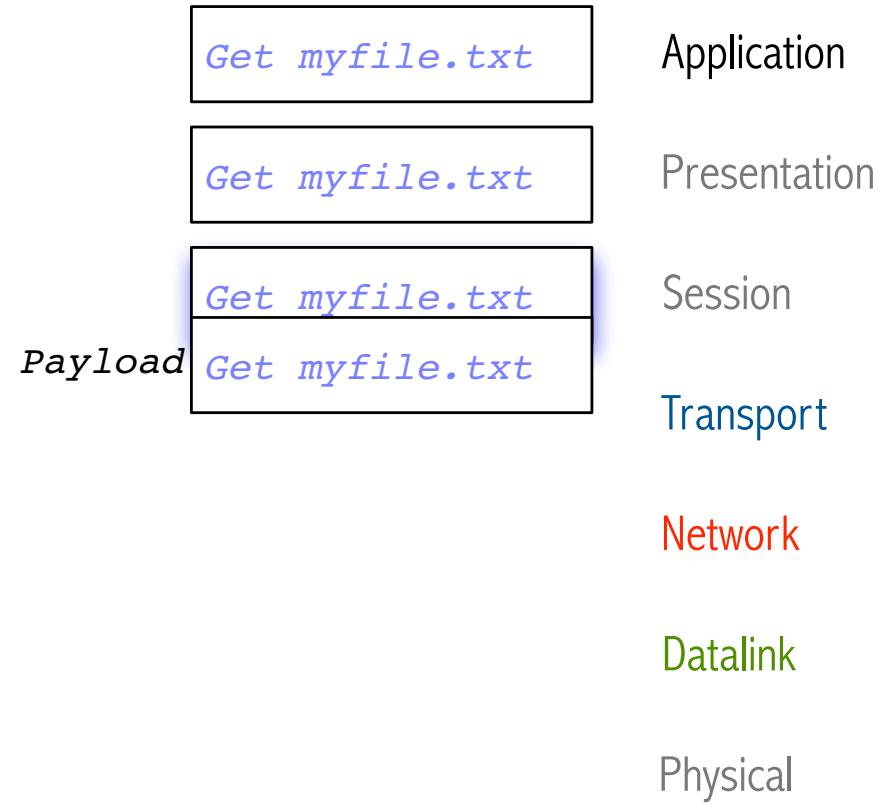
Network

Datalink

Physical









Get myfile.txt

Application

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Transport

Network

Datalink

Physical



Transport
Header

Get myfile.txt

Get myfile.txt

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Application

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Transport
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Get myfile.txt

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Get myfile.txt

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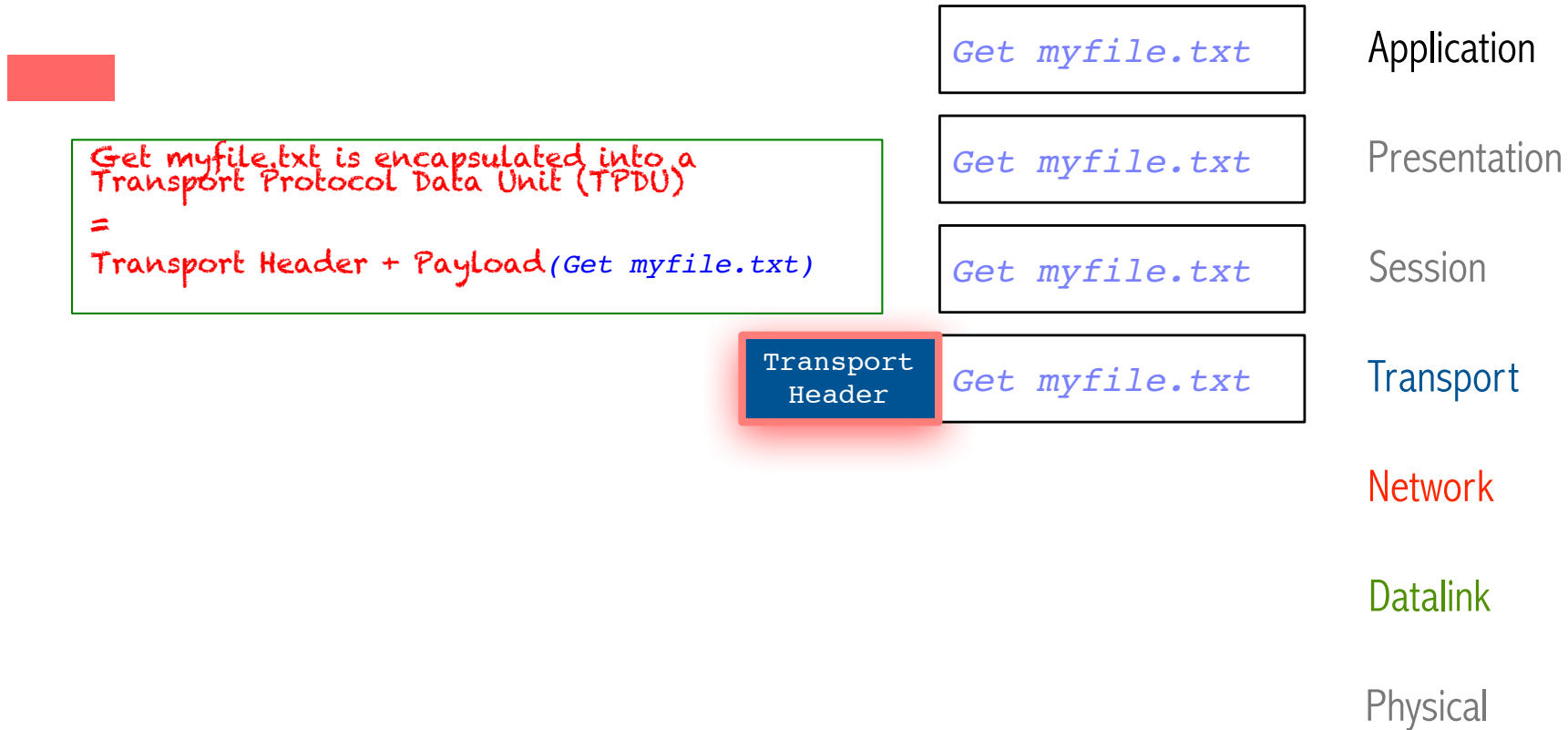
Transport

Network

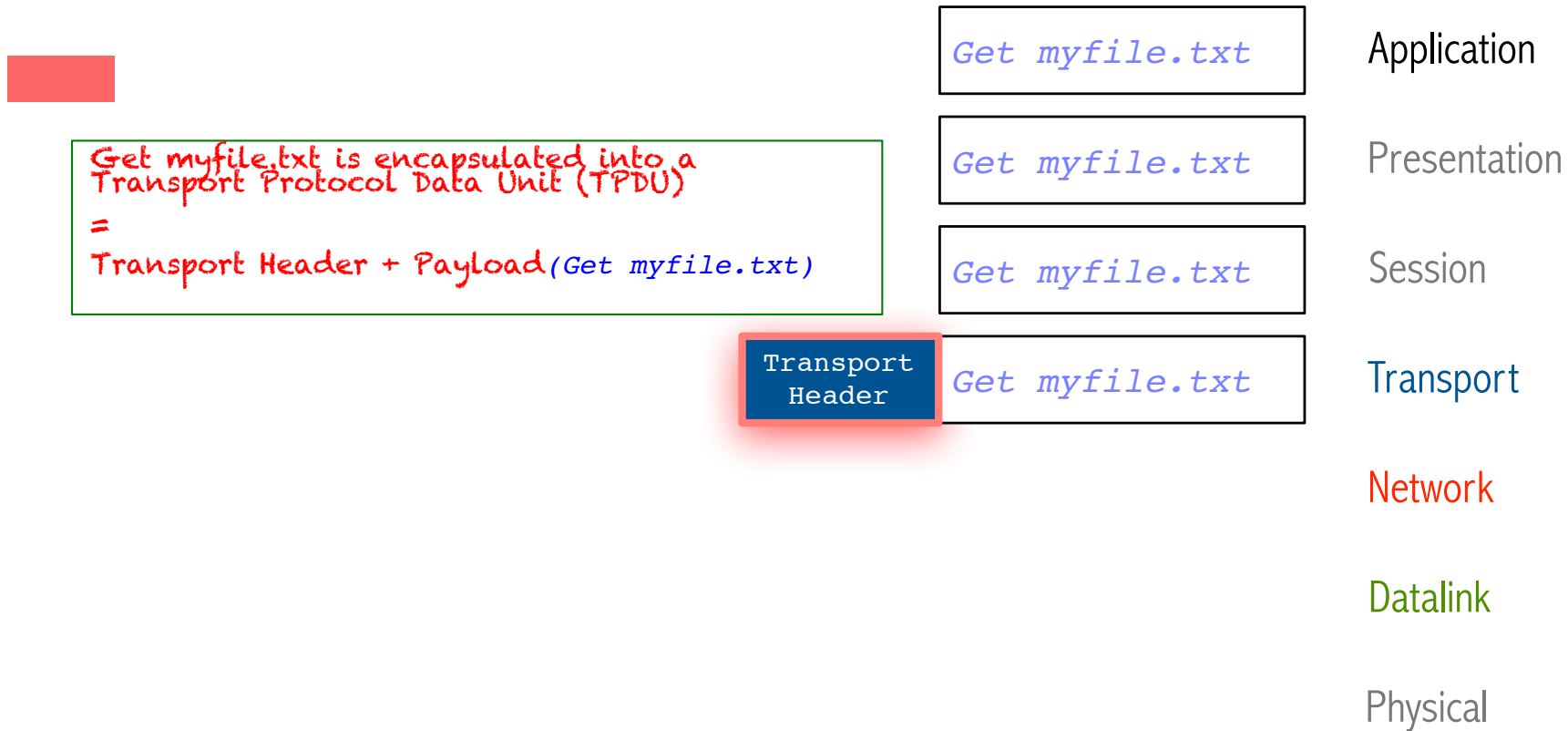
Datalink

Physical

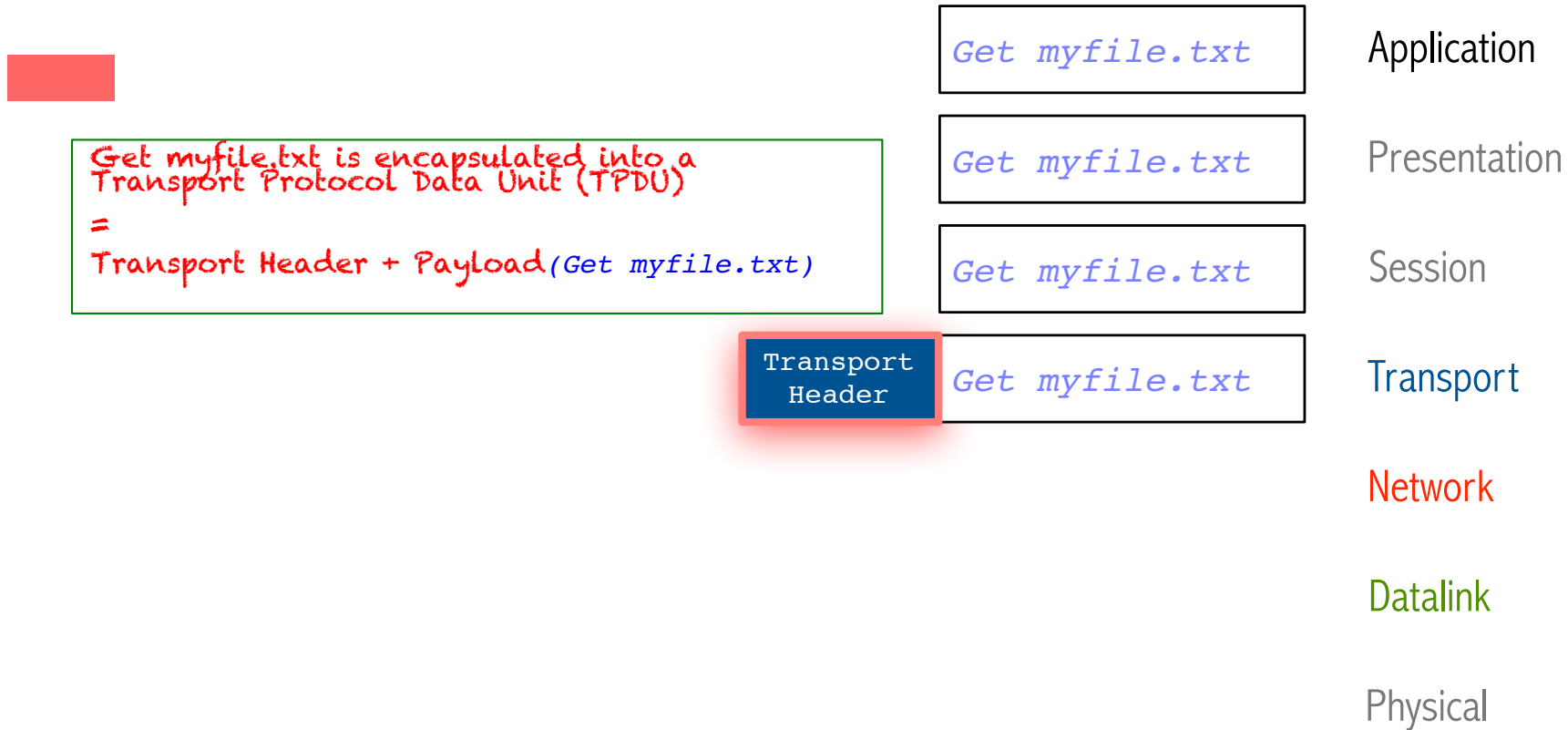
Encapsulation at layer 4

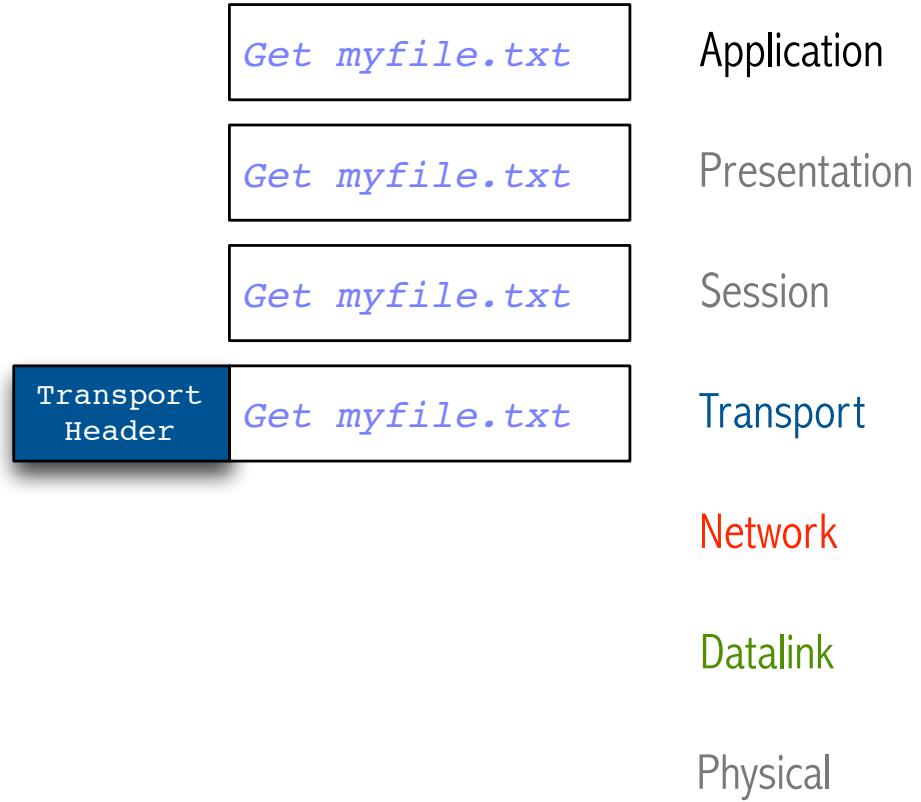


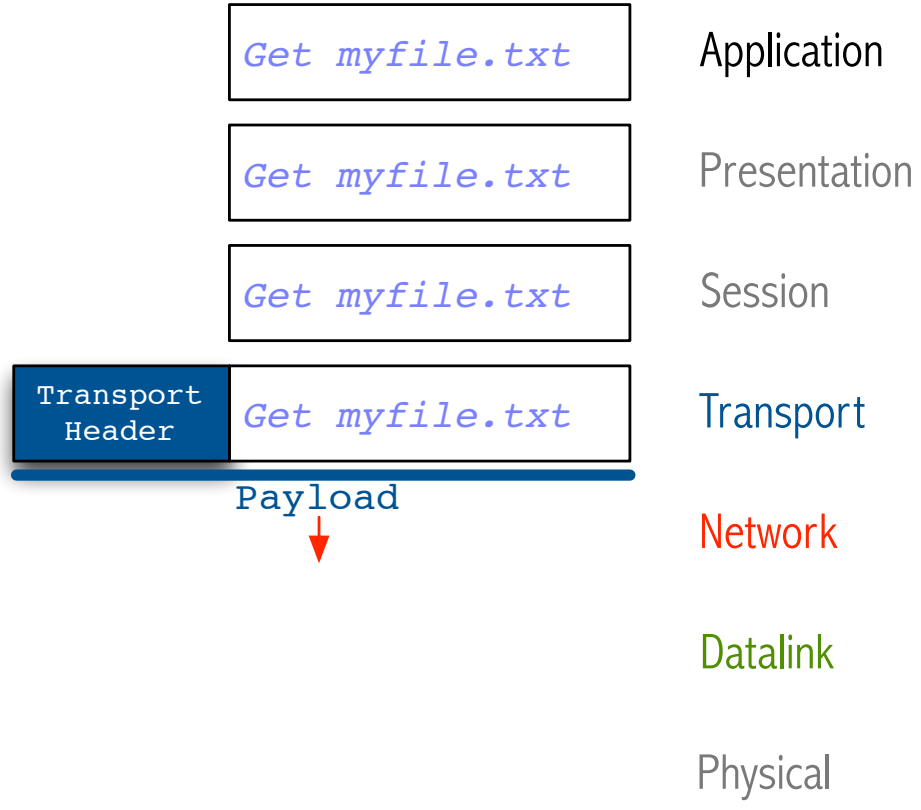
Encapsulation at layer 4

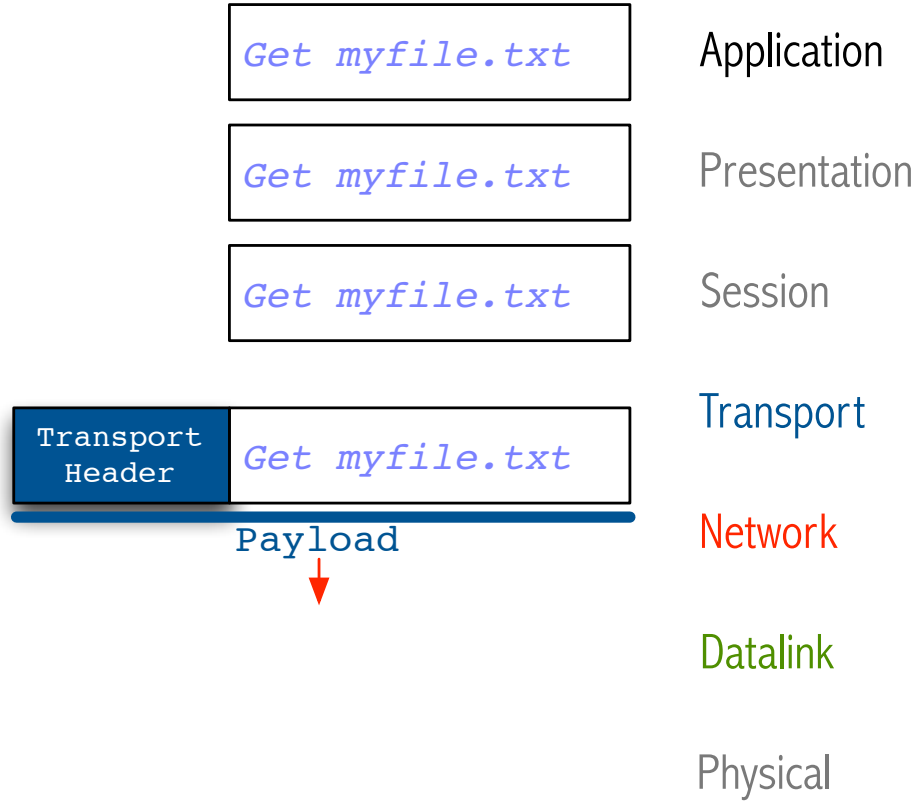


Encapsulation at layer 4











Get myfile.txt

Application

Get myfile.txt

Presentation

Get myfile.txt

Session

Transport

Transport Header	<i>Get myfile.txt</i>
------------------	-----------------------

Network

Payload

Datalink

Physical



Get myfile.txt

Application

Get myfile.txt

Presentation

Get myfile.txt

Session

Transport

Network
Header

Transport
Header

Get myfile.txt

Network

Datalink

Physical



Get myfile.txt

Application

Get myfile.txt

Presentation

Get myfile.txt

Session

Transport

Network
Header

Transport
Header

Get myfile.txt

Network

Datalink

Physical

Encapsulation at layer 3



Get myfile.txt

Application

Get myfile.txt

Presentation

Get myfile.txt

Session

Transport

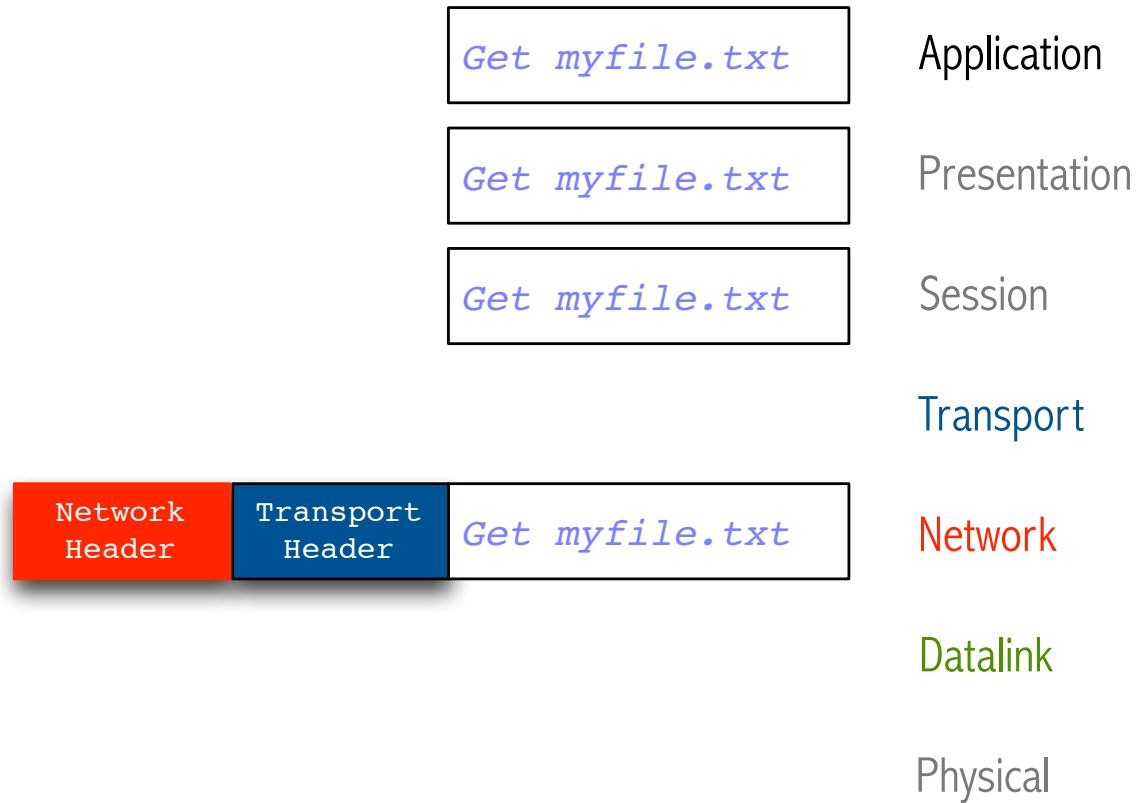
Network Header Transport Header *Get myfile.txt*

Network

Datalink

Physical

Encapsulation at layer 3





Get myfile.txt

Application

Get myfile.txt

Presentation

Get myfile.txt

Session

Transport

Network Header Transport Header *Get myfile.txt*

Network

Payload

Datalink



Physical





Get myfile.txt

Application

Get myfile.txt

Presentation

Get myfile.txt

Session

Transport



Network

Payload

Datalink



Physical





Get myfile.txt

Application

Get myfile.txt

Presentation

Get myfile.txt

Session

Transport



Network

Datalink

Physical





Get myfile.txt

Application

Get myfile.txt

Presentation

Get myfile.txt

Session

Transport

Network



Datalink

Payload



Physical





Get myfile.txt

Application

Get myfile.txt

Presentation

Get myfile.txt

Session

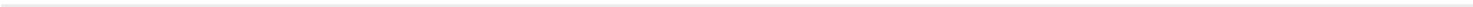
Transport

Network

Network Header Transport Header *Get myfile.txt*

Datalink

Physical





Get myfile.txt

Application

Get myfile.txt

Presentation

Get myfile.txt

Session

Transport

Network

Network Header Transport Header *Get myfile.txt*

Datalink

Physical



Get myfile.txt

Application

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Presentation

Get myfile.txt

Session

Transport

Network

Datalink
Header

Network
Header

Transport
Header

Get myfile.txt

Datalink

Physical



Get myfile.txt

Application

Get myfile.txt

Presentation

Get myfile.txt

Session

Transport

Network



Datalink

Physical

Encapsulation at layer 2



Get myfile.txt

Application

Get myfile.txt

Presentation

Get myfile.txt

Session

Transport

Network

Datalink Header Network Header Transport Header *Get myfile.txt*

Datalink

Physical

Encapsulation at layer 2



Get myfile.txt

Application

Get myfile.txt

Presentation

Get myfile.txt

Session

Transport

Network

Datalink Header Network Header Transport Header *Get myfile.txt*

Datalink

Physical

Encapsulation at layer 2



Application

Presentation

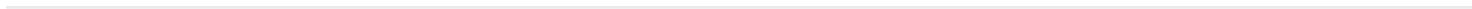
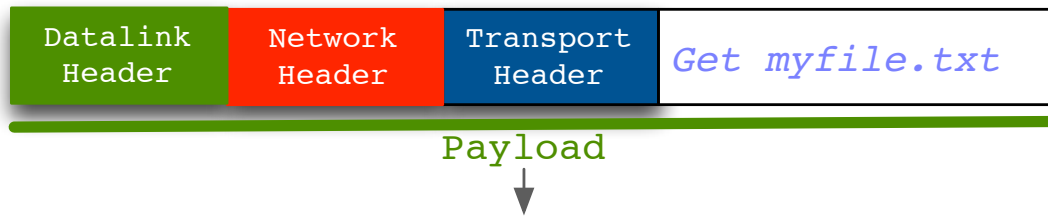
Session

Transport

Network

Datalink

Physical





Application

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Network

Datalink

Physical



Payload





Application

Presentation

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Datalink

Physical





Application

Presentation

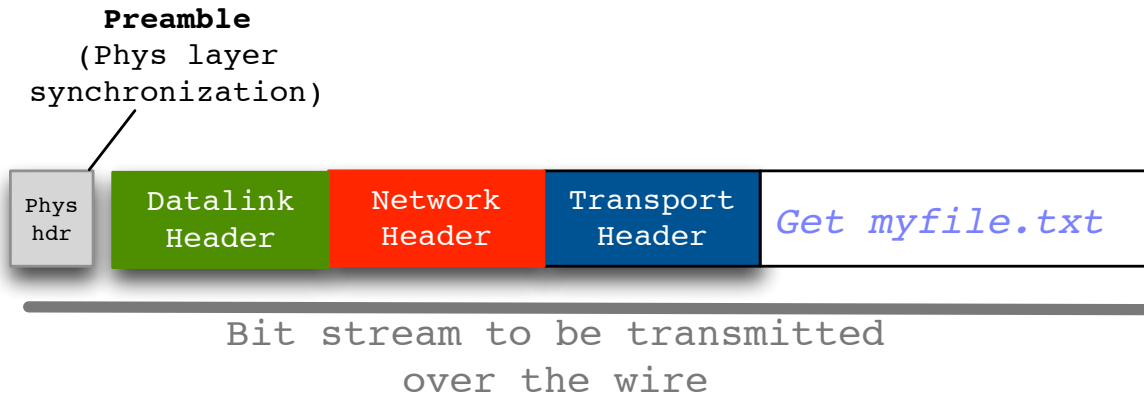
Session

Transport

Network

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Application

Presentation

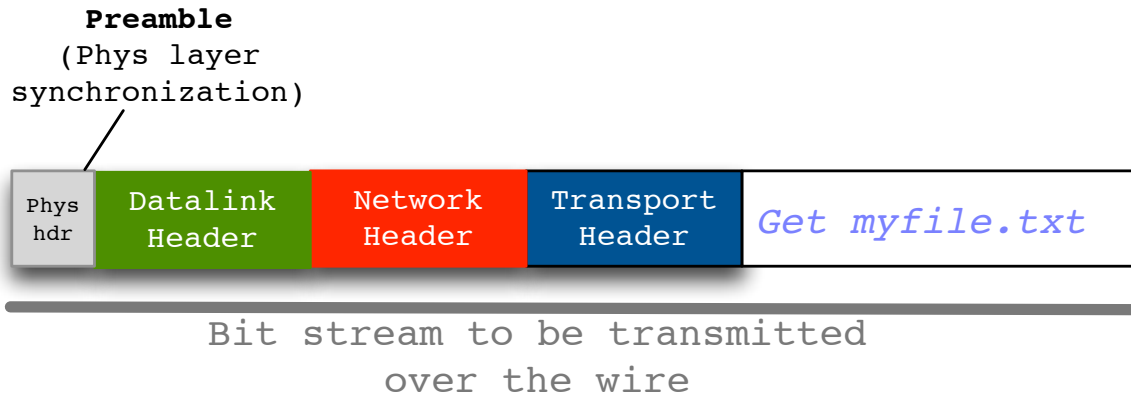
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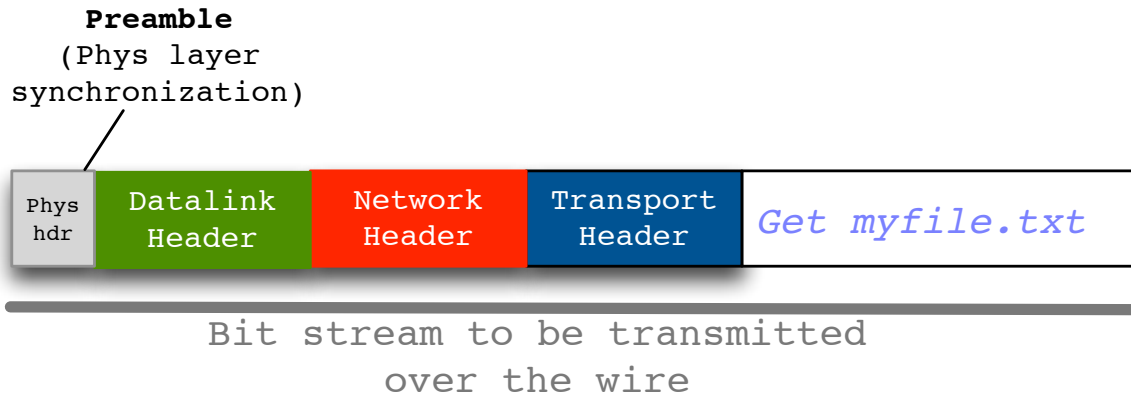
Session

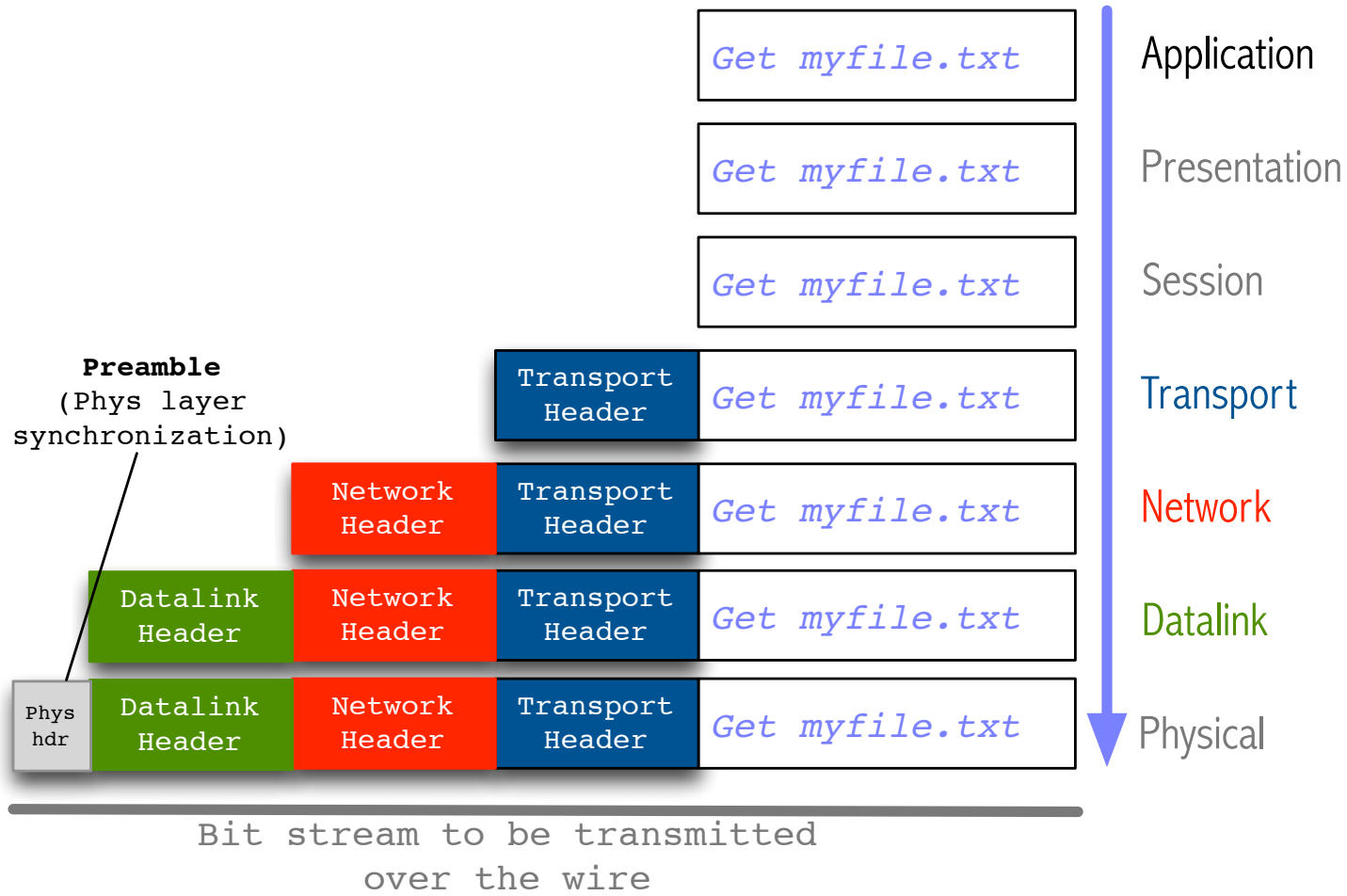
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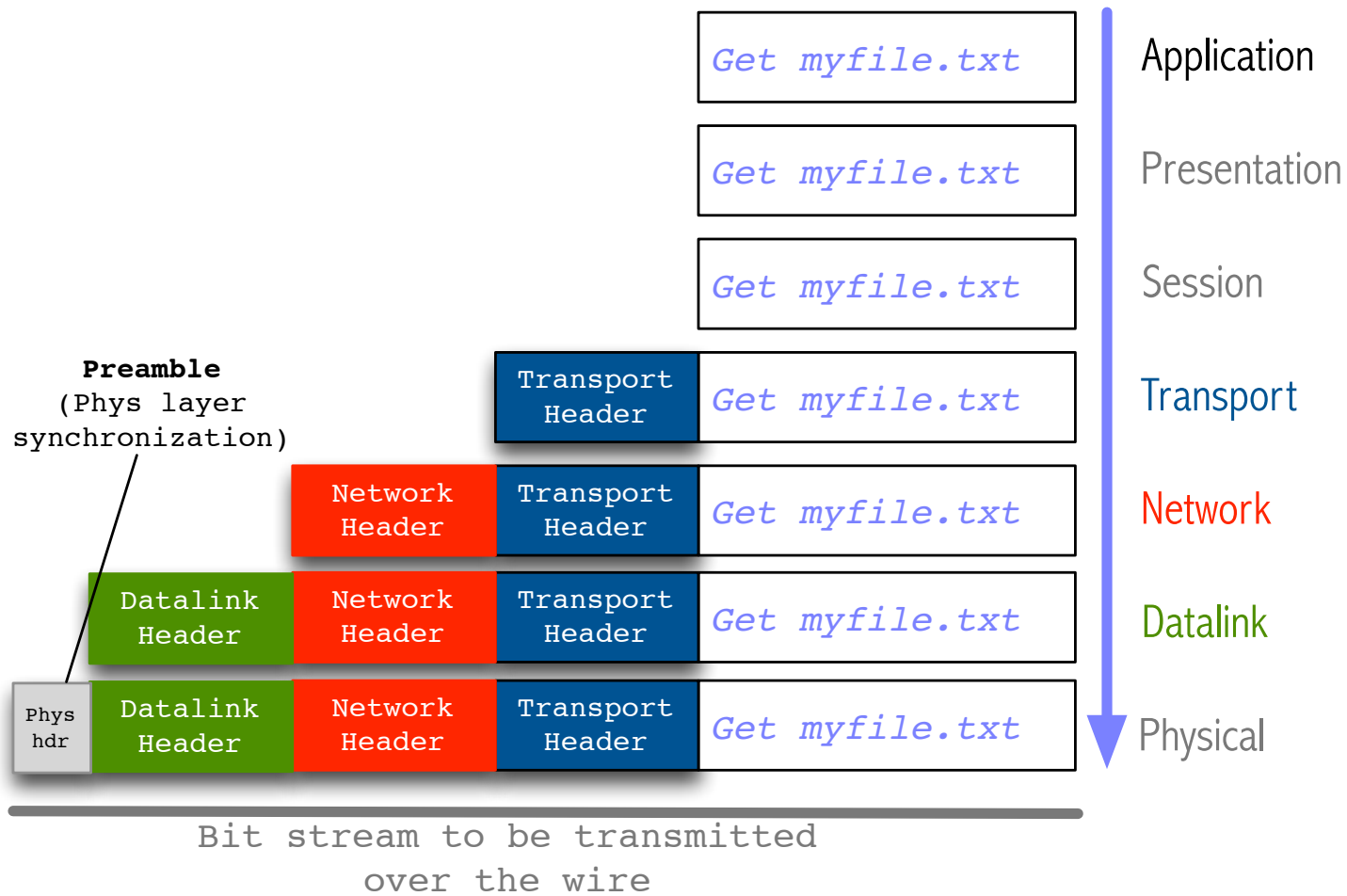
Network

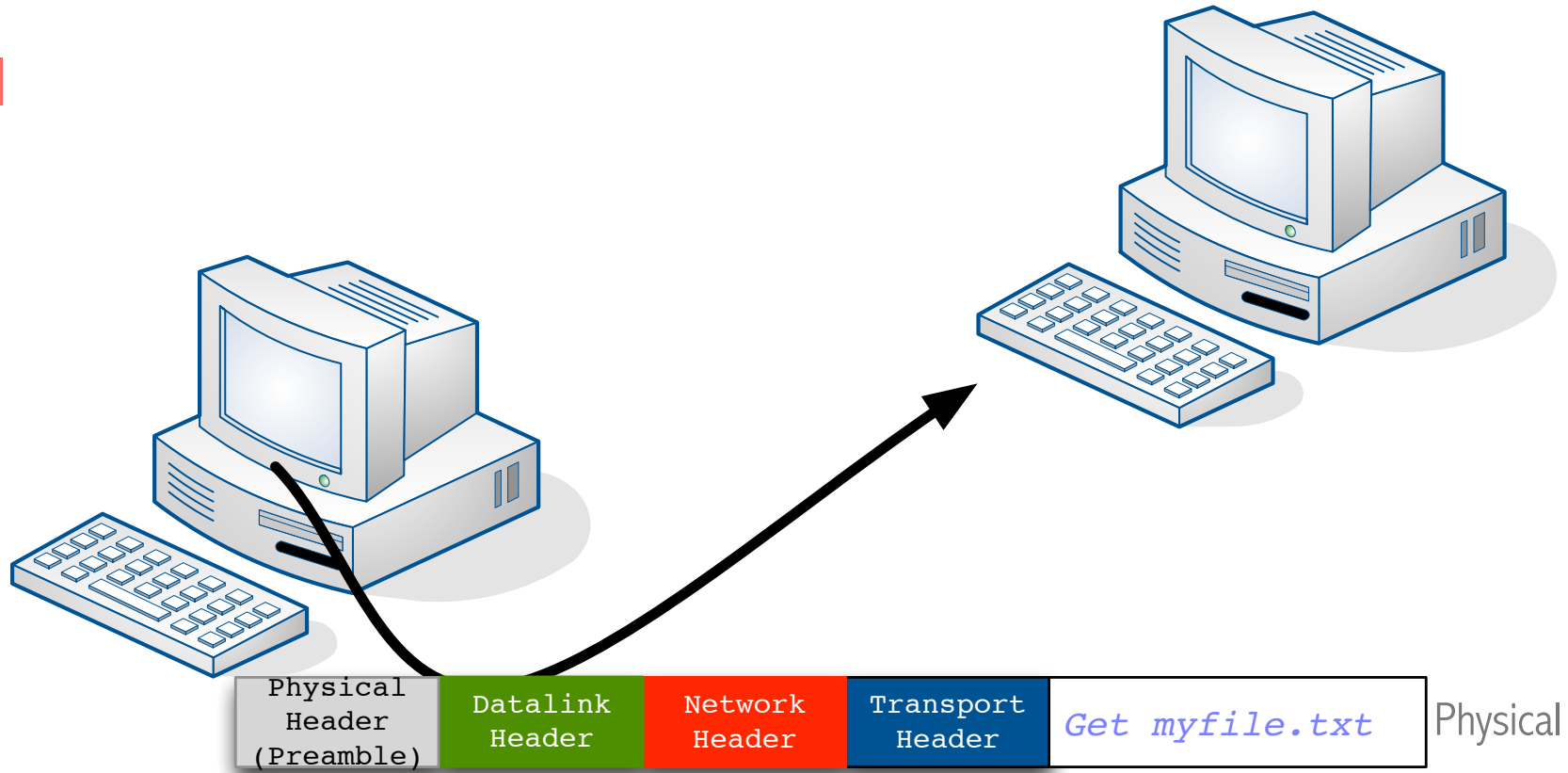
Datalink

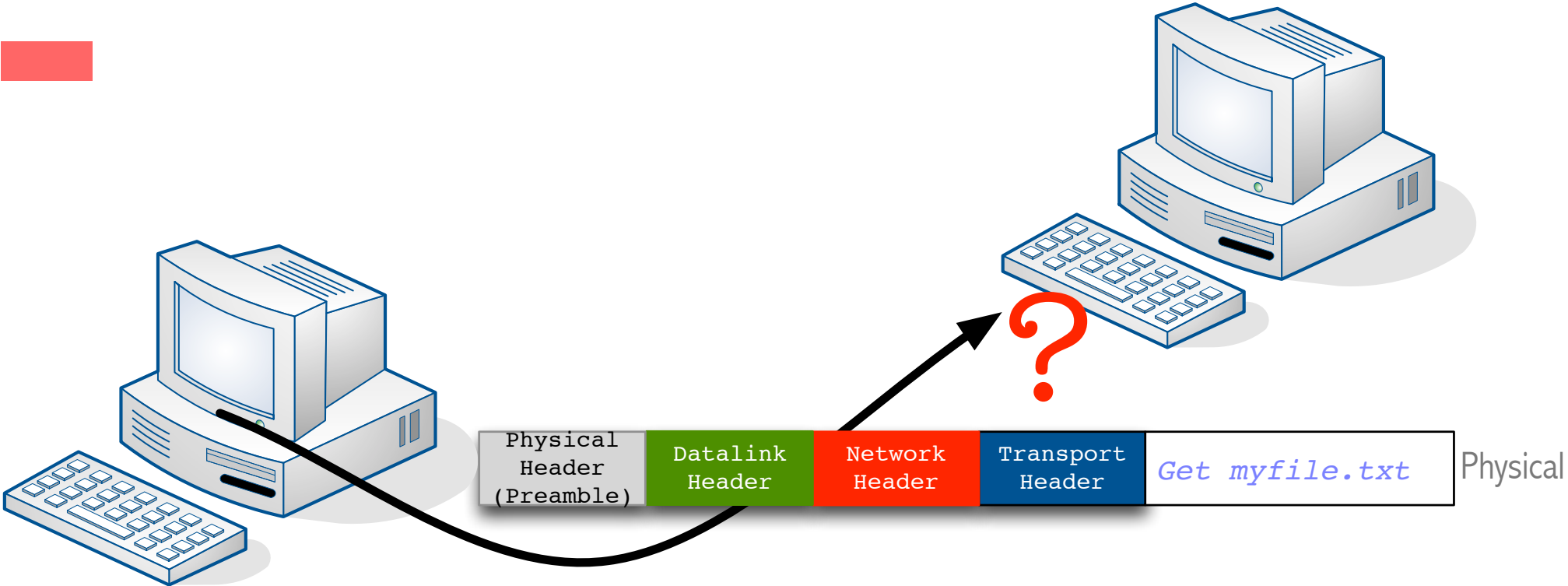
Physical

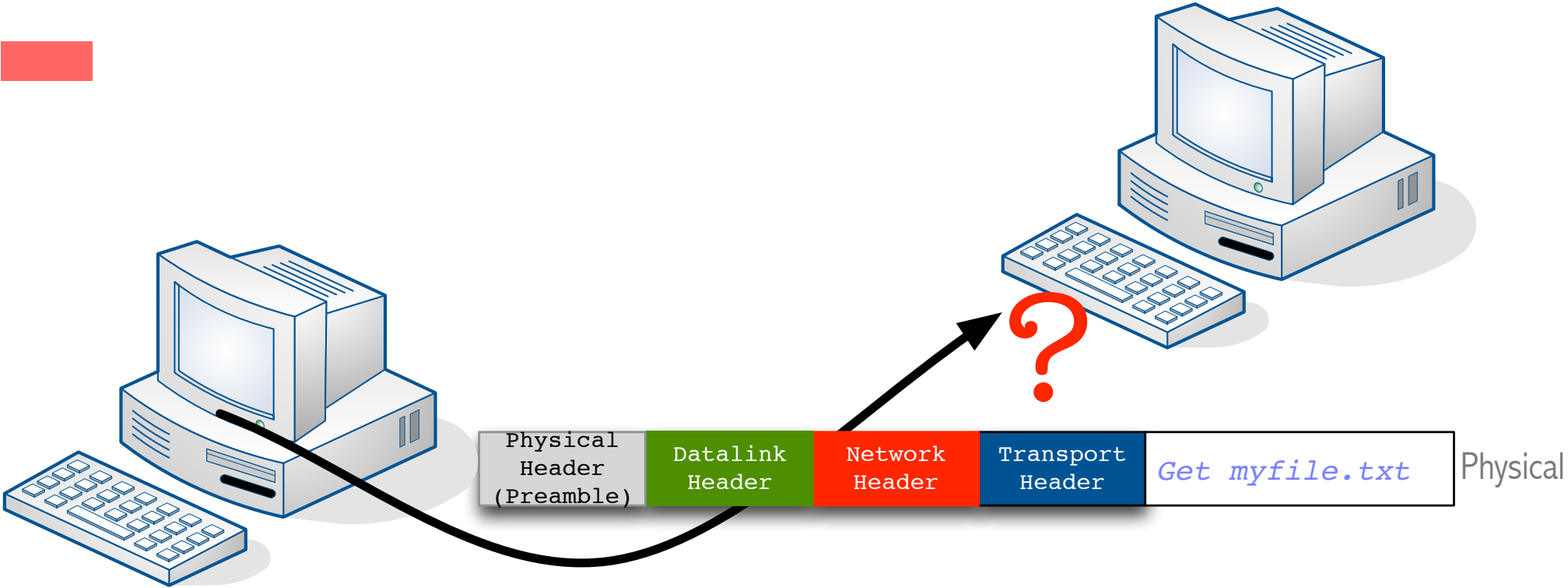




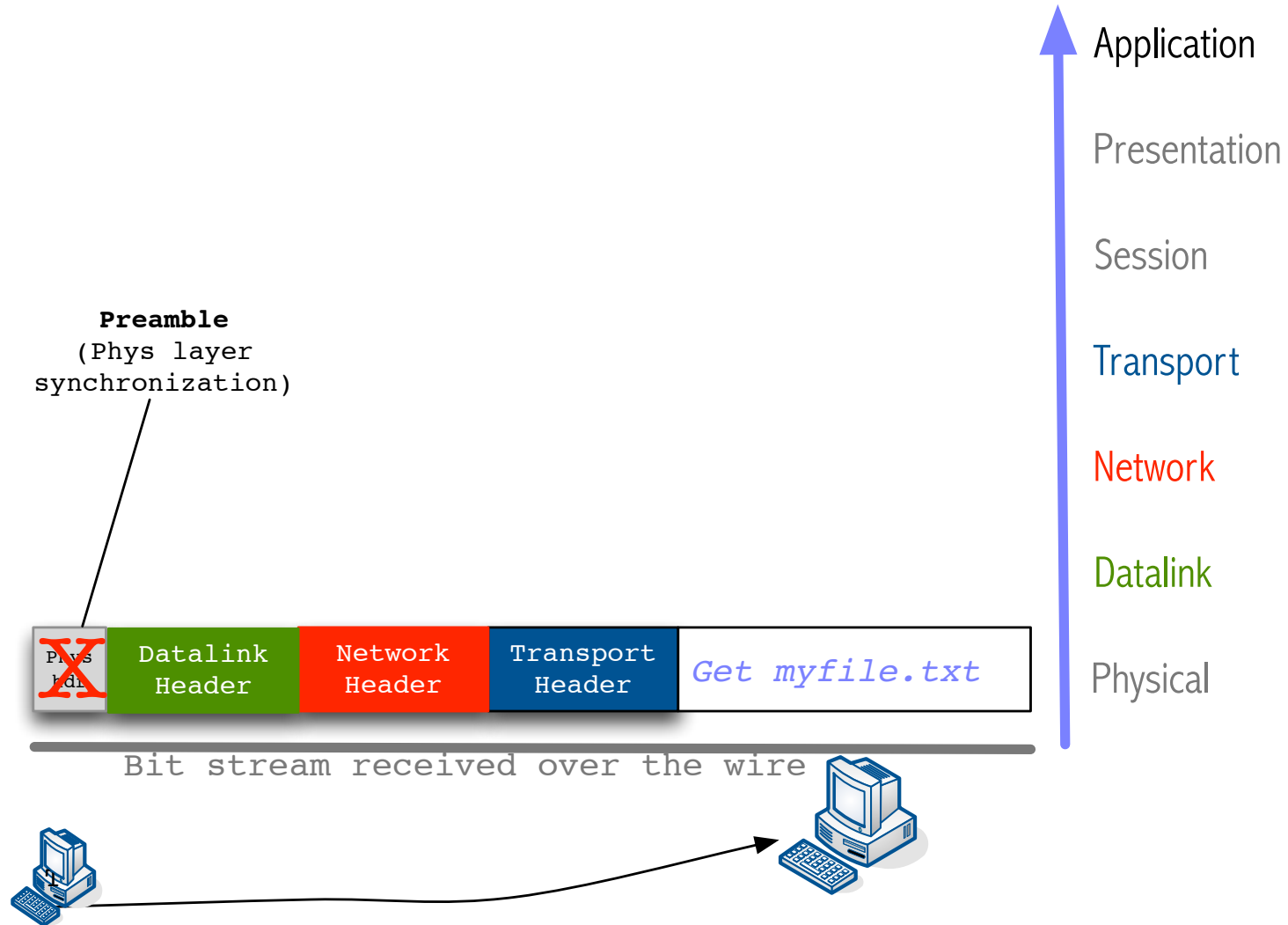




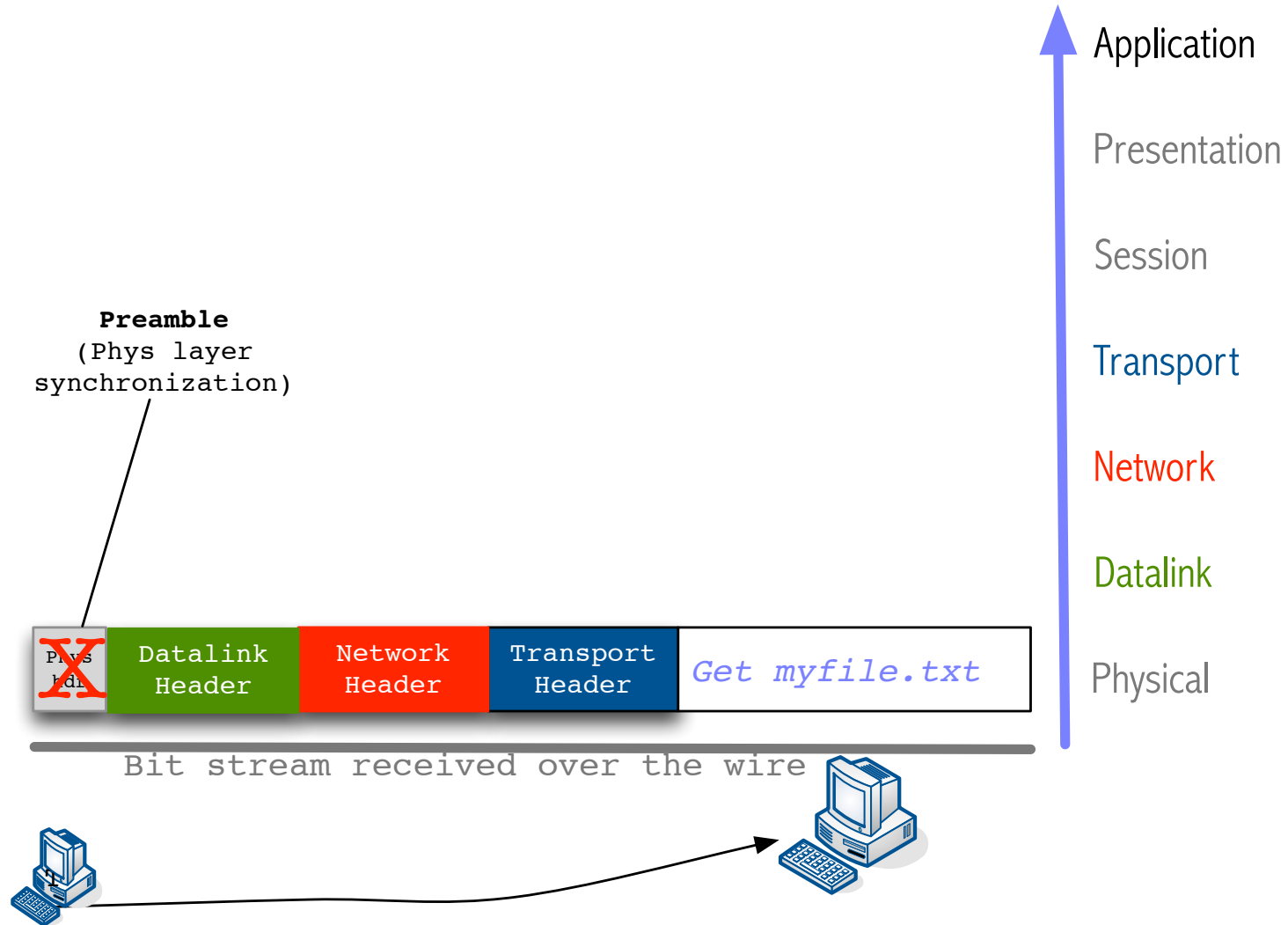


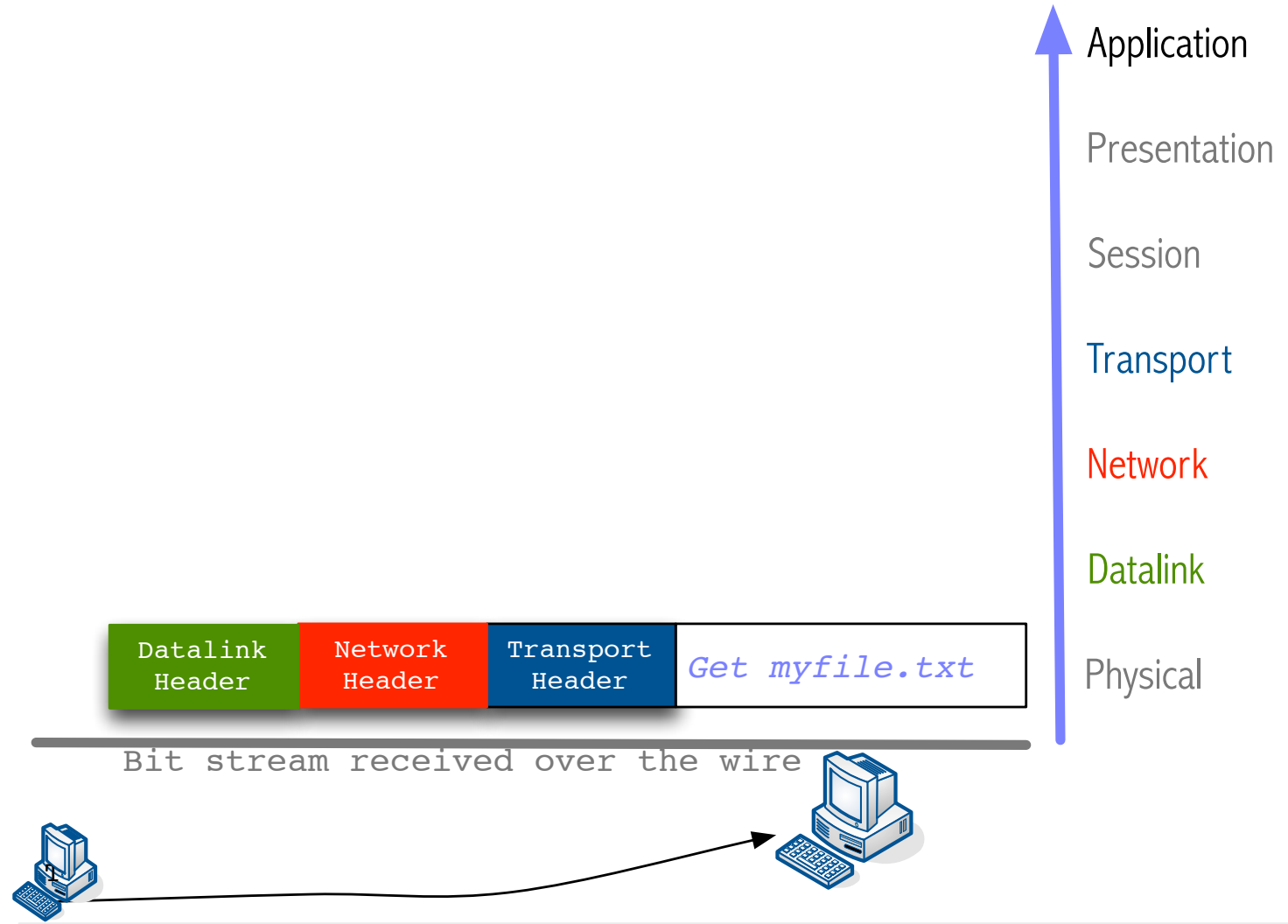


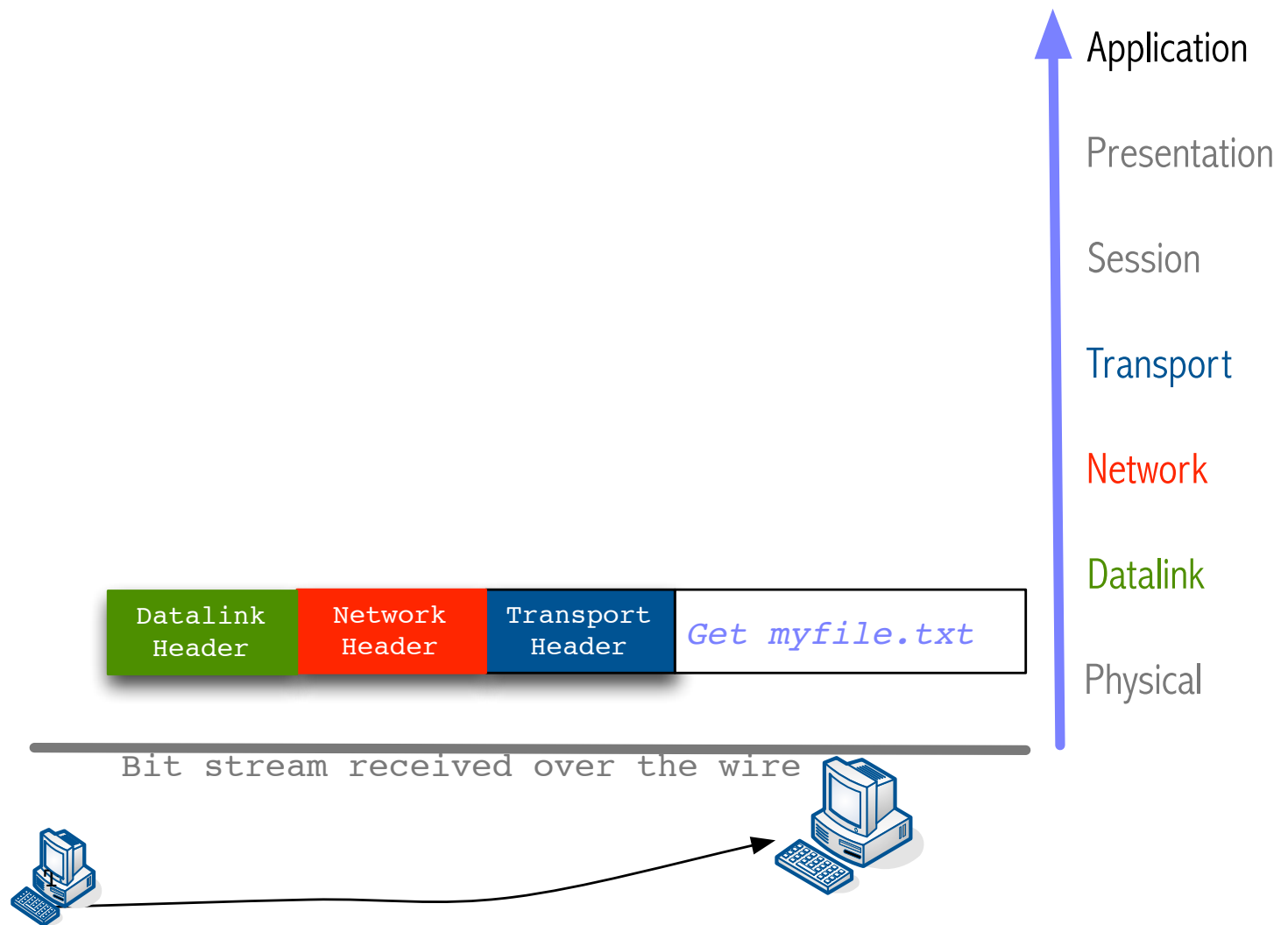
De-encapsulation at layer 1 (Supress the preamble)

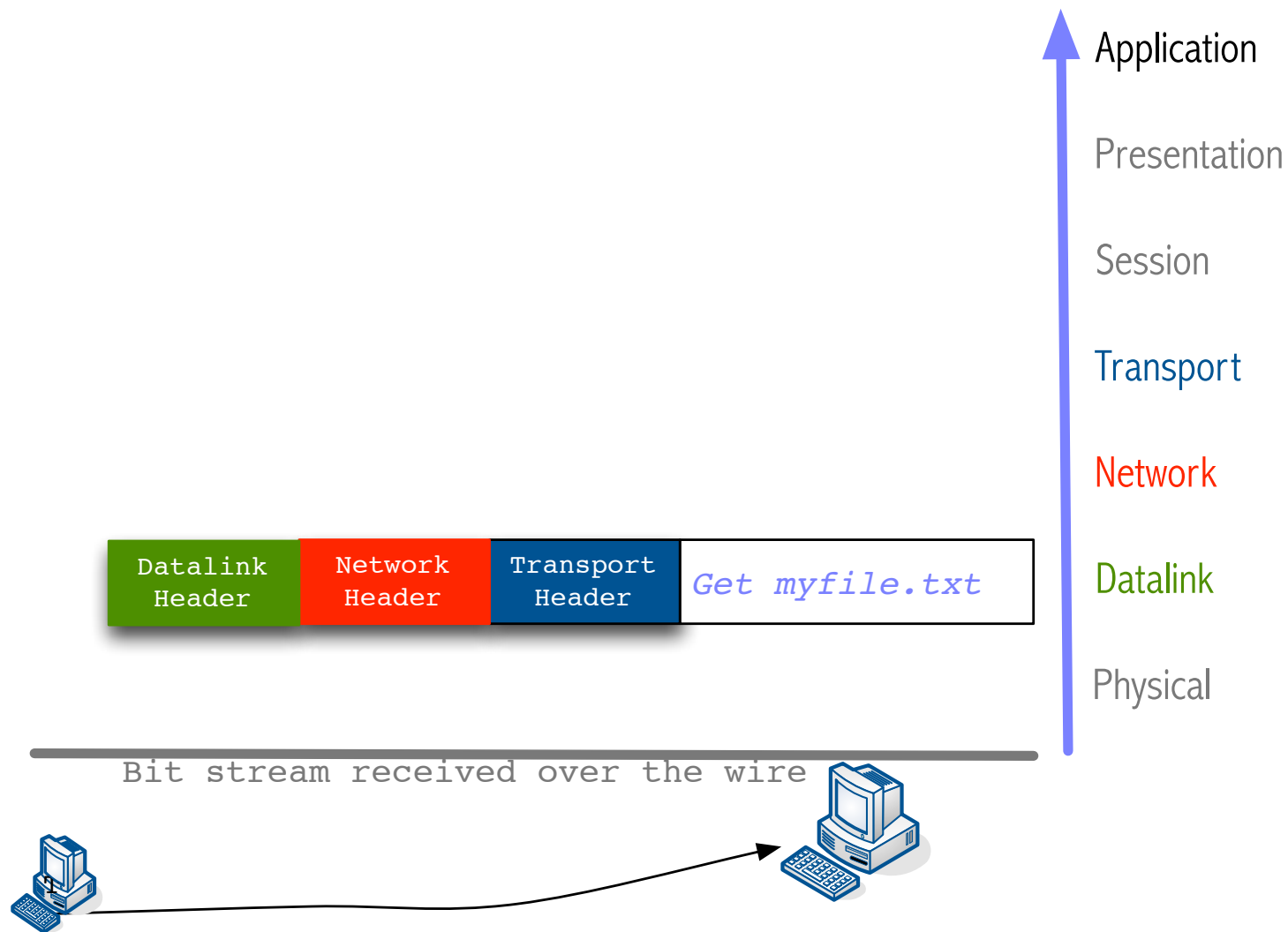


De-encapsulation at layer 1 (Supress the preamble)

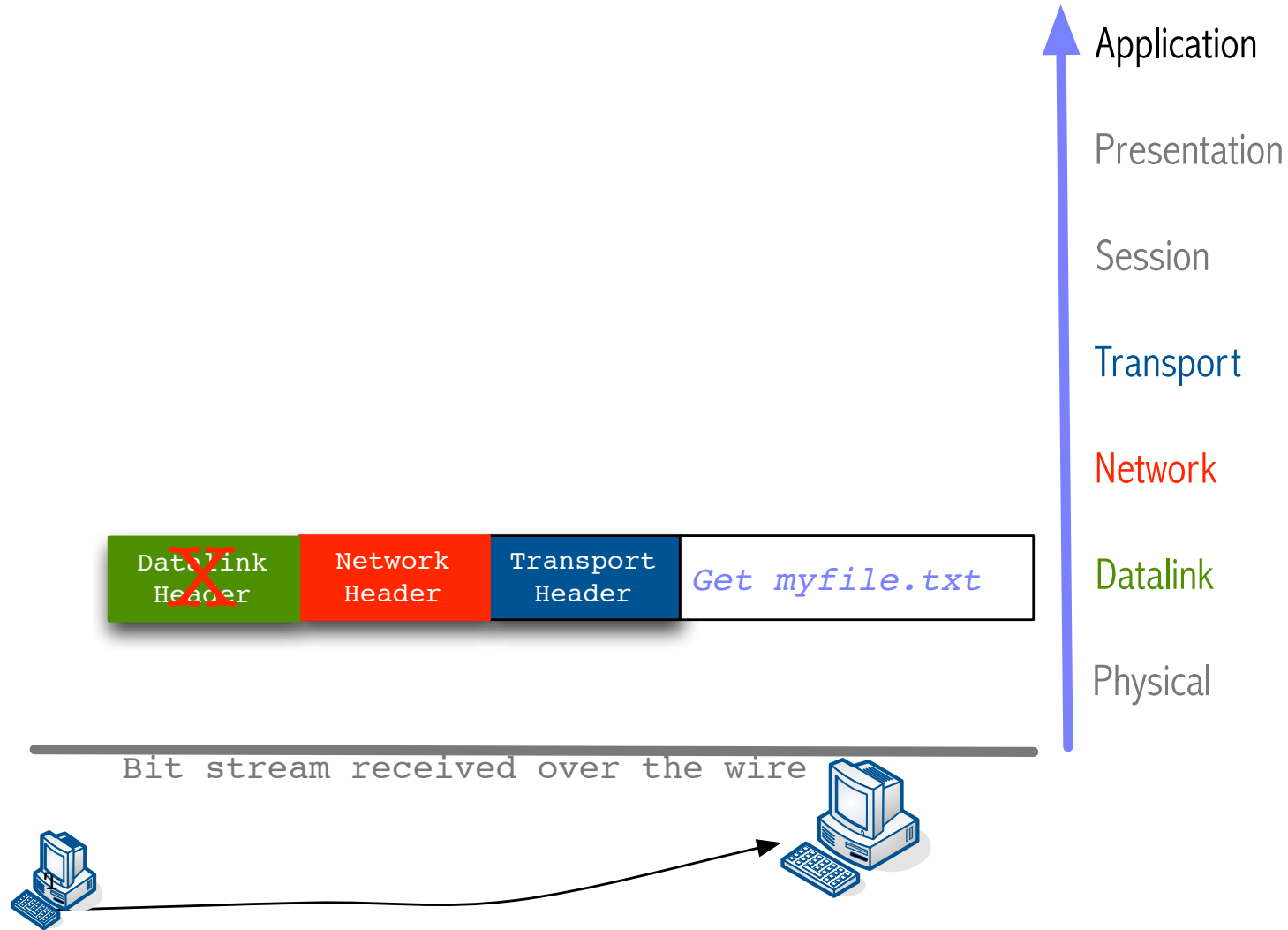


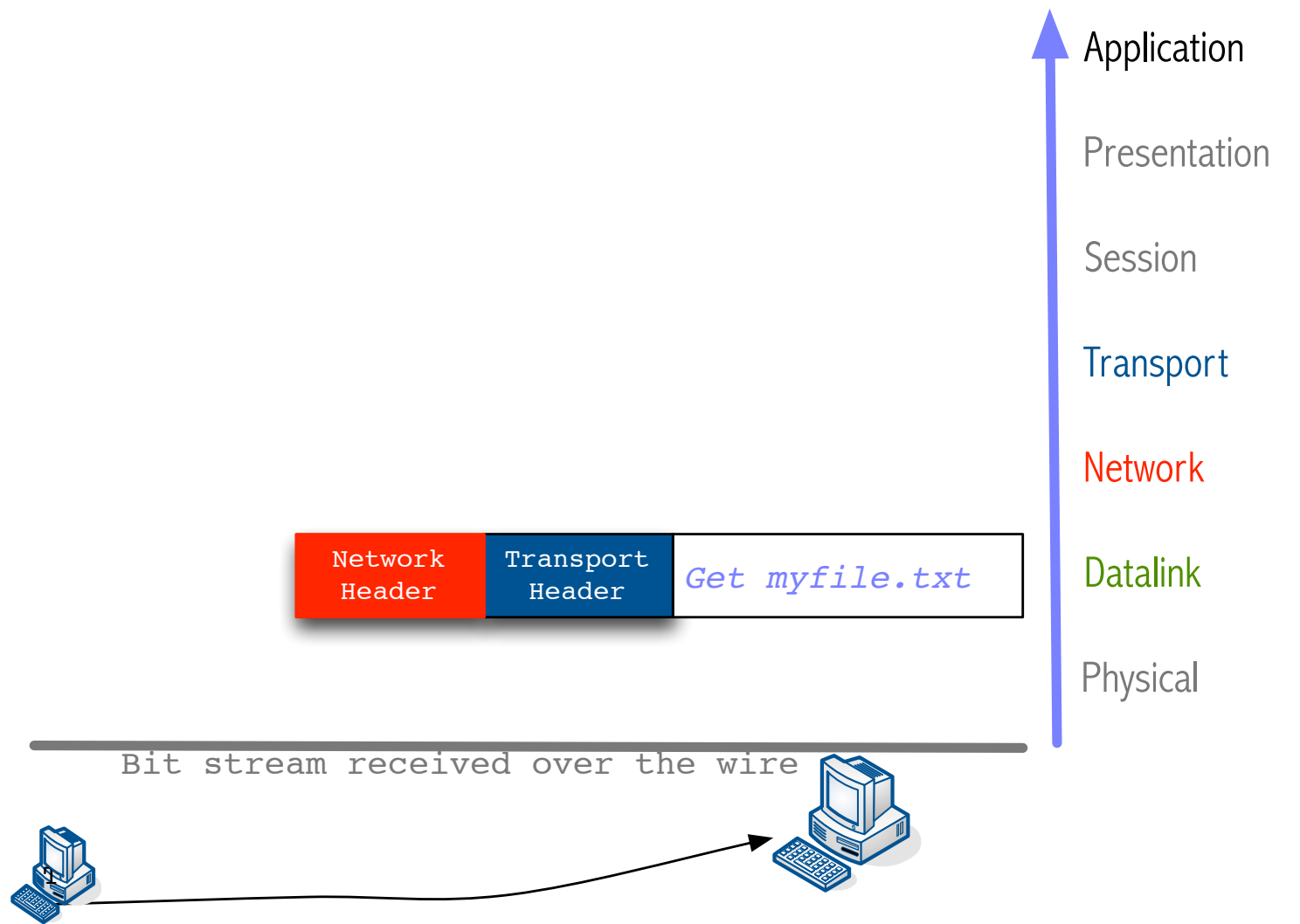


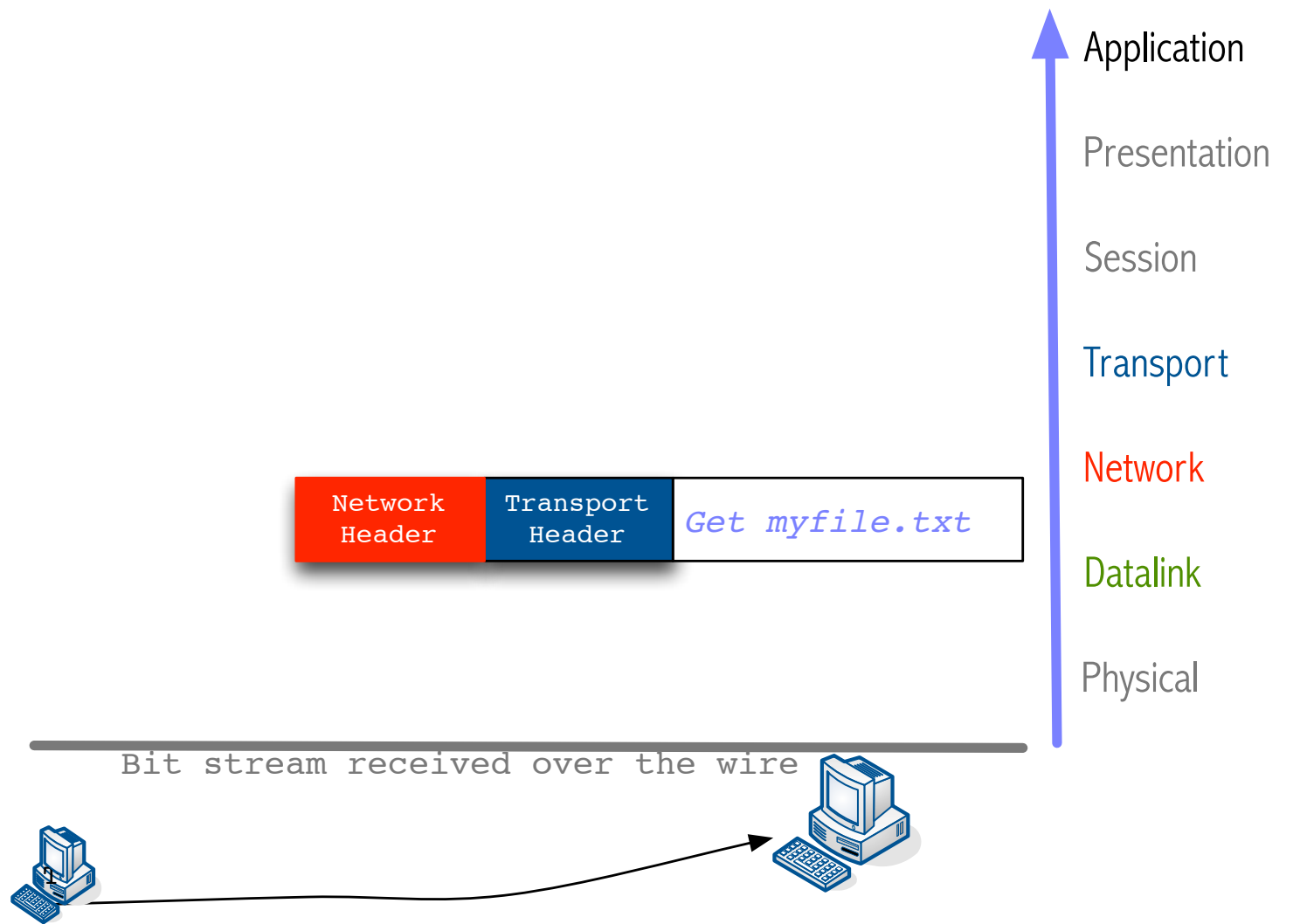


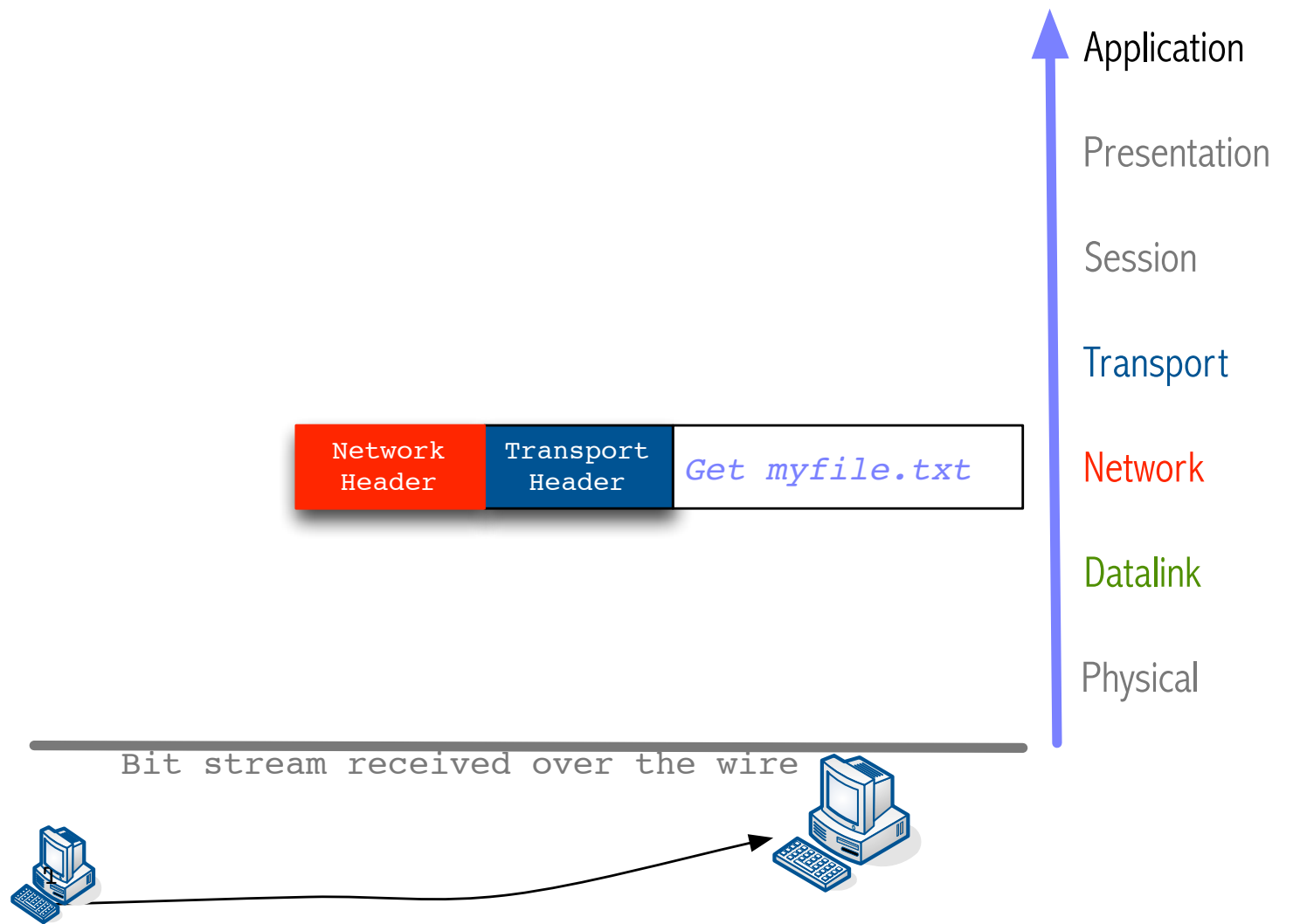


De-encapsulation at layer 2

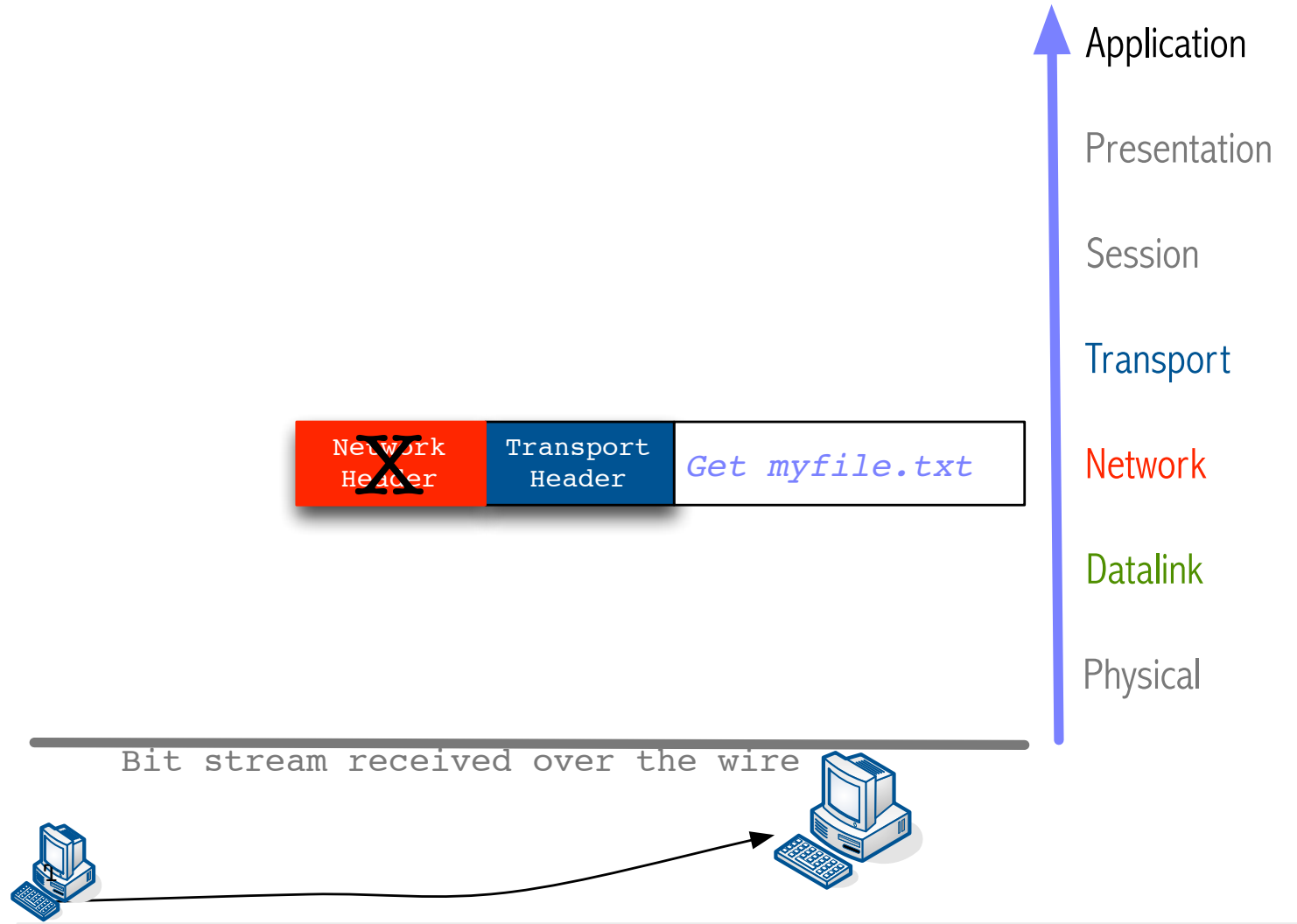


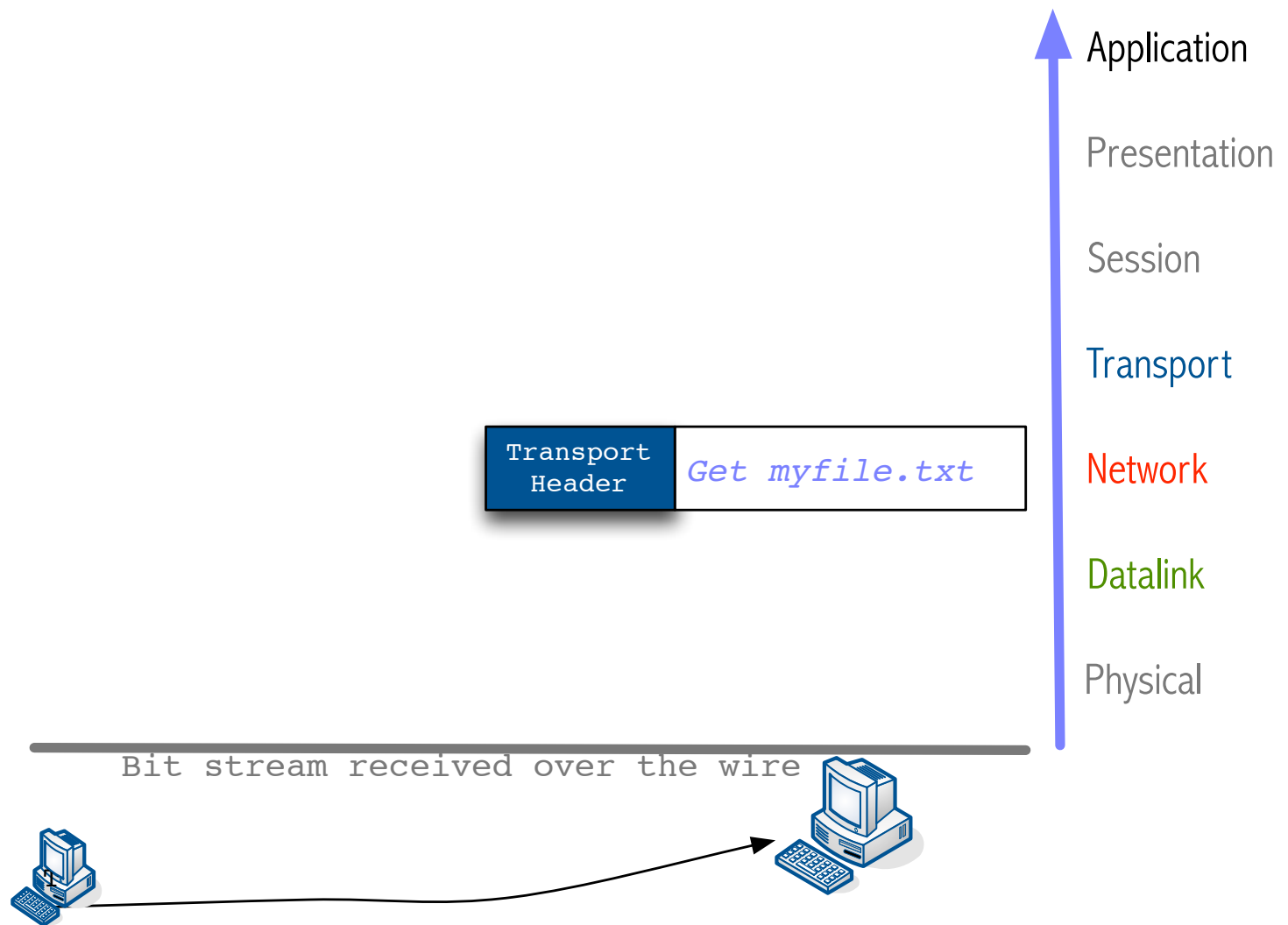


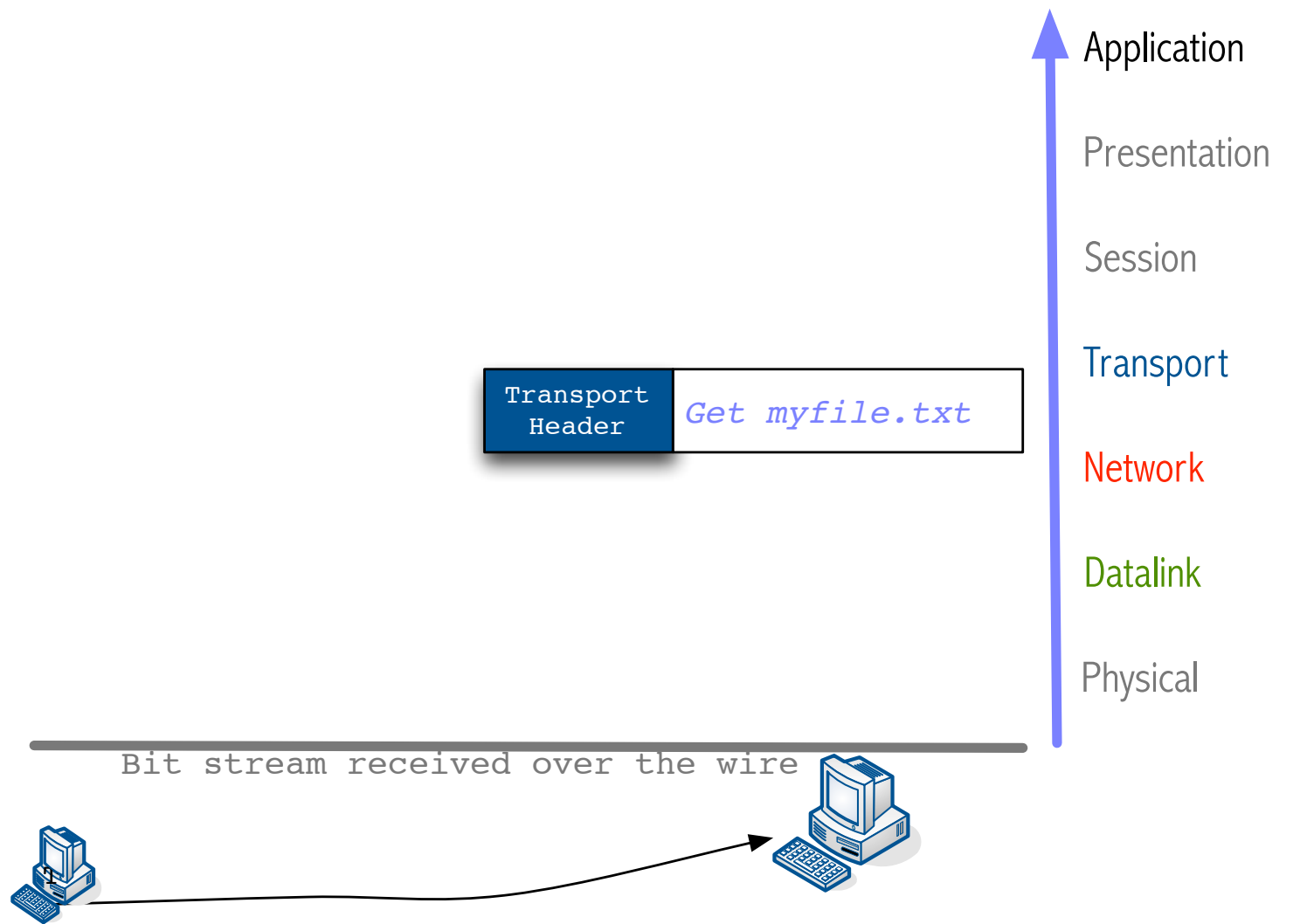


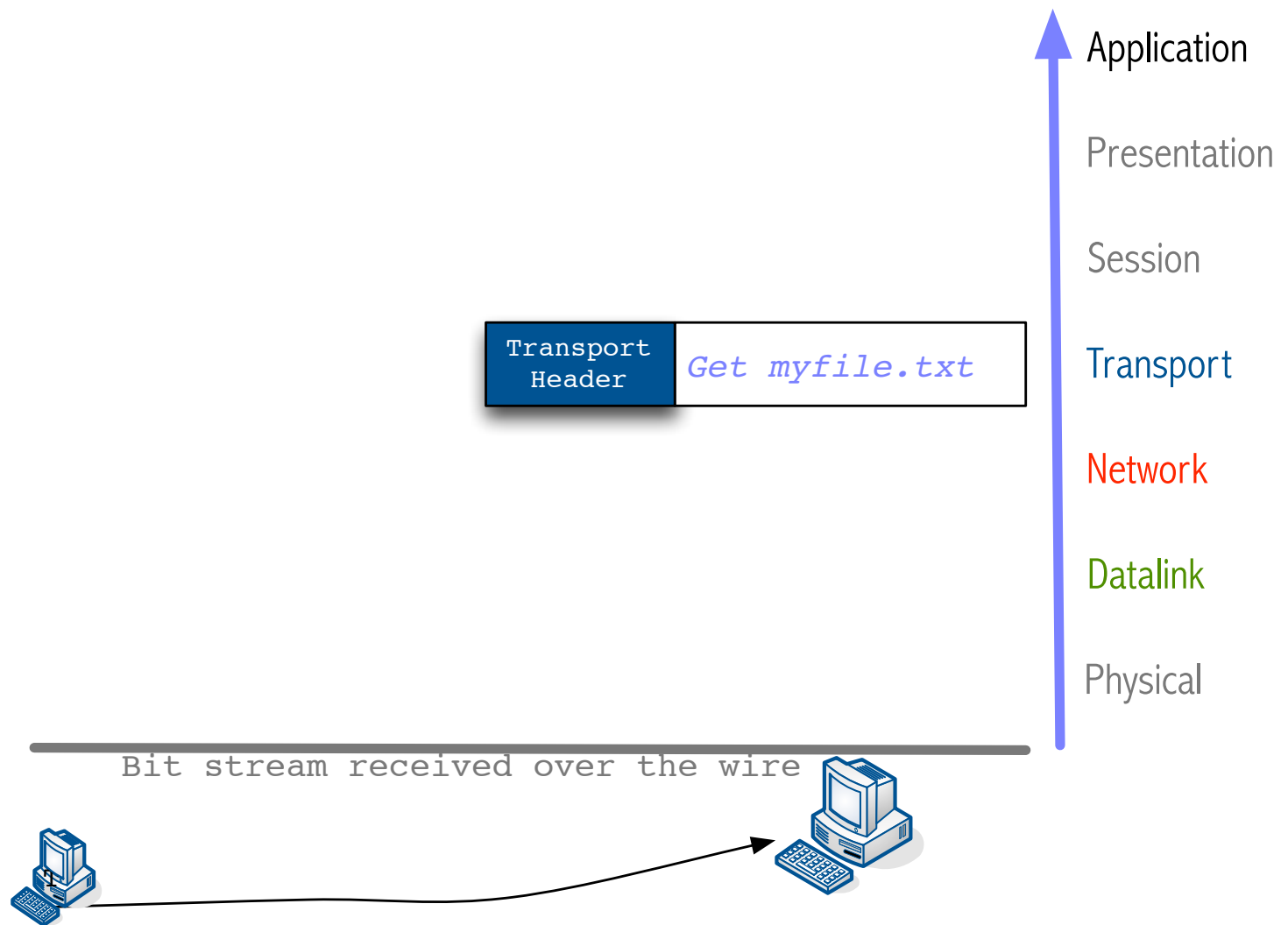


De-encapsulation at layer 3

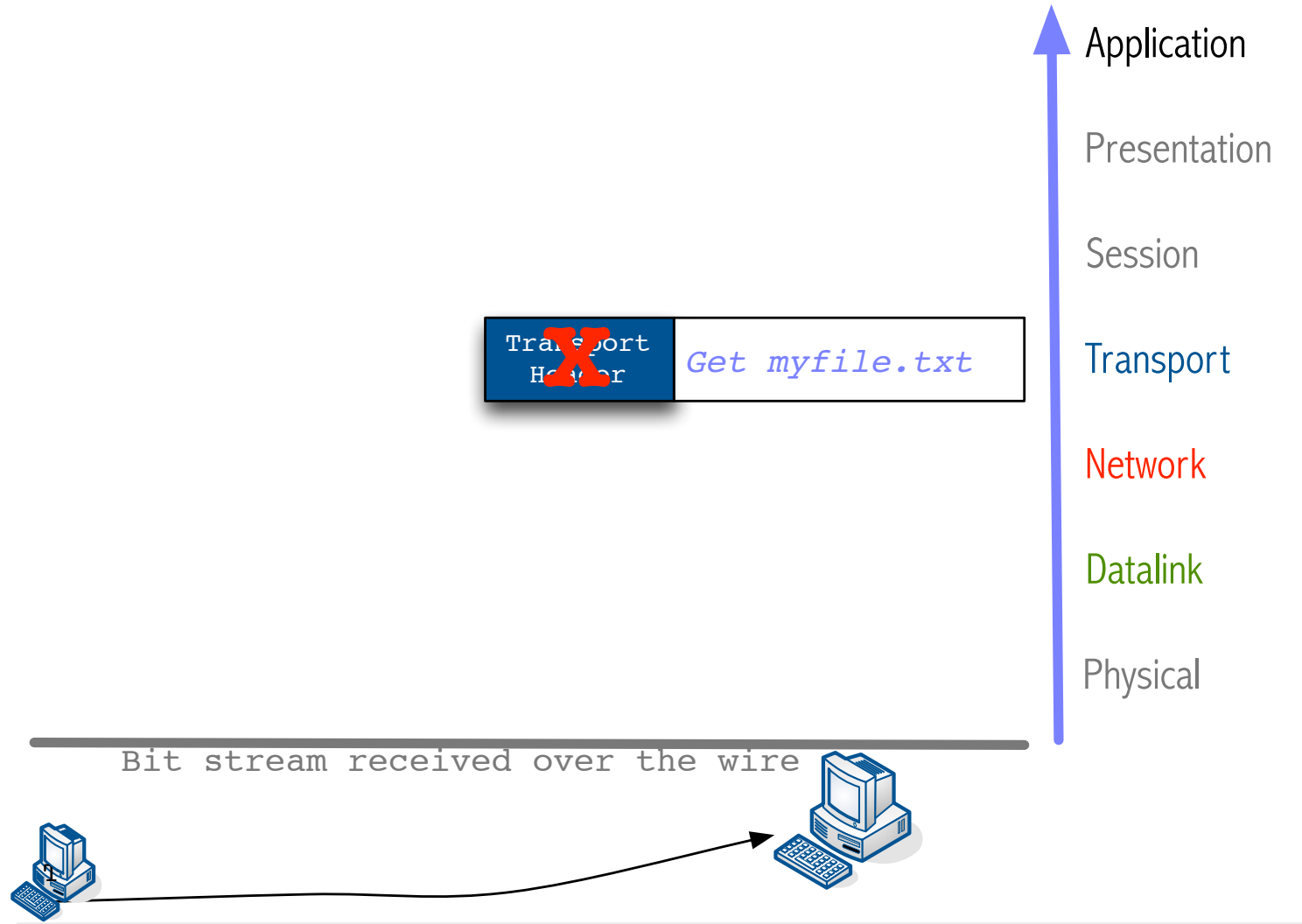


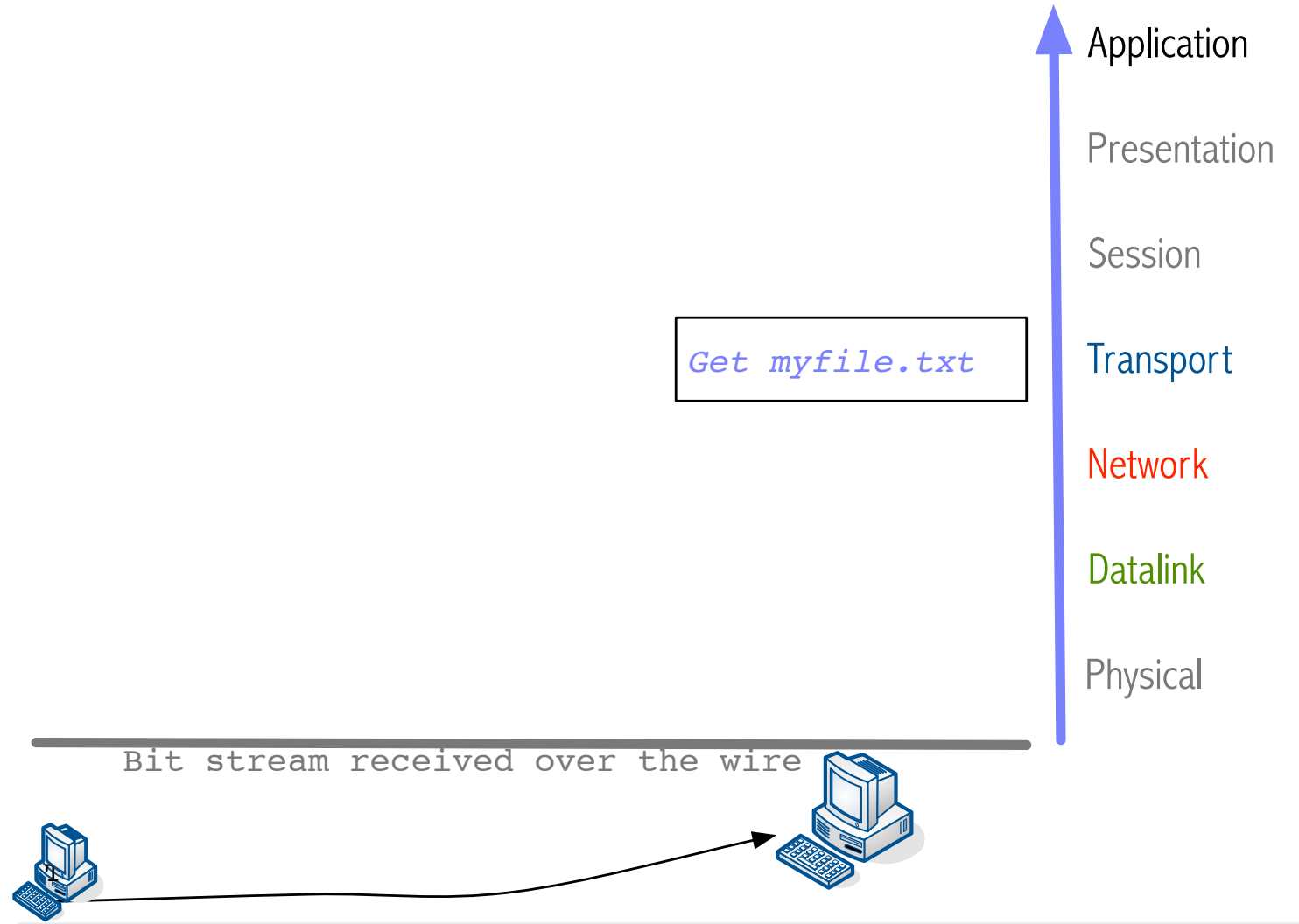


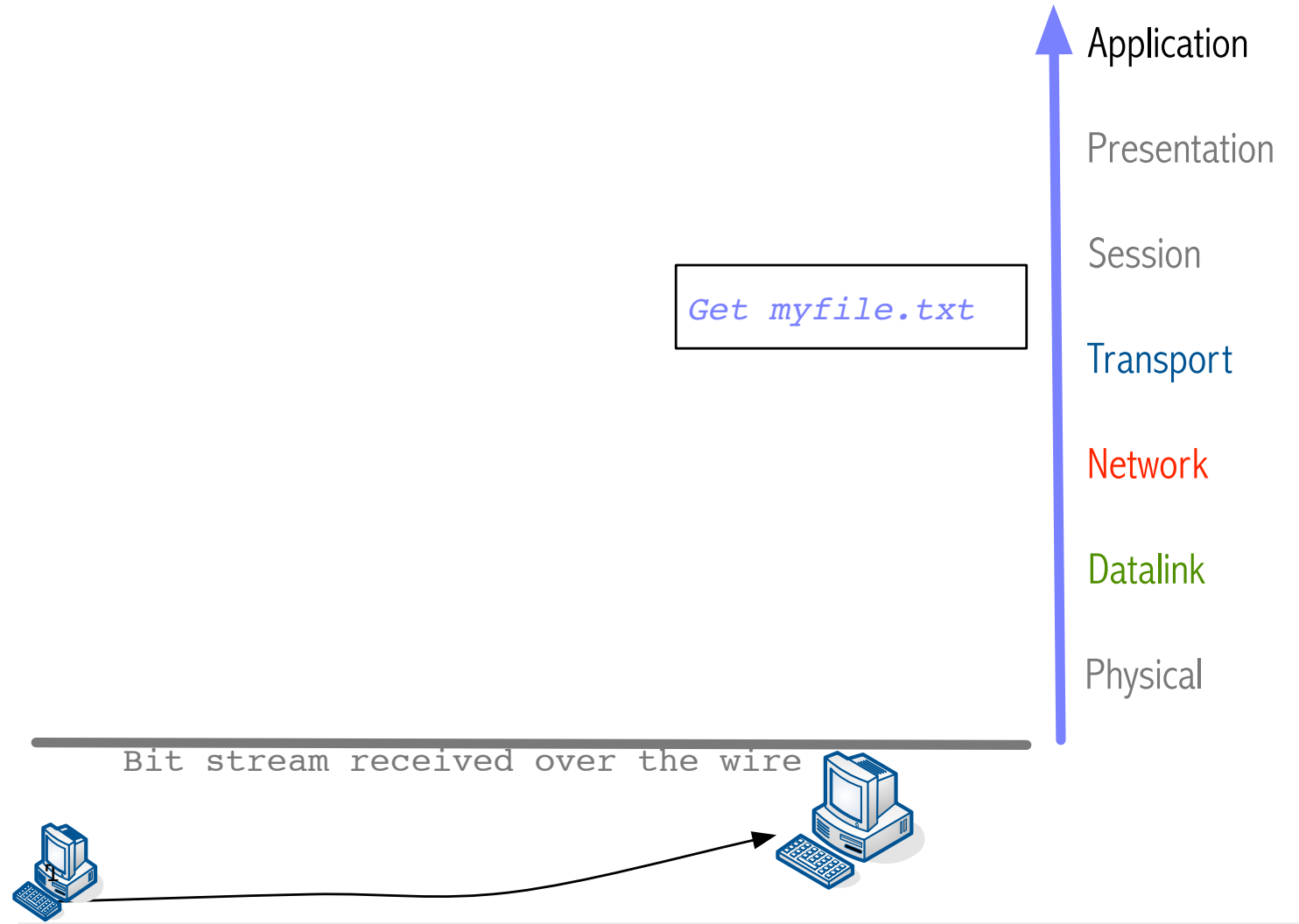




De-encapsulation at layer 4



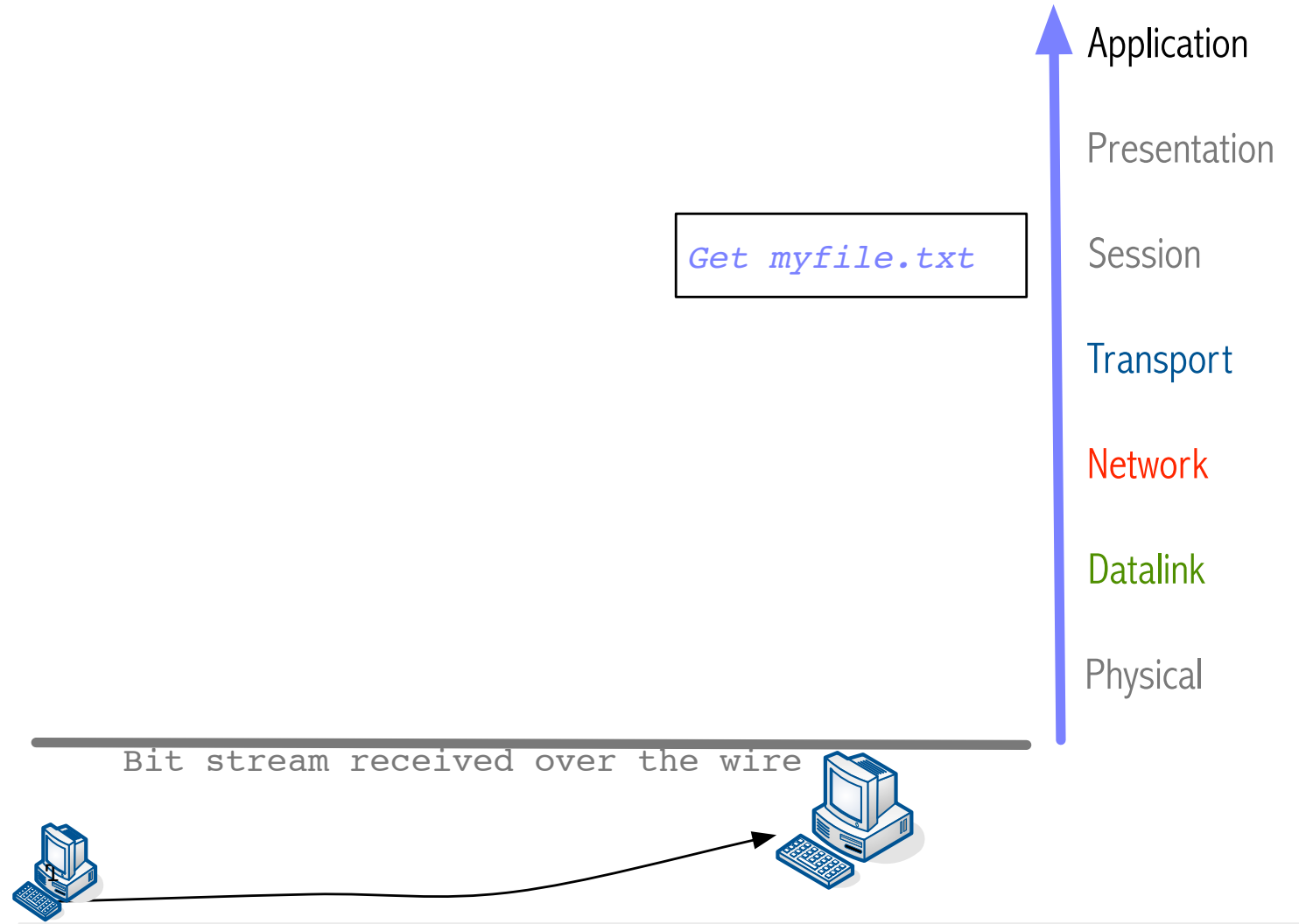




Bit stream received over the wire

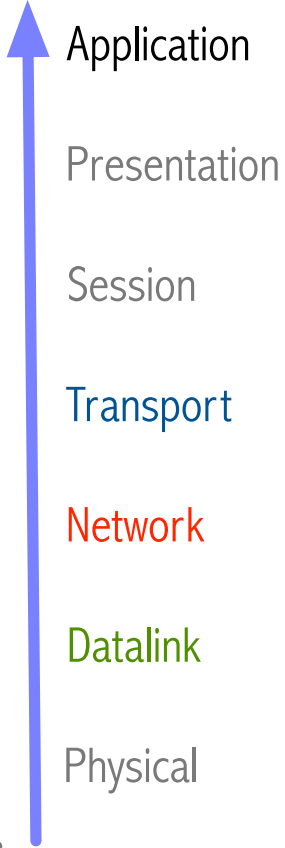
`Get myfile.txt`

- Application
- Presentation
- Session
- Transport
- Network
- Datalink
- Physical

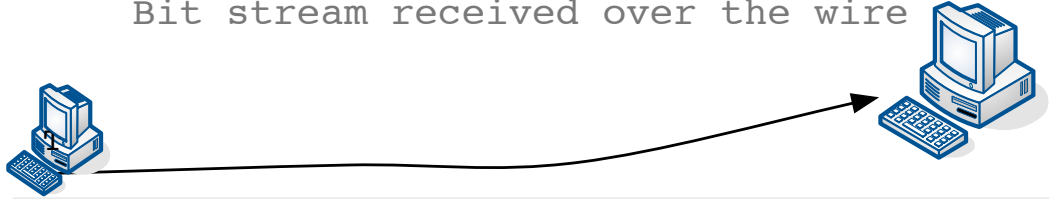




Get myfile.txt

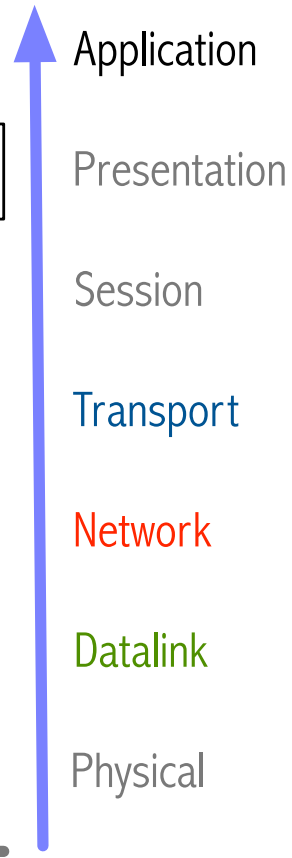


Bit stream received over the wire

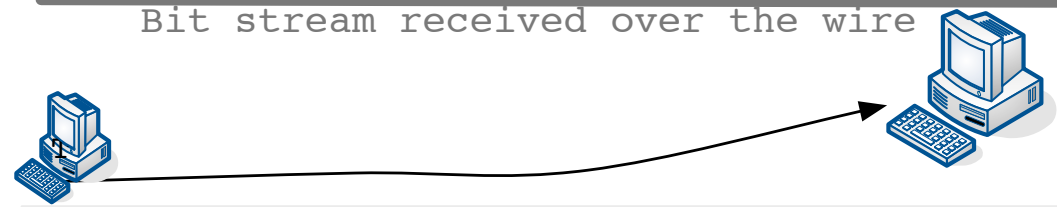




Get myfile.txt



Bit stream received over the wire





Get myfile.txt

Application

Presentation

Session

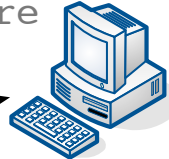
Transport

Network

Datalink

Physical

Bit stream received over the wire





Get myfile.txt

Application

Presentation

Session

Transport

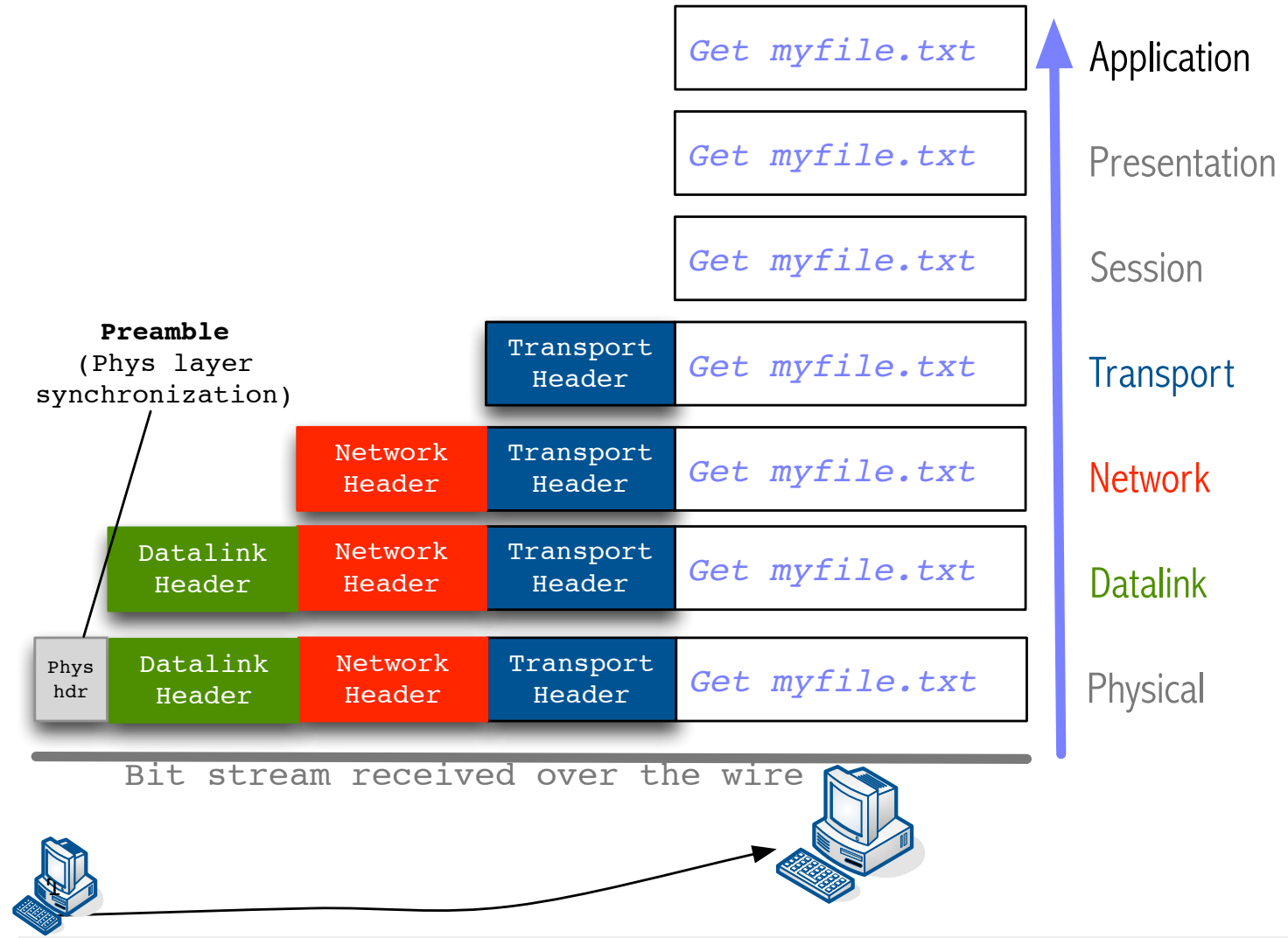
Network

Datalink

Physical

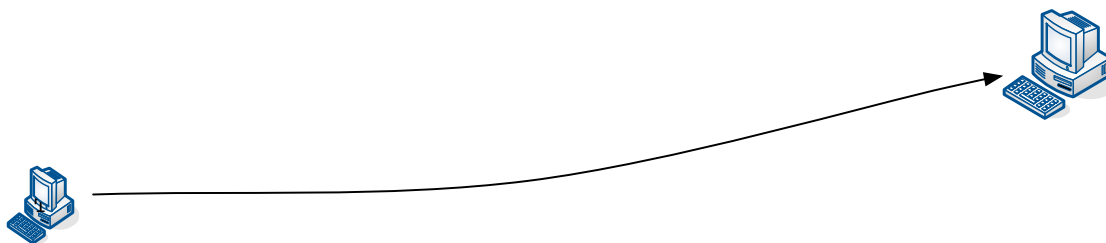
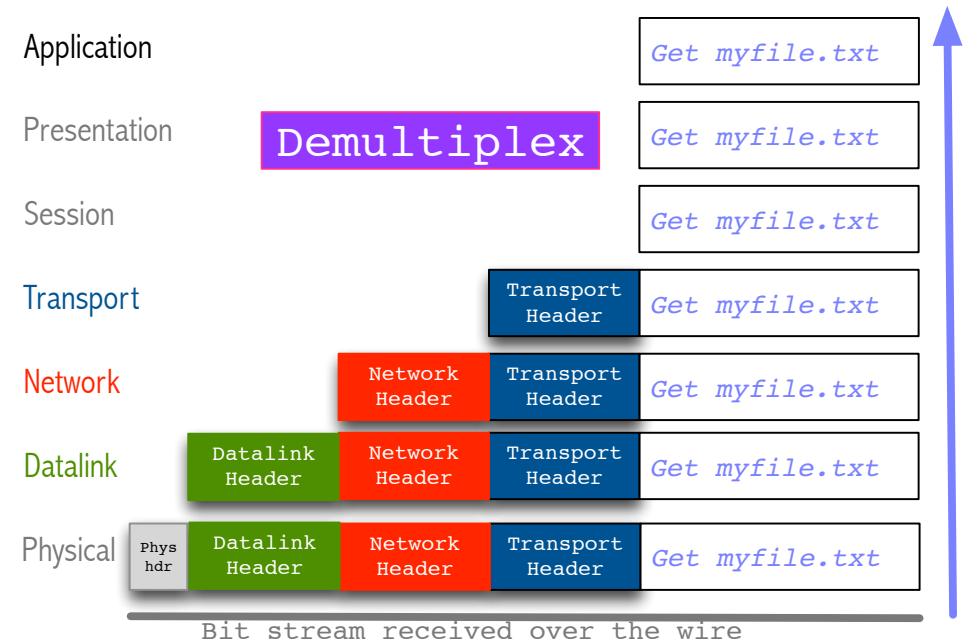
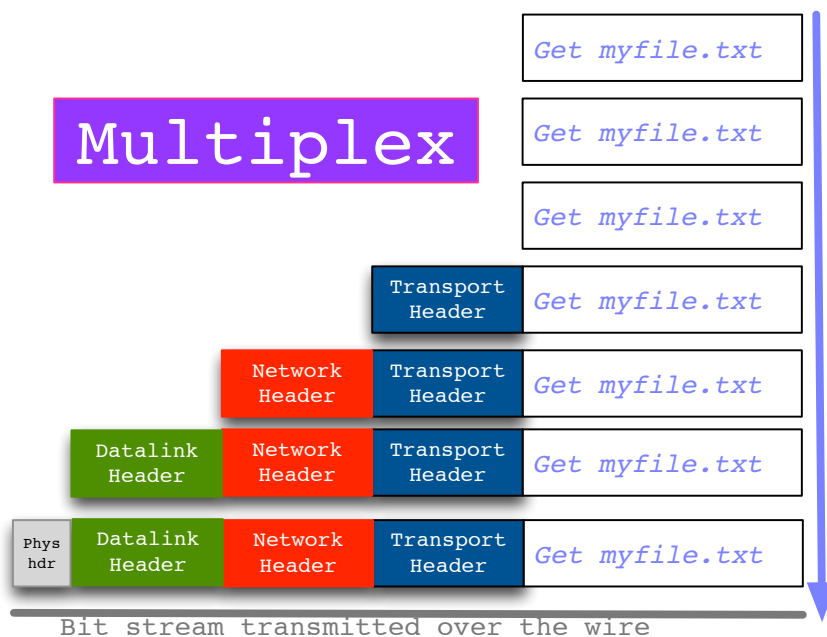
Bit stream received over the wire





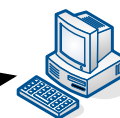
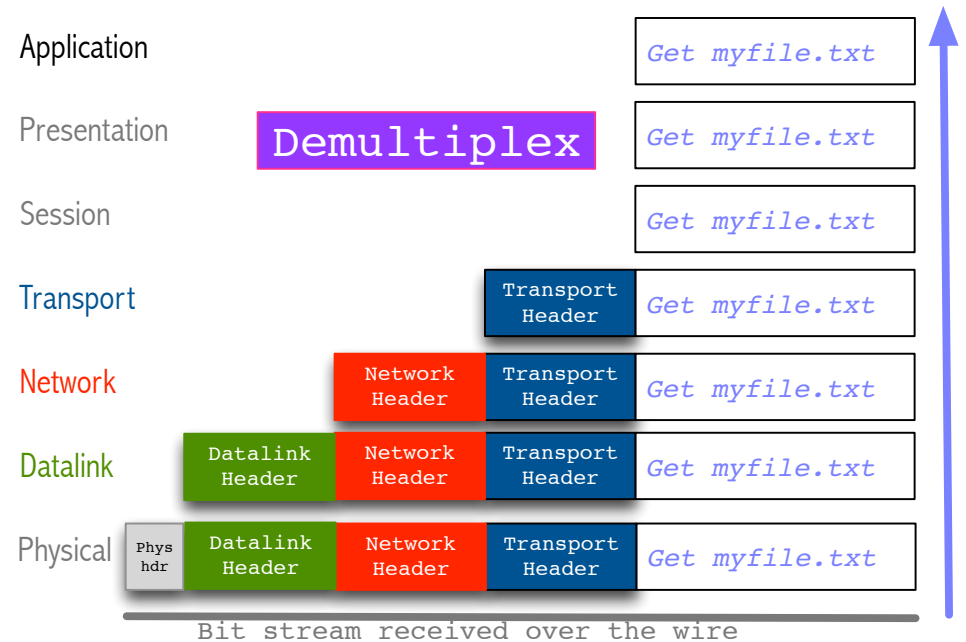
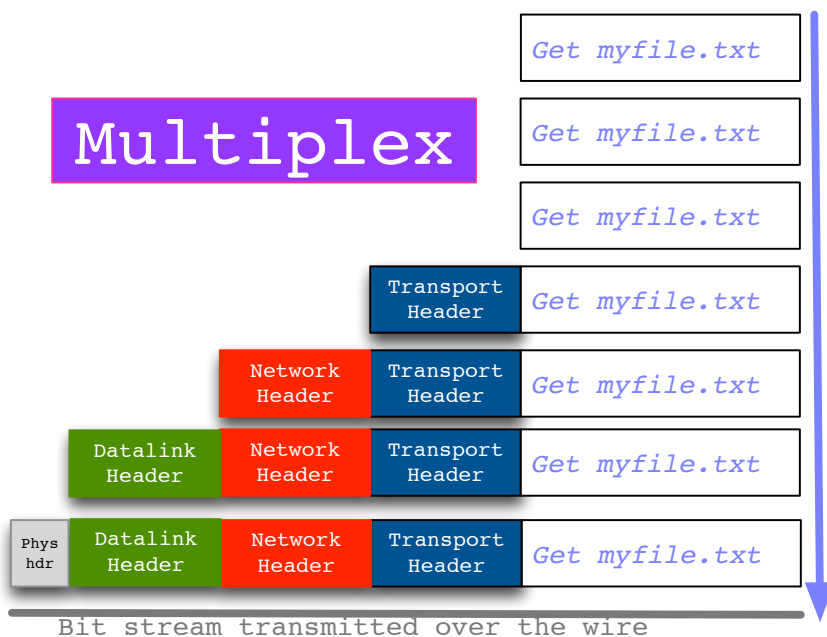
Multiplexing

- Transmitter multiplexes several flows by having each layer add its header which contains addressing information



Demultiplexing

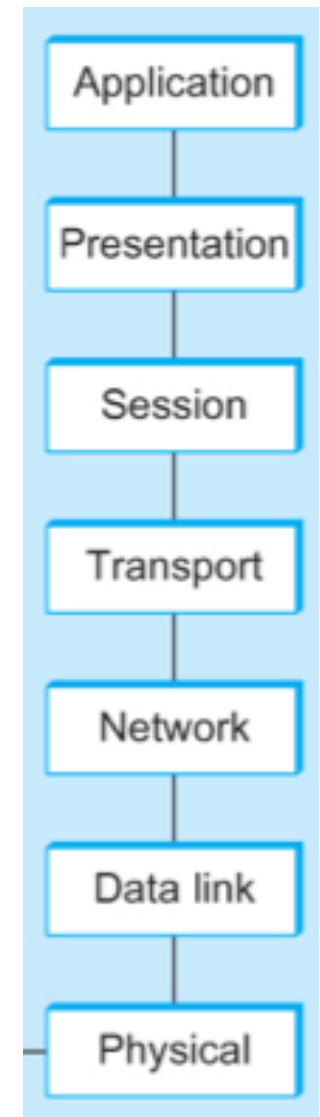
- Receiver demultiplexes several flows by having each layer analyze its header which contains addressing information about the upper-layer protocol that is to receive the payload



Description of Layers

- **Physical Layer**
 - ▣ Handles the transmission of raw bits over a communication link
- **Data Link Layer**
 - ▣ Collects a stream of bits into a larger aggregate called a *frame*
 - ▣ Network adaptor along with device driver in OS implement the protocol in this layer
 - ▣ Frames are actually delivered to hosts
- **Network Layer**
 - ▣ Handles routing among nodes within a packet-switched network
 - ▣ Unit of data exchanged between nodes in this layer is called a *packet*

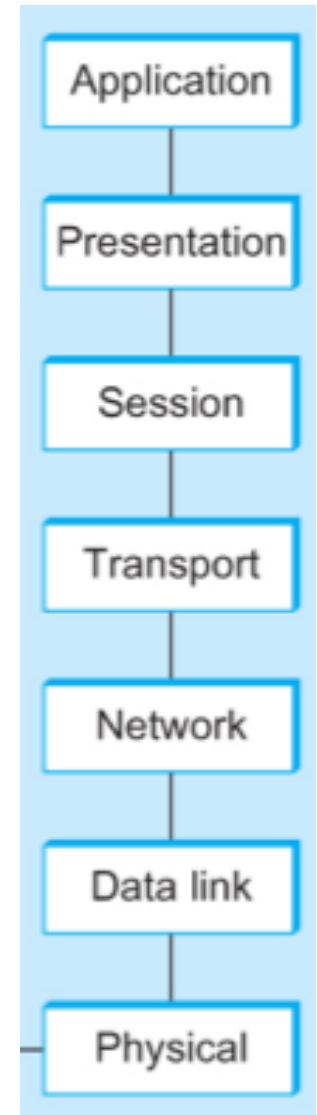
The lower three layers are implemented **on all network nodes**



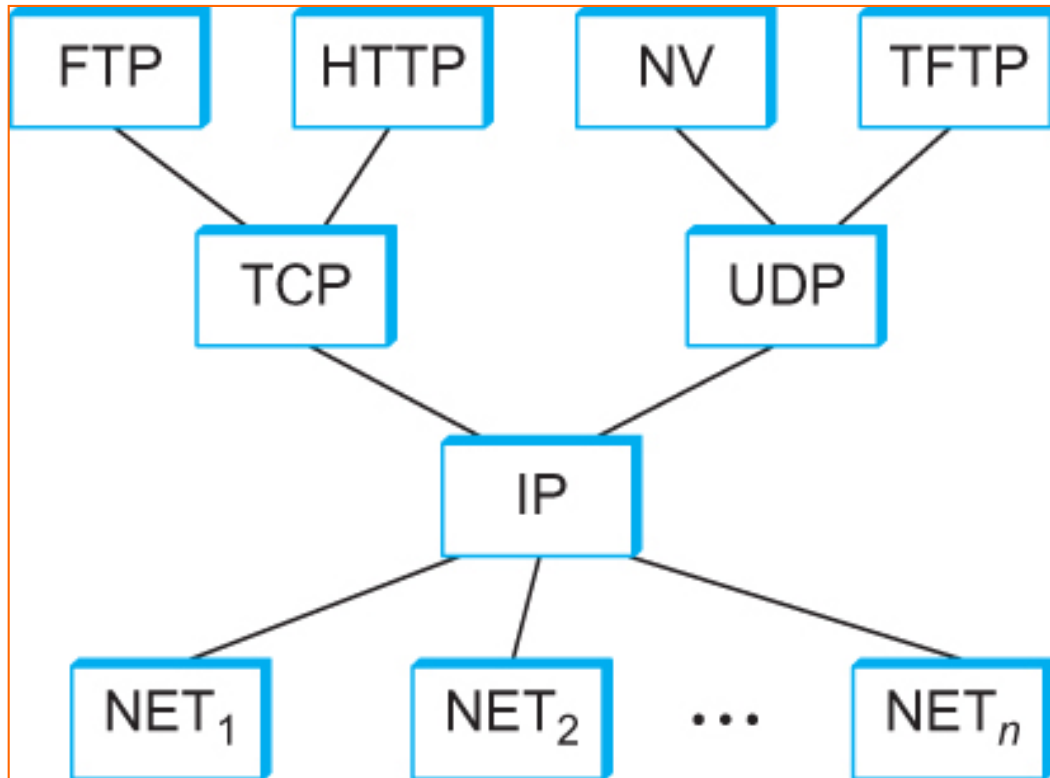
Description of Layers

- **Transport Layer**
 - ▣ Implements a process-to-process channel
 - ▣ Unit of data exchanges in this layer is called a *message*
- **Session Layer**
 - ▣ Provides a name space that is used to tie together the potentially different transport streams that are part of a single application
- **Presentation Layer**
 - ▣ Concerned about the format of data exchanged between peers
- **Application Layer**
 - ▣ Standardize common type of exchanges

The transport layer and the higher layers typically run only **on end-hosts** and not on the intermediate switches and routers

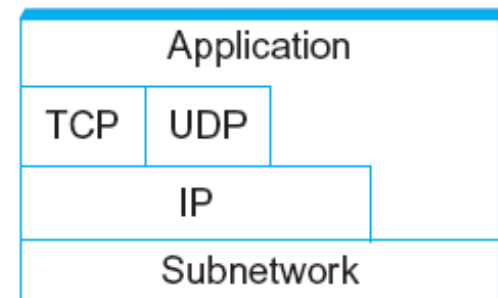


Internet Architecture



Example of an Internet Protocol Graph

TCP/IP architecture!



Alternative view of the Internet architecture. The “Network” layer shown here is sometimes referred to as the “sub-network” or “link” layer.

Internet Architecture

- Defined by **IETF** (Internet Engineering Task Force)
- Three main features
 - ▣ Does **not imply strict layering**. The application is free to bypass the defined transport layers and to directly use IP or other underlying networks
 - ▣ **An hour-glass shape** – wide at the top, narrow in the middle and wide at the bottom. IP serves as the focal point for the architecture
 - ▣ In order for a new protocol to be officially included in the architecture, there needs to be both a protocol specification and at least one (and preferably two) representative implementations of the specification

Network Performance: Networks must be fast

What are the essential network performance metrics: bandwidth and latency

Transmission media AS systems

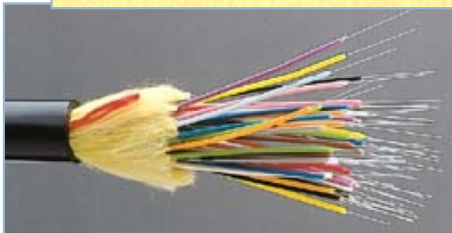
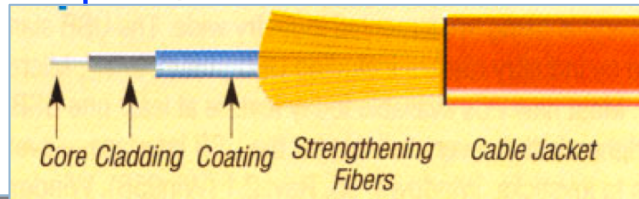


Input signal $f(t)$
Has bandwidth B

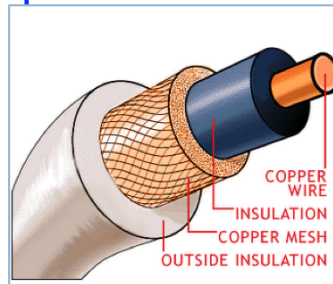
System with $H(\omega)$

Output signal $g(t)$

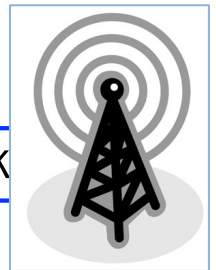
An optical fiber: $B = 100 \text{ GHz}$



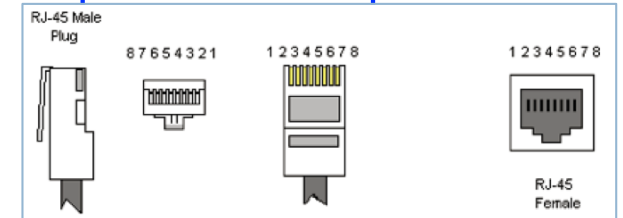
Coaxial cable: $B = 3 \text{ GHz}$



WiFi radio link



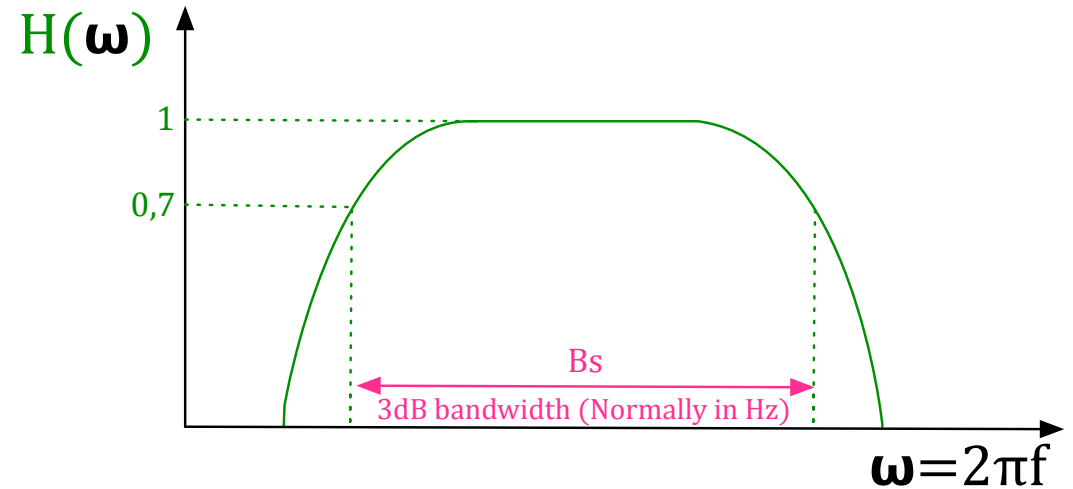
Twisted Pair



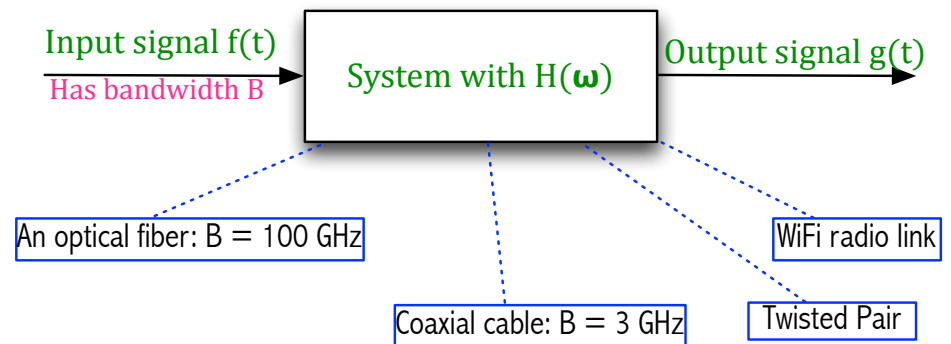
Bandwidth is a property of every system

The width of the accepted frequency band

- Measured in $\text{Hz} = 1/\text{s}$
- In networking bps means bits per second
 - ▣ bits/s
 - ▣ bits · Hz

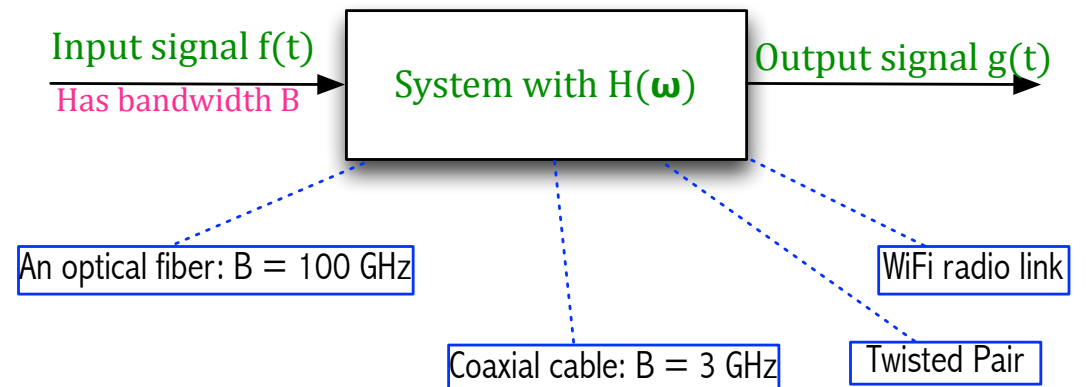
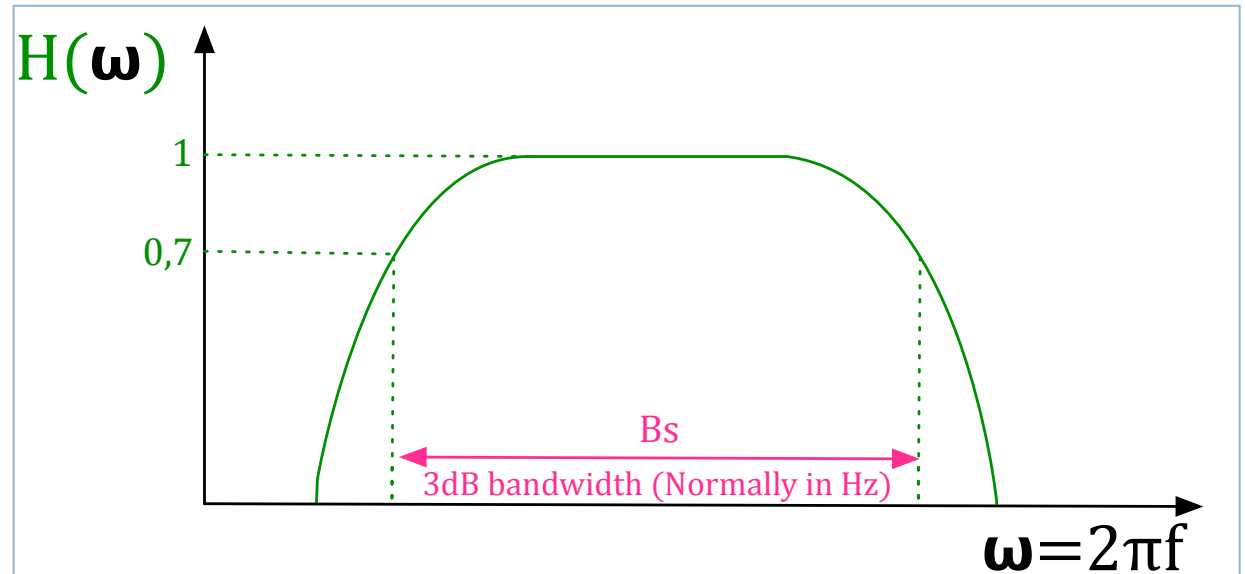


- Higher bandwidth (Hz) allows a higher transmission speed



Bandwidth

and
transmission
speed



If $B > B_s \Rightarrow$ output will have big distortion
If $B_s > B \Rightarrow$ output will have little distortion

- If the band can hold the signal, it passes cleanly on to the output
- B_s in Hz is the inverse of the signaling SPEED in seconds
- High bps needs high bandwidth \Rightarrow high speed demands high bandwidth
- High bandwidth permits high speed transmission

Performance



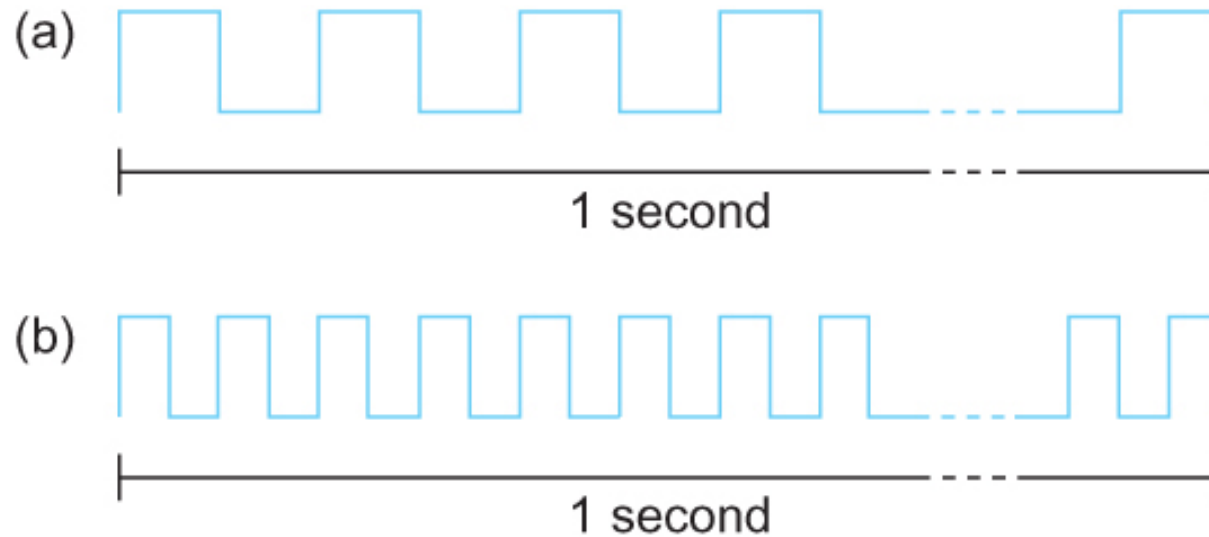
□ Bandwidth

- Width of the frequency band
- Number of bits per second that can be transmitted over a communication link in a given period of time
- The inverse of Bw is the time it takes to transmit one bit
 - $1\text{Mbps} = 1\text{M bits/s} = 10^6 \text{ bits/s} \rightarrow T_{\text{transm}} 1 \text{ bit} = 10^{-6} \text{ s/bit} \cdot 1 \text{ bit} = 1\mu\text{s}$
- Higher bandwidth (speed) means shorter transmission times

□ Multipliers used in expressing speeds (ratios)

- $1\text{Kbps} = 1\text{K bits/s} = 10^3 \text{ bits/s}$
- $1\text{Mbps} = 1\text{M bits/s} = 10^6 \text{ bits/s}$
- $1\text{Gbps} = 1\text{G bits/s} = 10^9 \text{ bits/s}$

Bandwidth



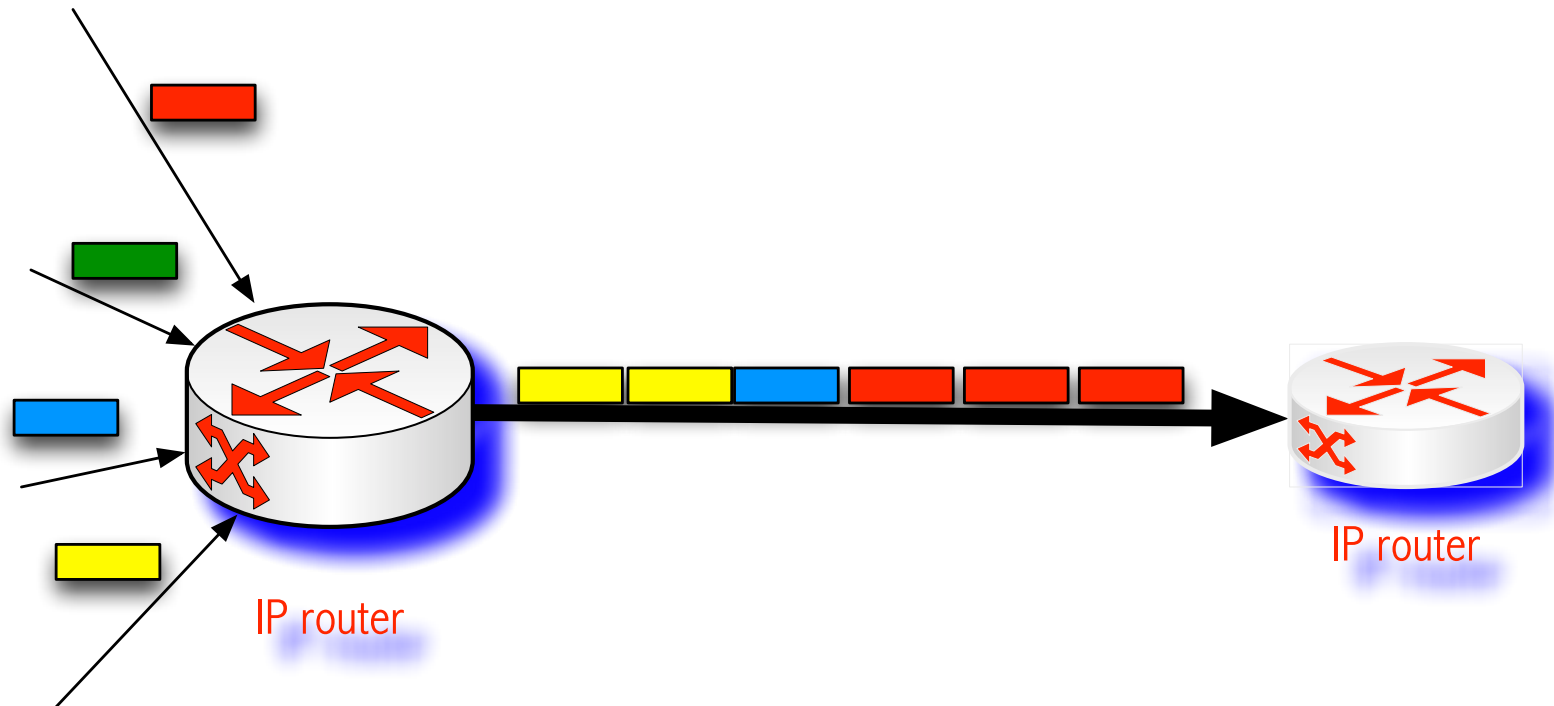
Bits transmitted at a particular bandwidth can be regarded as having some width:

(a) bits transmitted at 1Mbps (each bit $1\ \mu\text{s}$ wide);


(b) bits transmitted at 2Mbps (each bit $0.5\ \mu\text{s}$ wide).

Context for analyzing the components of performance in networks

An IP router is a form of statistical multiplexer



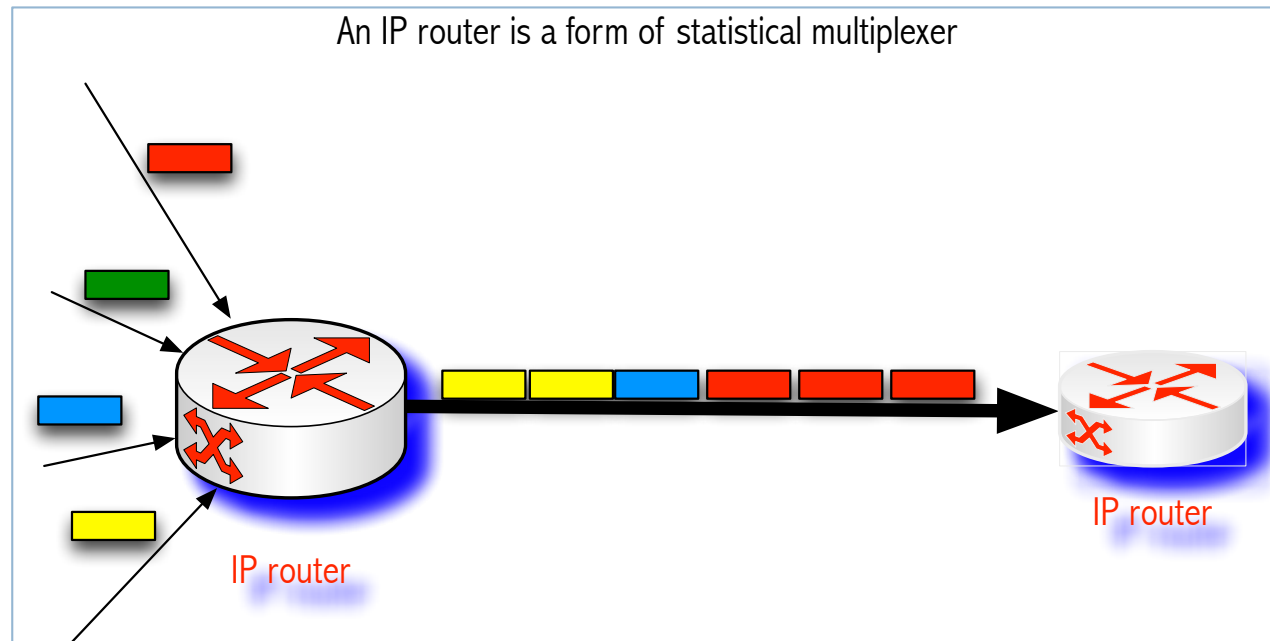
Performance

- 
- Latency = Propagation + transmit + queue
 - Propagation = distance/speed of light
 - Transmit = size/bandwidth

 - One bit transmission => propagation is important
 - Large bytes transmission => bandwidth is important

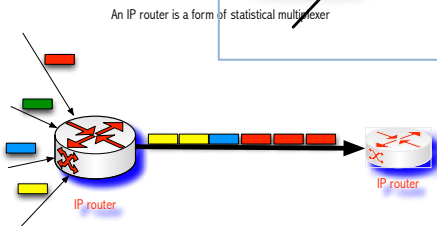
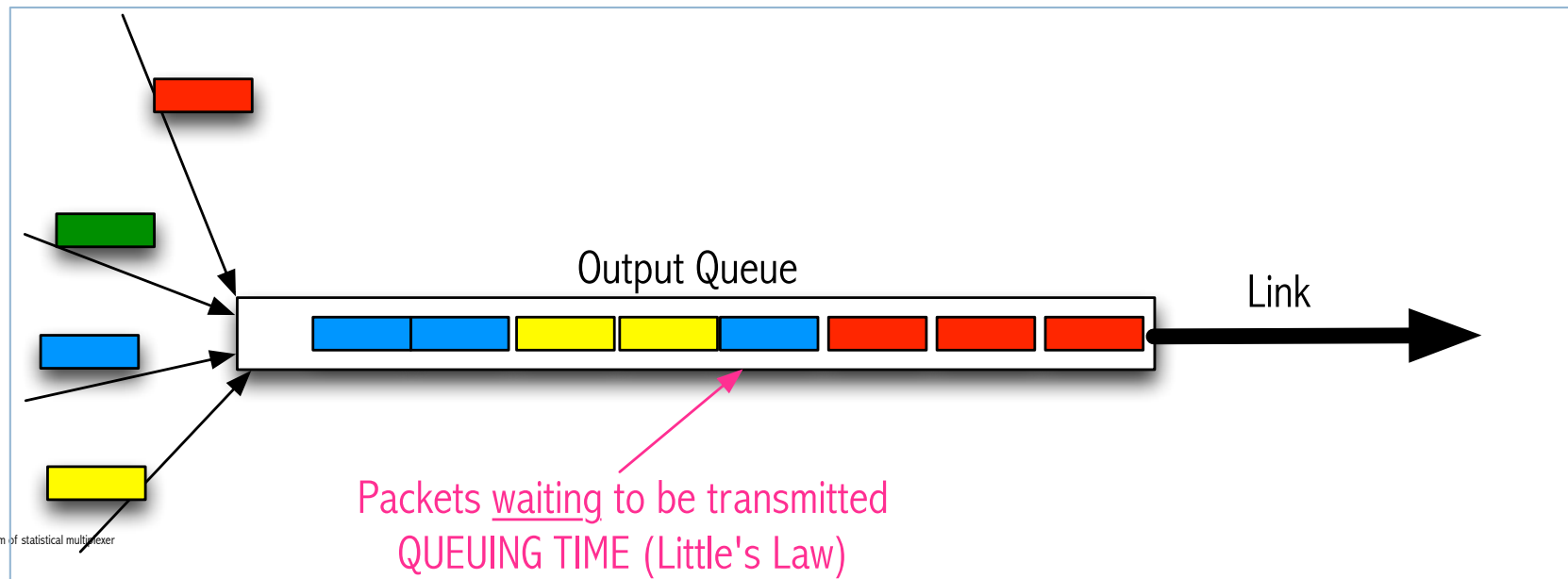
How much does a packet take to be *transferred* to its destination?

- Queuing time
- Transmission time
- Propagation time



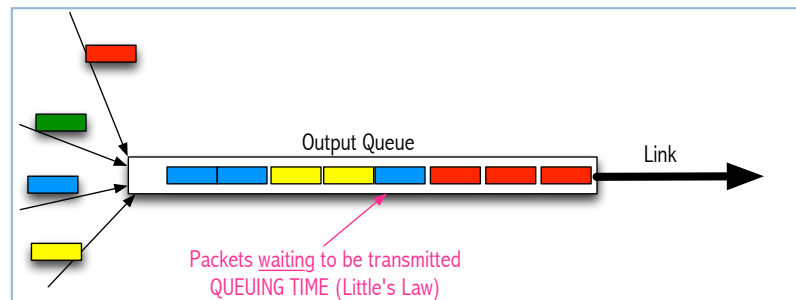
Queuing time

- The router has to enqueue the packets received from the various sources
 - ▣ Otherwise, they will be lost



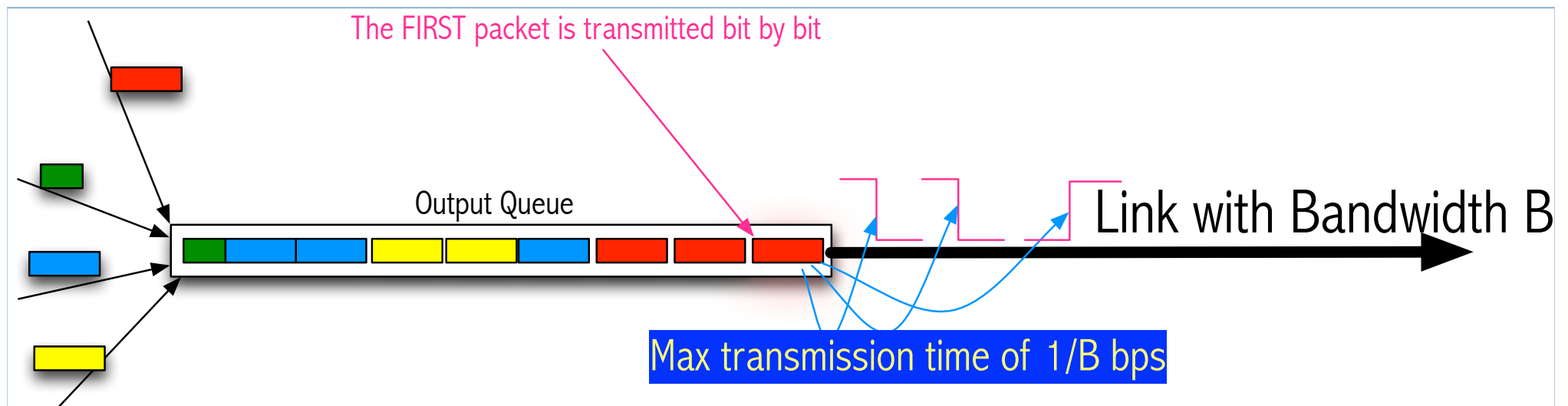
Transmission time

- The **packet** on the **first position in Queue** is transmitted
 - ▣ Packet bits are turned into signals (Electrical, Electromagnetic, Optical)
 - ▣ **1 bit** from packet -> **1 signal waveform**
 - ▣ At which **speed** (bps = bits/s) can this be done?
 - ▣ High Bw => High speed



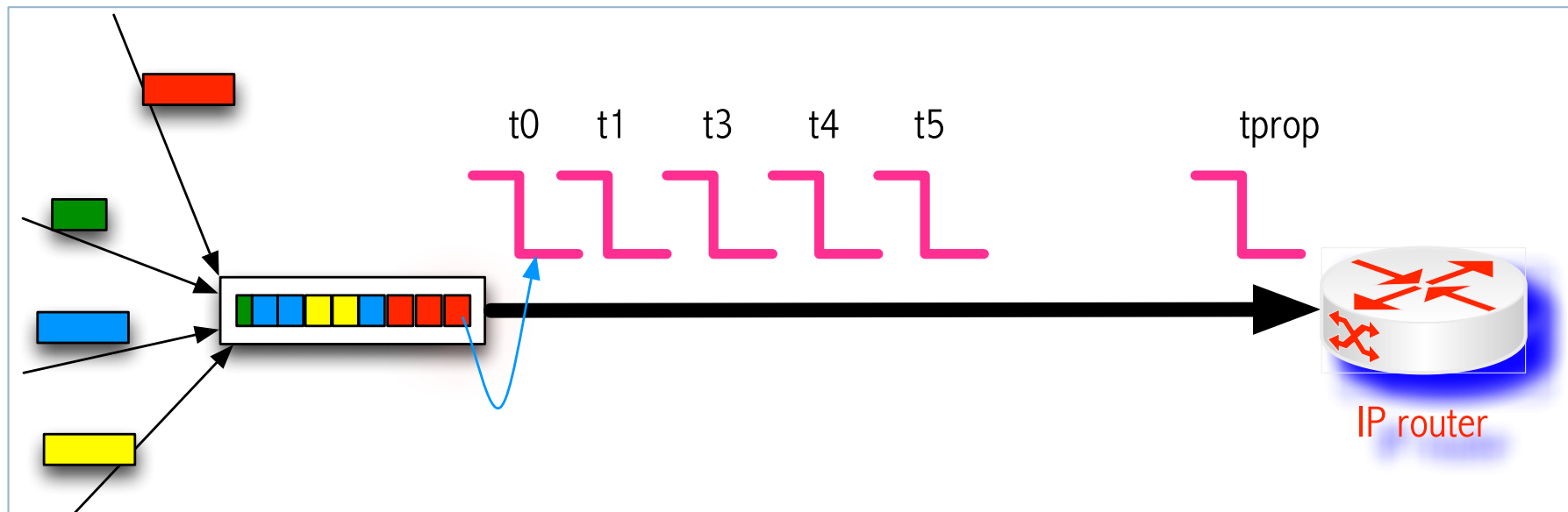
Transmission time

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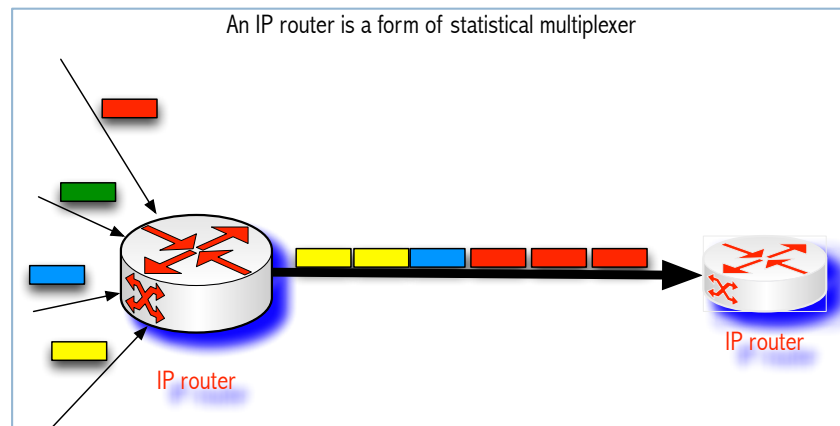
Propagation time (t_{prop})

- Now, each signal must propagate over the wire until it reaches the receiver
 - ▣ Speed of light in empty space $c = 3 \times 10^8$ m/s
 - ▣ Speed of light in other media
 - Copper: $2,3 \times 10^8$ m/s
 - Optical fibers: $2,0 \times 10^8$ m/s



Latency = total time to transfer one packet

- Latency = Propagation + transmit + queue
- Propagation = distance/speed of light
- Transmit = size/bandwidth
- If only one bit is transmitted => propagation is important
 - ▣ Or a small amount of bits
- If the amount of bits transmitted is large => bandwidth is important



Delay x Bandwidth

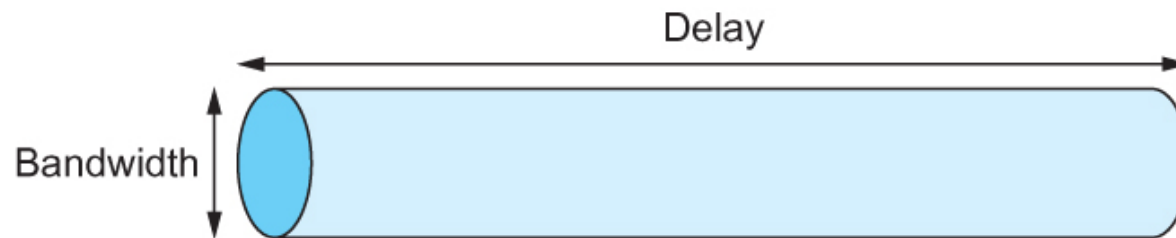
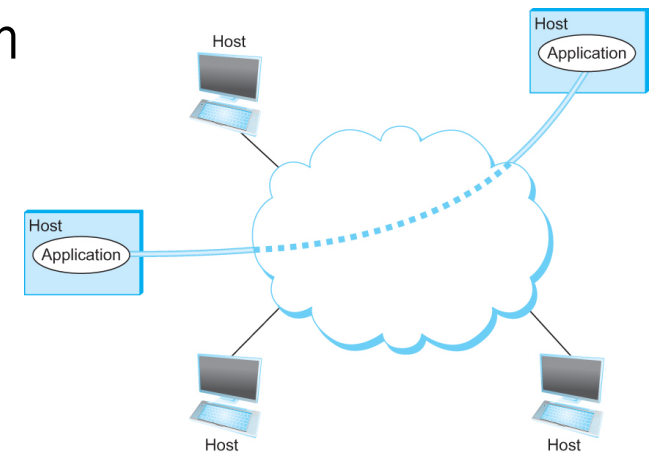
□ We think of the channel between a pair of processes as a hollow pipe

□ Latency (delay) length of the pipe and bandwidth

□ Delay of 50 ms and bandwidth of 45 Mbps


⇒ 50×10^{-3} seconds \times 45×10^6 bits/second

⇒ 2.25×10^6 bits = 280 KB data.



Network as a pipe

Delay X Bandwidth

- 
- Relative importance of **bandwidth** and **latency** depends on application
 - ▣ For large file transfer, bandwidth is critical
 - ▣ For small messages (HTTP, NFS, etc.), latency is critical
 - ▣ Variance in latency (**jitter**) can also affect some applications (*e.g.*, audio/video conferencing)

Delay X Bandwidth

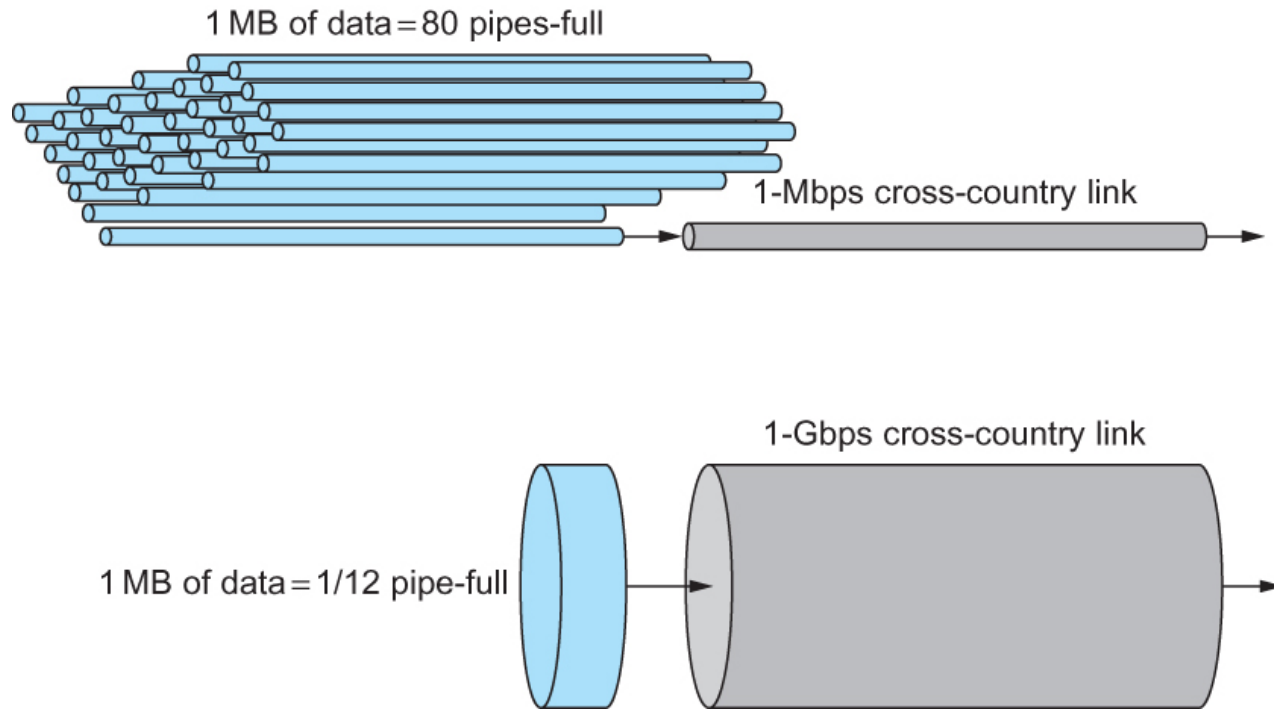
- How many bits the sender must transmit before the first bit arrives at the receiver if the sender keeps the pipe full
- Takes another one-way latency to receive a response from the receiver
- If the sender does not fill the pipe—send a whole delay \times bandwidth product's worth of data before it stops to wait for a signal—the sender will not fully utilize the network

Delay X Bandwidth

- Infinite bandwidth
 - ▣ RTT dominates
 - ▣ $\text{Throughput} = \text{TransferSize} / \text{TransferTime}$
 - ▣ $\text{TransferTime} = \text{RTT} + 1/\text{Bandwidth} \times \text{TransferSize}$


- It's all relative
 - ▣ 1-MB file to 1-Gbps link looks like a 1-KB packet to 1-Mbps link

Relationship between bandwidth and latency



A 1-MB file would fill the 1-Mbps link 80 times,
but only fill the 1-Gbps link 1/12 of one time

Summary

- 
- We have identified what we expect from a computer network
 - We have defined a layered architecture for computer network that will serve as a blueprint for our design
 - We have discussed the socket interface which will be used by applications for invoking the services of the network subsystem
 - We have discussed two performance metrics using which we can analyze the performance of computer networks

Exercises

What are the essential network performance metrics: bandwidth and latency